Featherweight OCL
A Proposal for a Machine-Checked Formal Semantics for OCL 2.5

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Abstract

The Unified Modeling Language (UML) is one of the few modeling languages that is widely used in industry. While UML is mostly known as diagrammatic modeling language (e.g., visualizing class models), it is complemented by a textual language, called Object Constraint Language (OCL). OCL is a textual annotation language, based on a three-valued logic, that turns UML into a formal language. Unfortunately the semantics of this specification language, captured in the “Annex A” of the OCL standard, leads to different interpretations of corner cases. Many of these corner cases had been subject to formal analysis since more than ten years.

The situation complicated when with version 2.3 the OCL was aligned with the latest version of UML: this led to the extension of the three-valued logic by a second exception element, called null. While the first exception element invalid has a strict semantics, null has a non strict semantic interpretation. These semantic difficulties lead to remarkable confusion for implementors of OCL compilers and interpreters.

In this paper, we provide a formalization of the core of OCL in HOL. It provides denotational definitions, a logical calculus and operational rules that allow for the execution of OCL expressions by a mixture of term rewriting and code compilation. Our formalization reveals several inconsistencies and contradictions in the current version of the OCL standard. They reflect a challenge to define and implement OCL tools in a uniform manner. Overall, this document is intended to provide the basis for a machine-checked text “Annex A” of the OCL standard targeting at tool implementors.
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Part I.

Introduction
1. Motivation

The Unified Modeling Language (UML) \[31, 32\] is one of the few modeling languages that is widely used in industry. UML is defined, in an open process, by the Object Management Group (OMG), i.e., an industry consortium. While UML is mostly known as diagrammatic modeling language (e.g., visualizing class models), it also comprises a textual language, called Object Constraint Language (OCL) \[33\]. OCL is a textual annotation language, originally conceived as a three-valued logic, that turns substantial parts of UML into a formal language. Unfortunately the semantics of this specification language, captured in the “Annex A” (originally, based on the work of Richters \[35\]) of the OCL standard leads to different interpretations of corner cases. Many of these corner cases had been subject to formal analysis since more than nearly fifteen years (see, e.g., \[5, 11, 19, 22, 26\]).

At its origins \[28, 35\], OCL was conceived as a strict semantics for undefinedness (e.g., denoted by the element \texttt{invalid}\[1\], with the exception of the logical connectives of type \texttt{Boolean} that constitute a three-valued propositional logic. At its core, OCL comprises four layers:

1. Operators (e.g., \texttt{and}, \texttt{+}) on built-in data structures such as \texttt{Boolean}, \texttt{Integer}, or typed sets (\texttt{Set(_)}).
2. Operators on the user-defined data model (e.g., defined as part of a UML class model) such as accessors, type casts and tests.
3. Arbitrary, user-defined, side-effect-free methods,
4. Specification for invariants on states and contracts for operations to be specified via pre- and post-conditions.

Motivated by the need for aligning OCL closer with UML, recent versions of the OCL standard \[30, 33\] added a second exception element. While the first exception element \texttt{invalid} has a strict semantics, \texttt{null} has a non strict semantic interpretation. Unfortunately, this extension results in several inconsistencies and contradictions. These problems are reflected in difficulties to define interpreters, code-generators, specification animators or theorem provers for OCL in a uniform manner and resulting incompatibilities of various tools.

For the OCL community, the semantics of \texttt{invalid} and \texttt{null} as well as many related issues resulted in the challenge to define a consistent version of the OCL standard that is well aligned with the recent developments of the UML. A syntactical and semantical

\[1\] In earlier versions of the OCL standard, this element was called \texttt{OclUndefined}. 

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consistent standard requires a major revision of both the informal and formal parts of
the standard. To discuss the future directions of the standard, several OCL experts
met in November 2013 in Aachen to discuss possible mid-term improvements of OCL,
strategies of standardization of OCL within the OMG, and a vision for possible long-term
developments of the language [15]. During this meeting, a Request for Proposals (RFP)
for OCL 2.5 was finalized and meanwhile proposed. In particular, this RFP requires
that the future OCL 2.5 standard document shall be generated from a machine-checked
source. This will ensure

- the absence of syntax errors,
- the consistency of the formal semantics,
- a suite of corner-cases relevant for OCL tool implementors.

In this document, we present a formalization using Isabelle/HOL [27] of a core
language of OCL. The semantic theory, based on a “shallow embedding”, is called
Featherweight OCL, since it focuses on a formal treatment of the key-elements of the
language (rather than a full treatment of all operators and thus, a “complete” implemen-
tation). In contrast to full OCL, it comprises just the logic captured in Boolean, the basic
data type Integer, the collection type Set, as well as the generic construction principle
of class models, which is instantiated and demonstrated for two examples (an automated
support for this type-safe construction is again out of the scope of Featherweight OCL).
This formal semantics definition is intended to be a proposal for the standardization
process of OCL 2.5, which should ultimately replace parts of the mandatory part of the
standard document [33] as well as replace completely its informative “Annex A.”
2. Background

2.1. A Guided Tour Through UML/OCL

The Unified Modeling Language (UML) [31, 32] comprises a variety of model types for describing static (e.g., class models, object models) and dynamic (e.g., state-machines, activity graphs) system properties. One of the more prominent model types of the UML is the class model (visualized as class diagram) for modeling the underlying data model of a system in an object-oriented manner. As a running example, we model a part of a conference management system.

Such a system usually supports the conference organizing process, e.g., creating a conference Website, reviewing submissions, registering attendees, organizing the different sessions and tracks, and indexing and producing the resulting proceedings. In this example, we constrain ourselves to the process of organizing conference sessions; Figure 2.1 shows the class model. We model the hierarchy of roles of our system as a hierarchy of classes (e.g., Hearer, Speaker, or Chair) using an inheritance relation (also called generalization). In particular, inheritance establishes a subtyping relationship, i.e., every Speaker (subclass) is also a Hearer (superclass).

A class does not only describe a set of instances (called objects), i.e., record-like data consisting of attributes such as name of class Session, but also operations defined over them. For example, for the class Session, representing a conference session, we model an operation findRole(p:Person):Role that should return the role of a Person in the context of a specific session; later, we will describe the behavior of this operation in more detail using UML. In the following, the term object describes a (run-time) instance of a class or one of its subclasses.

Figure 2.1.: A simple UML class model representing a conference system for organizing conference sessions: persons can participate, in different roles, in a session.
Relations between classes (called associations in UML) can be represented in a class diagram by connecting lines, e.g., Participant and Session or Person and Role. Associations may be labeled by a particular constraint called multiplicity, e.g., $0..*$ or $0..1$, which means that in a relation between participants and sessions, each Participant object is associated to at most one Session object, while each Session object may be associated to arbitrarily many Participant objects. Furthermore, associations may be labeled by projection functions like person and role; these implicit function definitions allow for OCL-expressions like self.person, where self is a variable of the class Role. The expression self.person denotes persons being related to the specific object self of type role. A particular feature of the UML are association classes (Participant in our example) which represent a concrete tuple of the relation within a system state as an object; i.e., associations classes allow also for defining attributes and operations for such tuples. In a class diagram, association classes are represented by a dotted line connecting the class with the association. Associations classes can take part in other associations. Moreover, UML supports also n-ary associations (not shown in our example).

We refine this data model using the Object Constraint Language (OCL) for specifying additional invariants, preconditions and postconditions of operations. For example, we specify that objects of the class Person are uniquely determined by the value of the name attribute and that the attribute name is not equal to the empty string (denoted by ‘ ’):

```ocl
class Person
  inv: name <> '' and Person::allInstances()->isUnique(p: Person | p.name)
end class Person
```

Moreover, we specify that every session has exactly one chair by the following invariant (called onlyOneChair) of the class Session:

```ocl
class Session
  inv onlyOneChair: self.participants->one(p: Participant | p.role.oclIsTypeOf(Chair))
end class Session
```

where p.role.oclIsTypeOf(Chair) evaluates to true, if p.role is of dynamic type Chair. Besides the usual static types (i.e., the types inferred by a static type inference), objects in UML and other object-oriented languages have a second dynamic type concept. This is a consequence of a family of casting functions (written $o_C$ for an object $o$ into another class type $C$) that allows for converting the static type of objects along the class hierarchy. The dynamic type of an object can be understood as its “initial static type” and is unchanged by casts. We complete our example by describing the behavior of the operation findRole as follows:

```ocl
class Session::findRole(person: Person): Role
  pre: self.participates.person->includes(person)
  post: result=self.participants->one(p: Participant | p.person = person ).role
       and self.participants = self.participants@pre
       and self.name = self.name@pre
end class Session::findRole
```

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where in post-conditions, the operator @pre allows for accessing the previous state.

In UML, classes can contain attributes of the type of the defining class. Thus, UML can represent (mutually) recursive datatypes. Moreover, OCL introduces also recursively specified operations.

A key idea of defining the semantics of UML and extensions like SecureUML [12] is to translate the diagrammatic UML features into a combination of more elementary features of UML and OCL expressions [21]. For example, associations are usually represented by collection-valued class attributes together with OCL constraints expressing the multiplicity. Thus, having a semantics for a subset of UML and OCL is tantamount for the foundation of the entire method.

2.2. Formal Foundation

2.2.1. Isabelle

Isabelle [27] is a generic theorem prover. New object logics can be introduced by specifying their syntax and natural deduction inference rules. Among other logics, Isabelle supports first-order logic, Zermelo-Fraenkel set theory and the instance for Church’s higher-order logic (HOL).

Isabelle’s inference rules are based on the built-in meta-level implication \( \Rightarrow \) allowing to form constructs like \( A_1 \Rightarrow \cdots \Rightarrow A_n \Rightarrow A_{n+1} \), which are viewed as a rule of the form “from assumptions \( A_1 \) to \( A_n \), infer conclusion \( A_{n+1} \)” and which is written in Isabelle as

\[
\langle A_1; \ldots; A_n \rangle \Rightarrow A_{n+1}
\]

or, in mathematical notation,

\[
\frac{A_1 \cdots A_n}{A_{n+1}}\quad (2.1)
\]

The built-in meta-level quantification \( \bigwedge x. x \) captures the usual side-constraints “\( x \) must not occur free in the assumptions” for quantifier rules; meta-quantified variables can be considered as “fresh” free variables. Meta-level quantification leads to a generalization of Horn-clauses of the form:

\[
\bigwedge x_1, \ldots, x_m. \langle A_1; \ldots; A_n \rangle \Rightarrow A_{n+1} \quad (2.2)
\]

Isabelle supports forward- and backward reasoning on rules. For backward-reasoning, a proof-state can be initialized and further transformed into others. For example, a proof of \( \phi \), using the Isar [38] language, will look as follows in Isabelle:

```
lemma label: \( \phi \)
  apply(case_tac)
  apply(simp_all)
done
```

(2.3)

This proof script instructs Isabelle to prove \( \phi \) by case distinction followed by a simplification of the resulting proof state. Such a proof state is an implicitly conjoint sequence.
of generalized Horn-clauses (called subgoals) $\phi_1, \ldots, \phi_n$ and a goal $\phi$. Proof states were usually denoted by:

\[
\text{label : } \phi \\
1. \phi_1 \\
\vdots \\
n. \phi_n
\]

Subgoals and goals may be extracted from the proof state into theorems of the form $[\phi_1; \ldots; \phi_n] \Rightarrow \phi$ at any time; this mechanism helps to generate test theorems. Further, Isabelle supports meta-variables (written $?x, ?y, \ldots$), which can be seen as “holes in a term” that can still be substituted. Meta-variables are instantiated by Isabelle’s built-in higher-order unification.

2.2.2. Higher-order Logic (HOL)

Higher-order logic (HOL) \cite{HOL} is a classical logic based on a simple type system. It provides the usual logical connectives like $\land$, $\lor$, $\Rightarrow$, $\neg$ as well as the object-logical quantifiers $\forall x. P x$ and $\exists x. P x$; in contrast to first-order logic, quantifiers may range over arbitrary types, including total functions $f : \alpha \Rightarrow \beta$. HOL is centered around extensional equality $=_e : \alpha \Rightarrow \alpha \Rightarrow \text{bool}$. HOL is more expressive than first-order logic, since, e.g., induction schemes can be expressed inside the logic. Being based on some polymorphically typed \(\lambda\)-calculus, HOL can be viewed as a combination of a programming language like SML or Haskell and a specification language providing powerful logical quantifiers ranging over elementary and function types.

Isabelle/HOL is a logical embedding of HOL into Isabelle. The (original) simple-type system underlying HOL has been extended by Hindley-Milner style polymorphism with type-classes similar to Haskell. While Isabelle/HOL is usually seen as proof assistant, we use it as symbolic computation environment. Implementations on top of Isabelle/HOL can re-use existing powerful deduction mechanisms such as higher-order resolution, tableau-based reasoners, rewriting procedures, Presburger arithmetic, and via various integration mechanisms, also external provers such as Vampire \cite{Vampire} and the SMT-solver Z3 \cite{Z3}.

Isabelle/HOL offers support for a particular methodology to extend given theories in a logically safe way: A theory-extension is conservative if the extended theory is consistent provided that the original theory was consistent. Conservative extensions can be constant definitions, type definitions, datatype definitions, primitive recursive definitions and wellfounded recursive definitions.

For instance, the library includes the type constructor $\tau_\bot := \bot \ | \ \langle \rangle : \alpha$ that assigns to each type $\tau$ a type $\tau_\bot$ disjointly extended by the exceptional element $\bot$. The function $\langle \rangle_\tau : \alpha_\bot \rightarrow \alpha$ is the inverse of $\langle \rangle_\tau$ (unspecified for $\bot$). Partial functions $\alpha \Rightarrow \beta$ are defined as functions $\alpha \Rightarrow \beta_\bot$ supporting the usual concepts of domain (dom $\bot$) and range (ran $\bot$).

As another example of a conservative extension, typed sets were built in the Isabelle libraries conservatively on top of the kernel of HOL as functions to bool; consequently,
the constant definitions for membership is as follows:

\[
\text{types} \quad \alpha \text{ set} = \alpha \Rightarrow \text{bool} \\
\text{definition} \quad \text{Collect} : (\alpha \Rightarrow \text{bool}) \Rightarrow \alpha \text{ set} \quad \text{— set comprehension} \\
\text{where} \quad \text{Collect } S \equiv S \\
\text{definition} \quad \text{member} : \alpha \Rightarrow \alpha \Rightarrow \text{bool} \quad \text{— membership test} \\
\text{where} \quad \text{member } s S \equiv S s
\]

Isabelle’s syntax engine is instructed to accept the notation \( \{ x \mid P \} \) for \( \text{Collect } \lambda x. P \) and the notation \( s \in S \) for \( \text{member } s S \). As can be inferred from the example, constant definitions are axioms that introduce a fresh constant symbol by some closed, non-recursive expressions; this type of axiom is logically safe since it works like an abbreviation. The syntactic side conditions of this axiom are mechanically checked, of course. It is straightforward to express the usual operations on sets like \( \cup \), \( \cap \), \( \cdot \) as conservative extensions, too, while the rules of typed set theory were derived by proofs from these definitions.

Similarly, a logical compiler is invoked for the following statements introducing the types option and list:

\[
\text{datatype} \quad \text{option} = \text{None} | \text{Some } \alpha \\
\text{datatype} \quad \alpha \text{ list} = \text{Nil} | \text{Cons } a l
\]

Here, \( [] \) or \( a # l \) are an alternative syntax for Nil or Cons \( a l \); moreover, \( [a,b,c] \) is defined as alternative syntax for \( a # b # c # [] \). These (recursive) statements were internally represented in by internal type and constant definitions. Besides the constructors None, Some, \( [] \) and Cons, there is the match operation

\[
\text{case } x \text{ of } \text{None } \Rightarrow F \mid \text{Some } a \Rightarrow G a
\]

respectively

\[
\text{case } x \text{ of } [] \Rightarrow F \mid \text{Cons } a r \Rightarrow G a r.
\]

From the internal definitions (not shown here) several properties were automatically derived. We show only the case for lists:

\[
\begin{align*}
(\text{case } [] \text{ of } \Rightarrow F) & | (a # r) \Rightarrow G a r = F \\
(\text{case } b # t \text{ of } \Rightarrow F) & | (a # r) \Rightarrow G a r = G b t \\
[a = [] \Rightarrow P \mid \exists x \ t. \ a = x # t \Rightarrow P] & \Longrightarrow P && \text{— distinctness} \\
[P]; \forall at. \ P t \Rightarrow P(a # t)] & \Longrightarrow Px && \text{— exhaust} \\
\end{align*}
\]

Finally, there is a compiler for primitive and wellfounded recursive function definitions. For example, we may define the sort operation of our running test example by:

\[
\begin{align*}
\text{fun} & \quad \text{ins} : [\alpha :: \text{linorder}, \alpha \text{ list}] \Rightarrow \alpha \text{ list} \\
\text{where} & \quad \text{ins } x [] = [x] \\
& \quad \text{ins } x (y # ys) = \text{if } x < y \text{ then } x # y # ys \text{ else } y # (\text{ins } x y s)
\end{align*}
\]

\(^1\)To increase readability, we use a slightly simplified presentation.
fun sort :: ((α :: linorder) list ⇒ α list)
where sort [] = []
sort (x # xs) = ins x (sort xs)

The internal (non-recursive) constant definition for these operations is quite involved; however, the logical compiler will finally derive all the equations in the statements above from this definition and make them available for automated simplification.

Thus, Isabelle/HOL also provides a large collection of theories like sets, lists, multisets, orderings, and various arithmetic theories which only contain rules derived from conservative definitions. In particular, Isabelle manages a set of executable types and operators, i.e., types and operators for which a compilation to SML, OCaml or Haskell is possible. Setups for arithmetic types such as int have been done; moreover any datatype and any recursive function were included in this executable set (providing that they only consist of executable operators). Similarly, Isabelle manages a large set of (higher-order) rewrite rules into which recursive function definitions were included. Provided that this rule set represents a terminating and confluent rewrite system, the Isabelle simplifier provides also a highly potent decision procedure for many fragments of theories underlying the constraints to be processed when constructing test theorems.

2.3. Featherweight OCL: Design Goals

Featherweight OCL is a formalization of the core of OCL aiming at formally investigating the relationship between the various concepts. At present, it does not attempt to define the complete OCL library. Instead, it concentrates on the core concepts of OCL as well as the types Boolean, Integer, and typed sets (Set(T)). Following the tradition of HOL-OCL [6, 8], Featherweight OCL is based on the following principles:

1. It is an embedding into a powerful semantic meta-language and environment, namely Isabelle/HOL [27].

2. It is a shallow embedding in HOL; types in OCL were injectively mapped to types in Featherweight OCL. Ill-typed OCL specifications cannot be represented in Featherweight OCL and a type in Featherweight OCL contains exactly the values that are possible in OCL. Thus, sets may contain null (Set{null} is a defined set) but not invalid (Set{invalid} is just invalid).

3. Any Featherweight OCL type contains at least invalid and null (the type Void contains only these instances). The logic is consequently four-valued, and there is a null-element in the type Set(A).

4. It is a strongly typed language in the Hindley-Milner tradition. We assume that a pre-process eliminates all implicit conversions due to subtyping by introducing explicit casts (e.g., oclAsType()). The details of such a pre-processing are described in [4]. Casts are semantic functions, typically injections, that may convert data between the different Featherweight OCL types.
5. All objects are represented in an object universe in the HOL-OCL tradition. The universe construction also gives semantics to type casts, dynamic type tests, as well as functions such as `oclAllInstances()`, or `oclIsNew()`.

6. Featherweight OCL types may be arbitrarily nested. For example, the expression `Set{Set{1,2}} = Set{Set{2,1}}` is legal and true.

7. For demonstration purposes, the set type in Featherweight OCL may be infinite, allowing infinite quantification and a constant that contains the set of all Integers. Arithmetic laws like commutativity may therefore be expressed in OCL itself. The iterator is only defined on finite sets.

8. It supports equational reasoning and congruence reasoning, but this requires a differentiation of the different equalities like strict equality, strong equality, meta-equality (HOL). Strict equality and strong equality require a subcalculus, “cp” (a detailed discussion of the different equalities as well as the subcalculus “cp”—for three-valued OCL 2.0—is given in [10]), which is nasty but can be hidden from the user inside tools.

2.4. The Theory Organization

The semantic theory is organized in a quite conventional manner in three layers. The first layer, called the denotational semantics comprises a set of definitions of the operators of the language. Presented as definitional axioms inside Isabelle/HOL, this part assures the logically consistency of the overall construction. The second layer, called logical layer, is derived from the former and centered around the notion of validity of an OCL formula \( P \) for a state-transition from pre-state \( \sigma \) to post-state \( \sigma' \), validity statements were written \((\sigma,\sigma') \models P\). The third layer, called algebraic layer, also derived from the former layers, tries to establish algebraic laws of the form \( P = P' \); such laws are amenable to equational reasoning and also help for automated reasoning and code-generation.

For space reasons, we will restrict ourselves in this paper to a few operators and make a traversal through all three layers to give a high-level description of our formalization. Especially, the details of the semantic construction for sets and the handling of objects and object universes were excluded from a presentation here.

2.4.1. Denotational Semantics

OCL is composed of

1. operators on built-in data structures such as `Boolean`, `Integer`, or `Set(A)`,

2. operators of the user-defined data-model such as accessors, type-casts and tests, and

3. user-defined, side-effect-free methods.
Conceptually, an OCL expression in general and Boolean expressions in particular (i.e., formulae) depends on the pair \((\sigma, \sigma')\) of pre- and post-state. The precise form of states is irrelevant for this paper (compare [13]) and will be left abstract in this presentation. We construct in Isabelle a type-class null that contains two distinguishable elements bot and null. Any type of the form \((\alpha, \bot)\) is an instance of this type-class with bot \(\equiv \bot\) and null \(\equiv [\bot]\). Now, any OCL type can be represented by an HOL type of the form:

\[ V(\alpha) := \text{state} \times \text{state} \rightarrow \alpha :: \text{null} \]

On this basis, we define \(V((\text{bool}, \bot))\) as the HOL type for the OCL type Boolean and define:

\[
\begin{align*}
I[\text{invalid} :: V(\alpha)]\tau &= \text{bot} \\
I[\text{null} :: V(\alpha)]\tau &= \text{null} \\
I[\text{true} :: \text{Boolean}]\tau &= [\text{true}] \\
I[\text{false}]\tau &= [\text{false}] \\
I[X.\text{oclIsUndefined()}]\tau &= (\text{if } I[X]\tau \in \{\text{bot}, \text{null}\} \text{ then } I[\text{true}]\tau \text{ else } I[\text{false}]\tau) \\
I[X.\text{oclIsInvalid()}]\tau &= (\text{if } I[X]\tau = \text{bot} \text{ then } I[\text{true}]\tau \text{ else } I[\text{false}]\tau)
\end{align*}
\]

where \(I[E]\) is the semantic interpretation function commonly used in mathematical textbooks and \(\tau\) stands for pairs of pre- and post state \((\sigma, \sigma')\). For reasons of conciseness, we will write \(\delta X\) for not \(X.\text{oclIsUndefined()}\) and \(\nu X\) for not \(X.\text{oclIsInvalid()}\) throughout this paper.

Due to the used style of semantic representation (a shallow embedding) \(I\) is in fact superfluous and defined semantically as the identity; instead of:

\[ I[\text{true} :: \text{Boolean}]\tau = [\text{true}] \]

we can therefore write:

\[ \text{true} :: \text{Boolean} = \lambda \tau. [\text{true}] \]

In Isabelle theories, this particular presentation of definitions paves the way for an automatic check that the underlying equation has the form of an axiomatic definition and is therefore logically safe. Since all operators of the assertion language depend on the context \(\tau = (\sigma, \sigma')\) and result in values that can be \(\bot\), all expressions can be viewed as evaluations from \((\sigma, \sigma')\) to a type \(\alpha\) which must possess a \(\bot\) and a null-element. Given that such constraints can be expressed in Isabelle/HOL via type classes (written: \(\alpha :: \kappa\)), all types for OCL-expressions are of a form captured by

\[ V(\alpha) := \text{state} \times \text{state} \rightarrow \alpha :: \{\text{bot, null}\}, \]

where state stands for the system state and state \(\times\) state describes the pair of pre-state and post-state and \(\_ := \_\) denotes the type abbreviation.

The current OCL semantics [29, Annex A] uses different interpretation functions for invariants and pre-conditions; we achieve their semantic effect by a syntactic transformation \(\_\pre\) which replaces, for example, all accessor functions \(\_a\) by their counterparts \(\_a@\pre\). For example, \((\text{self}.a > 5)\pre\) is just \((\text{self}.a@\pre > 5)\). This way, also invariants and pre-conditions can be interpreted by the same interpretation function and have the same type of an evaluation \(V(\alpha)\).
On this basis, one can define the core logical operators \texttt{not} and \texttt{and} as follows:

\[
I[\text{not } X] \tau = (\text{case } I[X] \tau \text{ of } \\
\bot \Rightarrow \bot \\
\lfloor \bot \rfloor \Rightarrow \lfloor \bot \rfloor \\
\lfloor \neg x \rfloor \Rightarrow \lfloor \neg x \rfloor)
\]

\[
I[X \text{ and } Y] \tau = (\text{case } I[X] \tau \text{ of } \\
\bot \Rightarrow (\text{case } I[Y] \tau \text{ of } \\
\bot \Rightarrow \bot \\
\lfloor \bot \rfloor \Rightarrow \lfloor \bot \rfloor \\
\lfloor \text{false} \rfloor \Rightarrow \lfloor \text{false} \rfloor) \\
\lfloor \bot \rfloor \Rightarrow (\text{case } I[Y] \tau \text{ of } \\
\bot \Rightarrow \bot \\
\lfloor \bot \rfloor \Rightarrow \lfloor \bot \rfloor \\
\lfloor \text{false} \rfloor \Rightarrow \lfloor \text{false} \rfloor \\
\lfloor \text{true} \rfloor \Rightarrow \lfloor \text{false} \rfloor) \\
\lfloor \text{true} \rfloor \Rightarrow (\text{case } I[Y] \tau \text{ of } \\
\bot \Rightarrow \bot \\
\lfloor \bot \rfloor \Rightarrow \lfloor \bot \rfloor \\
\lfloor \text{false} \rfloor \Rightarrow \lfloor \text{false} \rfloor \\
\lfloor y \rfloor \Rightarrow \lfloor y \rfloor) \\
\lfloor \text{false} \rfloor \Rightarrow \lfloor \text{false} \rfloor)
\]

These non-strict operations were used to define the other logical connectives in the usual classical way: \( X \lor Y \equiv (\text{not } X) \land (\text{not } Y) \) or \( X \impliedby Y \equiv (\text{not } X) \lor Y \).

The default semantics for an OCL library operator is strict semantics; this means that the result of an operation \( f \) is invalid if one of its arguments is invalid. For a semantics comprising null, we suggest to stay conform to the standard and define the addition for integers as follows:

\[
I[x + y] \tau = \begin{cases} 
\lfloor \lfloor \text{true} \rfloor \land I[\delta x] \tau = \lfloor \text{true} \rfloor \land I[\delta y] \tau = \lfloor \text{true} \rfloor \\
\text{then } \lfloor I[x] \tau \rfloor + \lfloor I[y] \tau \rfloor \rfloor \\
\text{else } \bot
\end{cases}
\]

where the operator “\( \times \)” on the left-hand side of the equation denotes the OCL addition of type \([\mathbb{V}(\mathbb{int}_1)_1, \mathbb{V}(\mathbb{int}_1)_1] \Rightarrow \mathbb{V}(\mathbb{int}_1)_1\) while the “\( + \)” on the right-hand side of the equation of type \([\mathbb{int}, \mathbb{int}] \Rightarrow \mathbb{int}\) denotes the integer-addition from the HOL library.

2.4.2. Logical Layer

The topmost goal of the logic for OCL is to define the \textit{validity statement}:

\[(\sigma, \sigma') \models P,\]
where $\sigma$ is the pre-state and $\sigma'$ the post-state of the underlying system and $P$ is a formula. Informally, a formula $P$ is valid if and only if its evaluation in $(\sigma, \sigma')$ (i.e., $\tau$ for short) yields true. Formally this means:

$$\tau \models P \equiv (I[P]\tau = [[true]]) .$$

On this basis, classical, two-valued inference rules can be established for reasoning over the logical connective, the different notions of equality, definedness and validity. Generally speaking, rules over logical validity can relate bits and pieces in various OCL terms and allow—via strong logical equality discussed below—the replacement of semantically equivalent sub-expressions. The core inference rules are:

$$\tau \models true \quad \neg(\tau \models false) \quad \neg(\tau \models invalid) \quad \neg(\tau \models null)$$

$$\tau \models not P \implies \neg(\tau \models P)$$

$$\tau \models P \land Q \implies \tau \models P \quad \tau \models P \land Q \implies \tau \models Q$$

$$\tau \models P \implies (if \ P \ then \ B_1 \ else \ B_2 \ endif)\tau = B_1 \ \tau$$

$$\tau \models not P \implies (if \ P \ then \ B_1 \ else \ B_2 \ endif)\tau = B_2 \ \tau$$

$$\tau \models P \implies \tau \models \delta P \quad \tau \models \delta X \implies \tau \models \nu X$$

By the latter two properties it can be inferred that any valid property $P$ (so for example: a valid invariant) is defined, which allows to infer for terms composed by strict operations that their arguments and finally the variables occurring in it are valid or defined.

We propose to distinguish the strong logical equality (written $\_ \triangleq \_\$), which follows the general principle that “equals can be replaced by equals,” from the strict referential equality (written $\_ \equiv \_$), which is an object-oriented concept that attempts to approximate and to implement the former. Strict referential equality, which is the default in the OCL language and is written $\_ \ = \ _\$ in the standard, is an overloaded concept and has to be defined for each OCL type individually; for objects resulting from class definitions, it is implemented by comparing the references to the objects. In contrast, strong logical equality is a polymorphic concept which is defined once and for all by:

$$I[X \triangleq Y]\tau \equiv [[I[X]\tau = I[Y]\tau]]$$

It enjoys nearly the laws of a congruence:

$$\tau \models (x \triangleq x)$$

$$\tau \models (x \triangleq y) \implies \tau \models (y \triangleq x)$$

$$\tau \models (x \triangleq y) \implies \tau \models (y \triangleq z) \implies \tau \models (x \triangleq z)$$

$$cp \ P \implies \tau \models (x \triangleq y) \implies \tau \models (\ P \ x \ ) \implies \tau \models (\ P \ y \ )$$

where the predicate $cp$ stands for context-passing, a property that is characterized by $P(X)$ equals $\lambda \tau. \ P(\lambda_\_ X\tau)\tau$. It means that the state tuple $(\sigma, \sigma')$ is passed unchanged from surrounding expressions to sub-expressions. it is true for all pure OCL expressions (but not arbitrary mixtures of OCL and HOL) in Featherweight OCL. The necessary side-calculus for establishing $cp$ can be fully automated.
The logical layer of the Featherweight OCL rules gives also a means to convert an OCL formula living in its four-valued world into a representation that is classically two-valued and can be processed by standard SMT solvers such as CVC3 24 or Z3 20. δ-closure rules for all logical connectives have the following format, e. g.:

\[
\begin{align*}
\tau \models \delta x & \implies (\tau \models \text{not } x) = (\neg(\tau \models x)) \\
\tau \models \delta x \implies \tau \models \delta y & \implies (\tau \models x \text{ and } y) = (\tau \models x \land \tau \models y) \\
\tau \models \delta x \implies \tau \models \delta y & \implies (\tau \models (x \implies y)) = ((\tau \models x) \rightarrow (\tau \models y))
\end{align*}
\]

Together with the general case-distinction

\[
\tau \models \delta x \lor \tau \models x \triangleq \text{invalid} \lor \tau \models x \triangleq \text{null}
\]

which is possible for any OCL type, a case distinction on the variables in a formula can be performed; due to strictness rules, formulae containing somewhere a variable \( x \) that is known to be invalid or null reduce usually quickly to contradictions. For example, we can infer from an invariant \( \tau \models x \equiv y - 3 \) that we have \( \tau \models x = y - 3 \land \tau \models \delta x \land \tau \models \delta y \). We call the latter formula the δ-closure of the former. Now, we can convert a formula like \( \tau \models x > 0 \text{ or } 3 \times y > x \times x \) into the equivalent formula \( \tau \models x > 0 \lor \tau \models 3 \times y > x \times x \) and thus internalize the OCL-logic into a classical (and more tool-conform) logic. This works—for the price of a potential, but due to the usually “rich” δ-closures of invariants rare—exponential blow-up of the formula for all OCL formulas.

2.4.3. Algebraic Layer

Based on the logical layer, we build a system with simpler rules which are amenable to automated reasoning. We restrict ourselves to pure equations on OCL expressions, where the used equality is the meta-(HOL-)equality.

Our denotational definitions on not and and can be re-formulated in the following ground equations:

\[
\begin{align*}
v \text{ invalid} & = \text{false} & v \text{ null} & = \text{true} \\
v \text{ true} & = \text{true} & v \text{ false} & = \text{true} \\
\delta \text{ invalid} & = \text{false} & \delta \text{ null} & = \text{false} \\
\delta \text{ true} & = \text{true} & \delta \text{ false} & = \text{true} \\
\text{not invalid} & = \text{invalid} & \text{not null} & = \text{null} \\
\text{not true} & = \text{false} & \text{not false} & = \text{true} \\
(\text{null and true}) & = \text{null} & (\text{null and false}) & = \text{false} \\
(\text{null and null}) & = \text{null} & (\text{null and invalid}) & = \text{invalid} \\
(\text{false and true}) & = \text{false} & (\text{false and false}) & = \text{false} \\
(\text{false and null}) & = \text{false} & (\text{false and invalid}) & = \text{false}
\end{align*}
\]
\( (\text{true and true}) = \text{true} \) \quad \( (\text{true and false}) = \text{false} \) \\
\( (\text{true and null}) = \text{null} \) \quad \( (\text{true and invalid}) = \text{invalid} \) \\
\( (\text{invalid and true}) = \text{invalid} \) \\
\( (\text{invalid and false}) = \text{false} \) \\
\( (\text{invalid and null}) = \text{invalid} \) \\
\( (\text{invalid and invalid}) = \text{invalid} \)

On this core, the structure of a conventional lattice arises:

\[
\begin{align*}
X \land X &= X \\
X \land Y &= Y \land X \\
\text{false} \land X &= \text{false} \\
X \land \text{false} &= \text{false} \\
\text{true} \land X &= X \\
X \land \text{true} &= X \\
X \land (Y \land Z) &= X \land Y \land Z
\end{align*}
\]

as well as the dual equalities for \( \_ \lor \_ \) and the De Morgan rules. This wealth of algebraic properties makes the understanding of the logic easier as well as automated analysis possible: it allows for, for example, computing a DNF of invariant systems (by clever term-rewriting techniques) which are a prerequisite for \( \delta \)-closures.

The above equations explain the behavior for the most-important non-strict operations. The clarification of the exceptional behaviors is of key-importance for a semantic definition the standard and the major deviation point from HOL-OCL \([6, 8]\), to Featherweight OCL as presented here. The standard expresses at many places that most operations are strict, i.e., enjoy the properties (exemplary for \( \_ + \_ \)):

\[
\begin{align*}
\text{invalid} + X &= \text{invalid} \\
X + \text{invalid} &= \text{invalid} \\
X + \text{null} &= \text{invalid} \\
\text{null} + X &= \text{invalid} \\
\text{null.oclAsType}(X) &= \text{invalid}
\end{align*}
\]

besides “classical” exceptional behavior:

\[
\begin{align*}
1 / 0 &= \text{invalid} \\
1 / \text{null} &= \text{invalid} \\
\text{null->isEmpty}() &= \text{true}
\end{align*}
\]

Moreover, there is also the proposal to use \text{null} as a kind of “don’t know” value for all strict operations, not only in the semantics of the logical connectives. Expressed in algebraic equations, this semantic alternative (this is \textit{not} Featherweight OCL at present) would boil down to:

\[
\begin{align*}
\text{invalid} + X &= \text{invalid} \\
X + \text{invalid} &= \text{invalid} \\
X + \text{null} &= \text{null} \\
\text{null} + X &= \text{null} \\
\text{null.oclAsType}(X) &= \text{null} \\
1 / 0 &= \text{invalid} \\
1 / \text{null} &= \text{null} \\
\text{null->isEmpty}() &= \text{null}
\end{align*}
\]

While this is logically perfectly possible, while it can be argued that this semantics is “intuitive”, and although we do not expect a too heavy cost in deduction when computing
δ-closures, we object that there are other, also “intuitive” interpretations that are even more wide-spread: In classical spreadsheet programs, for example, the semantics tends to interpret null (representing empty cells in a sheet) as the neutral element of the type, so 0 or the empty string, for example. This semantic alternative (this is not Featherweight OCL at present) would yield:

\[
\begin{align*}
\text{invalid} + X &= \text{invalid} & \text{X} + \text{invalid} &= \text{invalid} \\
X + \text{null} &= X & \text{null} + X &= X \\
\text{null}.\text{oclAsType}(X) &= \text{invalid} \\
1 / 0 &= \text{invalid} & 1 / \text{null} &= \text{invalid} \\
\text{null} \rightarrow \text{isEmpty}() &= \text{true}
\end{align*}
\]

Algebraic rules are also the key for execution and compilation of Featherweight OCL expressions. We derived, e. g.:

\[
\begin{align*}
\delta \text{Set}() &= \text{true} \\
\delta (X \rightarrow \text{including}(x)) &= \delta X \text{ and } \delta x \\
\text{Set}() \rightarrow \text{includes}(x) &= (\text{if } \nu x \text{ then false} \\
& \quad \text{else invalid endif}) \\
(X \rightarrow \text{including}(x) \rightarrow \text{includes}(y)) &= \\
& (\text{if } \delta X \\
& \quad \text{then if } x = y \\
& \quad \text{then true} \\
& \quad \text{else } X \rightarrow \text{includes}(y) \\
& \quad \text{endif} \\
& \quad \text{else invalid} \\
& \quad \text{endif})
\end{align*}
\]

As \text{Set\{1,2\}} is only syntactic sugar for

\[
\text{Set}() \rightarrow \text{including}(1) \rightarrow \text{including}(2)
\]

an expression like \text{Set\{1,2\}}\rightarrow\text{includes}(null) becomes decidable in Featherweight OCL by a combination of rewriting and code-generation and execution. The generated documentation from the theory files can thus be enriched by numerous “test-statements” like:

value "τ |= (Set\{Set\{2,null\}\} ∋ Set\{Set\{null,2\}\})"

which have been machine-checked and which present a high-level and in our opinion fairly readable information for OCL tool manufacturers and users.

\[2\]In spreadsheet programs the interpretation of null varies from operation to operation; e. g., the average function treats null as non-existing value and not as 0.
2.5. Object-oriented Datatype Theories

As mentioned earlier, the OCL is composed of

1. operators on built-in data structures such as Boolean, Integer or Set(\_), and

2. operators of the user-defined data model such as accessors, type casts and tests.

In the following, we will refine the concepts of a user-defined data-model (implied by a class-model, visualized by a class-diagram) as well as the notion of state used in the previous section to much more detail. In contrast to wide-spread opinions, UML class diagrams represent in a compact and visual manner quite complex, object-oriented data-types with a surprisingly rich theory. It is part of our endeavor here to make this theory explicit and to point out corner cases. A UML class diagram—underlying a given OCL formula—produces several implicit operations which become accessible via appropriate OCL syntax:

1. Classes and class names (written as \(C_1, \ldots, C_n\)), which become types of data in OCL. Class names declare two projector functions to the set of all objects in a state: \(C_i.allInstances()\) and \(C_i.allInstances@pre()\),

2. an inheritance relation \(_<_\) on classes and a collection of attributes \(A\) associated to classes,

3. two families of accessors; for each attribute \(a\) in a class definition (denoted \(X.a :: C_i \rightarrow A\) and \(X.a@pre :: C_i \rightarrow A\) for \(A \in \{V(\_\_), C_1, \ldots, C_n\}\)),

4. type casts that can change the static type of an object of a class \((X.oclAsType(C_i)\) of type \(C_j \rightarrow C_i)\)

5. two dynamic type tests \((X.oclIsTypeOf(C_i)\) and \(X.oclIsKindOf(C_i)\)),

6. and last but not least, for each class name \(C_i\) there is an instance of the overloaded referential equality (written \(_=_\)).

Assuming a strong static type discipline in the sense of Hindley-Milner types, Featherweight OCL has no “syntactic subtyping.” This does not mean that subtyping cannot be expressed semantically in Featherweight OCL; by giving a formal semantics to type-casts, subtyping becomes an issue of the front-end that can make implicit type-coerions explicit by introducing explicit type-casts. Our perspective shifts the emphasis on the semantic properties of casting, and the necessary universe of object representations (induced by a class model) that allows to establish them.

2.5.1. Object Universes

It is natural to construct system states by a set of partial functions \(f\) that map object identifiers oid to some representations of objects:

\[
\text{typedef } \alpha \text{ state } := \{\sigma : \text{oid} \rightarrow \alpha \mid \text{inv}_\sigma(\sigma)\} \tag{2.12}
\]
where \( \text{inv}_\sigma \) is a to be discussed invariant on states.

The key point is that we need a common type \( \alpha \) for the set of all possible object representations. Object representations model “a piece of typed memory,” i.e., a kind of record comprising administration information and the information for all attributes of an object: here, the primitive types as well as collections over them are stored directly in the object representations, class types and collections over them are represented by oid’s (respectively lifted collections over them).

In a shallow embedding which must represent UML types injectively by HOL types, there are two fundamentally different ways to construct such a set of object representations, which we call an object universe \( \mathfrak{A} \):

1. an object universe can be constructed for a given class model, leading to closed world semantics, and
2. an object universe can be constructed for a given class model and all its extensions by new classes added into the leaves of the class hierarchy, leading to an open world semantics.

For the sake of simplicity, we chose the first option for Featherweight OCL, while HOL-OCL [7] used an involved construction allowing the latter.

A naïve attempt to construct \( \mathfrak{A} \) would look like this: the class type \( C_i \) induced by a class will be the type of such an object representation:

\[
C_i := \text{oid} \times A_{i1} \times \cdots \times A_{ik} \times (C_{j1} + \cdots + C_{jm}) \bot
\]

where \( A_{ik} \) ranges over the local attribute types of \( C_i \) and \( C_{jl} \) ranges over all class type extensions of the subclass \( C_j \) of \( C_i \).

It is possible to define constructors, accessors, and the referential equality on this object universe. However, the treatment of type casts and type tests cannot be faithful with common object-oriented semantics, be it in UML or Java: casting up along the class hierarchy can only be implemented by loosing information, such that casting up and casting down will not give the required identity:

\[
X.\text{oclIsTypeOf}(C_k) \implies X.\text{oclAsType}(C_i).\text{oclAsType}(C_k) = X
\]

whenever \( C_k < C_i \) and \( X \) is valid. (2.14)

(2.15)

To overcome this limitation, we introduce an auxiliary type \( C_{\text{ext}} \) for class type extension; together, they were inductively defined for a given class diagram:

Let \( C_i \) be a class with a possibly empty set of subclasses \( \{C_{j1}, \ldots, C_{jm}\} \).

- Then the class type extension \( C_{\text{ext}} \) associated to \( C_i \) is \( A_{i1} \times \cdots \times A_{in} \times (C_{j1\text{ext}} + \cdots + C_{jm\text{ext}}) \bot \) where \( A_{ik} \) ranges over the local attribute types of \( C_i \) and \( C_{jl\text{ext}} \) ranges over all class type extensions of the subclass \( C_j \) of \( C_i \).
Then the class type for $C_i$ is $oid \times A_{i_1} \times \cdots \times A_{i_n} \times (C_{j_1}^{\text{ext}} + \cdots + C_{j_m}^{\text{ext}})_{\perp}$ where $A_{i_k}$ ranges over the inherited and local attribute types of $C_i$ and $C_{j_l}^{\text{ext}}$ ranges over all class type extensions of the subclass $C_j$ of $C_i$.

Example instances of this scheme—outlining a compiler—can be found in Section 6.1 and Section 7.1. This construction can not be done in HOL itself since it involves quantifications and iterations over the “set of class-types”; rather, it is a meta-level construction. Technically, this means that we need a compiler to be done in SML on the syntactic “meta-model”-level of a class model.

With respect to our semantic construction here, which above all means is intended to be type-safe, this has the following consequences:

- there is a generic theory of states, which must be formulated independently from a concrete object universe,
- there is a principle of translation (captured by the inductive scheme for class type extensions and class types above) that converts a given class model into an concrete object universe,
- there are fixed principles that allow to derive the semantic theory of any concrete object universe, called the object-oriented datatype theory.

We will work out concrete examples for the construction of the object-universes in Section 6.1 and Section 7.1 and the derivation of the respective datatype theories. While an automatization is clearly possible and desirable for concrete applications of Featherweight OCL, we consider this out of the scope of this paper which has a focus on the semantic construction and its presentation.

### 2.5.2. Accessors on Objects and Associations

Our choice to use a shallow embedding of OCL in HOL and, thus, having an injective mapping from OCL types to HOL types, results in type-safety of Featherweight OCL. Arguments and results of accessors are based on type-safe object representations and not oid’s. This implies the following scheme for an accessor:

- The evaluation and extraction phase. If the argument evaluation results in an object representation, the oid is extracted, if not, exceptional cases like invalid are reported.
- The dereferentiation phase. The oid is interpreted in the pre- or post-state, the resulting object is casted to the expected format. The exceptional case of nonexistence in this state must be treated.
- The selection phase. The corresponding attribute is extracted from the object representation.
• The re-construction phase. The resulting value has to be embedded in the adequate HOL type. If an attribute has the type of an object (not value), it is represented by an optional (set of) oid, which must be converted via dereferentiation in one of the states to produce an object representation again. The exceptional case of nonexistence in this state must be treated.

The first phase directly translates into the following formalization:

definition

\[
\text{eval}_{\text{extract}} X f = (\lambda \tau. \text{case } X \tau \text{ of } \perp \Rightarrow \text{invalid } \tau \text{ exception } \mid \perp_j \Rightarrow \text{invalid } \tau \text{ deref. null } \mid \uparrow \text{oid}_j \Rightarrow f \text{ (oid of obj) } \tau)
\]

(2.16)

For each class \( C \), we introduce the dereferentiation phase of this form:

definition deref_{oid} C \text{ fst snd f oid} = (\lambda \tau. \text{case (heap (fst snd ) oid of } \mid \text{in}_C \text{ obj}_j \Rightarrow f \text{ obj } \tau \mid \perp_j \Rightarrow \text{invalid } \tau)

(2.17)

The operation yields undefined if the oid is uninterpretable in the state or referencing an object representation not conforming to the expected type.

We turn to the selection phase: for each class \( C \) in the class model with at least one attribute, and each attribute \( a \) in this class, we introduce the selection phase of this form:

definition select \( a \) f = (\lambda \text{mk}_C \text{ oid } \cdots \perp \cdots C_X \text{ ext} \Rightarrow \text{null } \mid \text{mk}_C \text{ oid } \cdots a_j \cdots C_X \text{ ext} \Rightarrow f (\lambda x \perp \text{oid}_j \text{ oid } x) a)

(2.18)

This works for definitions of basic values as well as for object references in which the \( a \) is of type oid. To increase readability, we introduce the functions:

definition \text{in}_{\text{pre state}} = \text{fst} \quad \text{first component }

definition \text{in}_{\text{post state}} = \text{snd} \quad \text{second component }

definition \text{reconst}_{\text{basetype}} = \text{id} \quad \text{identity function }

(2.19)

Let \( \_ \text{.getBase} \) be an accessor of class \( C \) yielding a value of base-type \( A_{\text{base}} \). Then its definition is of the form:

definition \_ \text{.getBase} :: C \Rightarrow A_{\text{base}}

where \( X \text{.getBase} = \text{eval}_{\text{extract}} X (\text{deref}_{\text{oid}} C \text{ in}_{\text{post state}} (\text{select}_{\text{getBase}} \text{ reconst}_{\text{basetype}})) \)

(2.20)

Let \( \_ \text{.getObject} \) be an accessor of class \( C \) yielding a value of object-type \( A_{\text{object}} \). Then its definition is of the form:

definition \_ \text{.getObject} :: C \Rightarrow A_{\text{object}}

where \( X \text{.getObject} = \text{eval}_{\text{extract}} X (\text{deref}_{\text{oid}} C \text{ in}_{\text{post state}} (\text{select}_{\text{getObject}} (\text{deref}_{\text{oid}} C \text{ in}_{\text{post state}}))) \)

(2.21)
The variant for an accessor yielding a collection is omitted here; its construction follows by the application of the principles of the former two. The respective variants were produced when in_post_state is replaced by in_pre_state.

Examples for the construction of accessors via associations can be found in Section 6.1.8, the construction of accessors via attributes in Section 7.1.8. The construction of casts and type tests ->oclIsTypeOf() and ->oclIsKindOf() is similarly.

In the following, we discuss the role of multiplicities on the types of the accessors. Depending on the specified multiplicity, the evaluation of an attribute can yield just a value (multiplicity 0..1 or 1) or a collection type like Set or Sequence of values (otherwise). A multiplicity defines a lower bound as well as a possibly infinite upper bound on the cardinality of the attribute’s values.

**Single-Valued Attributes**

If the upper bound specified by the attribute’s multiplicity is one, then an evaluation of the attribute yields a single value. Thus, the evaluation result is not a collection. If the lower bound specified by the multiplicity is zero, the evaluation is not required to yield a non-null value. In this case an evaluation of the attribute can return null to indicate an absence of value.

To facilitate accessing attributes with multiplicity 0..1, the OCL standard states that single values can be used as sets by calling collection operations on them. This implicit conversion of a value to a Set is not defined by the standard. We argue that the resulting set cannot be constructed the same way as when evaluating a Set literal. Otherwise, null would be mapped to the singleton set containing null, but the standard demands that the resulting set is empty in this case. The conversion should instead be defined as follows:

```oclass
context OclAny::asSet():T
  post: if self = null then result = Set{}
    else result = Set{self} endif
```

**Collection-Valued Attributes**

If the upper bound specified by the attribute’s multiplicity is larger than one, then an evaluation of the attribute yields a collection of values. This raises the question whether null can belong to this collection. The OCL standard states that null can be owned by collections. However, if an attribute can evaluate to a collection containing null, it is not clear how multiplicity constraints should be interpreted for this attribute. The question arises whether the null element should be counted or not when determining the cardinality of the collection. Recall that null denotes the absence of value in the case of a cardinality upper bound of one, so we would assume that null is not counted. On the other hand, the operation size defined for collections in OCL does count null.

We propose to resolve this dilemma by regarding multiplicities as optional. This point of view complies with the UML standard, that does not require lower and upper bounds
to be defined for multiplicities. In case a multiplicity is specified for an attribute, i.e., a lower and an upper bound are provided, we require any collection the attribute evaluates to not contain null. This allows for a straightforward interpretation of the multiplicity constraint. If bounds are not provided for an attribute, we consider the attribute values to not be restricted in any way. Because in particular the cardinality of the attribute’s values is not bounded, the result of an evaluation of the attribute is of collection type. As the range of values that the attribute can assume is not restricted, the attribute can evaluate to a collection containing null. The attribute can also evaluate to invalid. Allowing multiplicities to be optional in this way gives the modeler the freedom to define attributes that can assume the full ranges of values provided by their types. However, we do not permit the omission of multiplicities for association ends, since the values of association ends are not only restricted by multiplicities, but also by other constraints enforcing the semantics of associations. Hence, the values of association ends cannot be completely unrestricted.

**The Precise Meaning of Multiplicity Constraints**

We are now ready to define the meaning of multiplicity constraints by giving equivalent invariants written in OCL. Let \( a \) be an attribute of a class \( C \) with a multiplicity specifying a lower bound \( m \) and an upper bound \( n \). Then we can define the multiplicity constraint on the values of attribute \( a \) to be equivalent to the following invariants written in OCL:

\[
\begin{align*}
\text{context } C \quad &\text{inv lowerBound: } a->\text{size()} \geq m \\
&\text{inv upperBound: } a->\text{size()} \leq n \\
&\text{inv notNull: } \text{not } a->\text{includes}(\text{null})
\end{align*}
\]

If the upper bound \( n \) is infinite, the second invariant is omitted. For the definition of these invariants we are making use of the conversion of single values to sets described in Section 2.5.2. If \( n \leq 1 \), the attribute \( a \) evaluates to a single value, which is then converted to a Set on which the size operation is called.

If a value of the attribute \( a \) includes a reference to a non-existent object, the attribute call evaluates to invalid. As a result, the entire expressions evaluate to invalid, and the invariants are not satisfied. Thus, references to non-existent objects are ruled out by these invariants. We believe that this result is appropriate, since we argue that the presence of such references in a system state is usually not intended and likely to be the result of an error. If the modeler wishes to allow references to non-existent objects, she can make use of the possibility described above to omit the multiplicity.

### 2.5.3. Other Operations on States

Defining \( _.\text{allInstances()} \) is straightforward; the only difference is the property \( T\text{.allInstances()}->\text{excludes}(\text{null}) \) which is a consequence of the fact that null’s are values and do not “live” in the state. In our semantics which admits states with

---

3We are however aware that a well-formedness rule of the UML standard does define a default bound of one in case a lower or upper bound is not specified.
“dangling references,” it is possible to define a counterpart to \_.oclIsNew() called \_.oclIsDeleted() which asks if an object id (represented by an object representation) is contained in the pre-state, but not the post-state.

OCL does not guarantee that an operation only modifies the path-expressions mentioned in the postcondition, i.e., it allows arbitrary relations from pre-states to post-states. This framing problem is well-known (one of the suggested solutions is \[23\]). We define

\[
(S : \text{Set}(\text{OclAny})) \rightarrow \text{oclIsModifiedOnly}() : \text{Boolean}
\]

where \(S\) is a set of object representations, encoding a set of oid’s. The semantics of this operator is defined such that for any object whose oid is not represented in \(S\) and that is defined in pre and post state, the corresponding object representation will not change in the state transition. A simplified presentation is as follows:

\[
I[X \rightarrow \text{oclIsModifiedOnly}()](\sigma, \sigma') \equiv \begin{cases} 
\bot & \text{if } X' = \bot \lor \text{null} \in X' \\
\forall i \in M. \sigma_i = \sigma'_i & \text{otherwise}.
\end{cases}
\]

where \(X' = I[X](\sigma, \sigma')\) and \(M = (\text{dom} \sigma \cap \text{dom} \sigma') - \{\text{OidOf } x \mid x \in [X']\}\). Thus, if we require in a postcondition \(\text{Set}{} \rightarrow \text{oclIsModifiedOnly}()\) and exclude via \_.oclIsNew() and \_.oclIsDeleted() the existence of new or deleted objects, the operation is a query in the sense of the OCL standard, i.e., the \text{isQuery} property is true. So, whenever we have \(\tau \models X \rightarrow \text{excluding}(s.a) \rightarrow \text{oclIsModifiedOnly}()\) and \(\tau \models X \rightarrow \forall x (\text{not}(x = s.a))\), we can infer that \(\tau \models s.a \equiv s.a \_\text{pre}\).

### 2.6. A Machine-checked Annex A

Isabelle, as a framework for building formal tools \[37\], provides the means for generating formal documents. With formal documents (such as the one you are currently reading) we refer to documents that are machine-generated and ensure certain formal guarantees. In particular, all formal content (e.g., definitions, formulae, types) are checked for consistency during the document generation.

For writing documents, Isabelle supports the embedding of informal texts using a \LaTeX-based markup language within the theory files. To ensure the consistency, Isabelle supports to use, within these informal texts, antiquotations that refer to the formal parts and that are checked while generating the actual document as PDF. For example, in an informal text, the antiquotation @\{thm "not_not"\} will instruct Isabelle to lock-up the (formally proven) theorem of name ocl\_not\_not and to replace the antiquotation with the actual theorem, i.e., \text{not( not x)} = x.

Figure 2.2 illustrates this approach. Figure 2.2a shows the jEdit-based development environment of Isabelle with an excerpt of one of the core theories of Featherweight OCL. Figure 2.2b shows the generated PDF document where all antiquotations are replaced. Moreover, the document generation tools allows for defining syntactic sugar as well as skipping technical details of the formalization.

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Thus, applying the Featherweight OCL approach to writing an updated Annex A that provides a formal semantics of the most fundamental concepts of OCL would ensure

1. that all formal context is syntactically correct and well-typed, and
2. all formal definitions and the derived logical rules are semantically consistent.

Overall, this would contribute to one of the main goals of the OCL 2.5 RFP, as discussed at the OCL meeting in Aachen [15].
Part II.

A Proposal for Formal Semantics of OCL 2.5
3. Formalization I: Core Definitions

theory
  OCL-core
imports
  Main
begin

3.1. Preliminaries

3.1.1. Notations for the Option Type

First of all, we will use a more compact notation for the library option type which occur all over in our definitions and which will make the presentation more like a textbook:

notation Some (⌊(-)⌋)
notation None (⊥)

The following function (corresponding to the in the Isabelle/HOL library) is defined as the inverse of the injection Some.

fun drop :: 'a option ⇒ 'a (⌈(-)⌉)
where drop-lift[simp]: ⌈v⌉ = v

3.1.2. Minimal Notions of State and State Transitions

Next we will introduce the foundational concept of an object id (oid), which is just some infinite set.

In order to assure executability of as much as possible formulas, we fixed the type of object id’s to just natural numbers.

type-synonym oid = nat

We refrained from the alternative:

type-synonym oid = ind

which is slightly more abstract but non-executable.

States are just a partial map from oid’s to elements of an object universe 'A, and state transitions pairs of states . . .

record ('A)state =
  heap :: oid → 'A
  assoc2 :: oid → (oid × oid) list
  assoc3 :: oid → (oid × oid × oid) list
type-synonym \((\forall)st = \forall state \times \forall state\)

### 3.1.3. Prerequisite: An Abstract Interface for OCL Types

To have the possibility to nest collection types, such that we can give semantics to expressions like `Set\{Set\{2\},null\}`, it is necessary to introduce a uniform interface for types having the `invalid` (= bottom) element. The reason is that we impose a data-invariant on raw-collection types which assures that the `invalid` element is not allowed inside the collection; all raw-collections of this form were identified with the `invalid` element itself. The construction requires that the new collection type is not comparable with the raw-types (consisting of nested option type constructions), such that the data-invariant must be expressed in terms of the interface. In a second step, our base-types will be shown to be instances of this interface.

This uniform interface consists in a type class requiring the existence of a bot and a null element. The construction proceeds by abstracting the null (defined by `[\bot]` on `\forall a option option`) to a `null` element, which may have an arbitrary semantic structure, and an undefinedness element `\bot` to an abstract undefinedness element `bot` (also written `\bot` whenever no confusion arises). As a consequence, it is necessary to redefine the notions of invalid, defined, valuation etc. on top of this interface.

This interface consists in two abstract type classes `bot` and `null` for the class of all types comprising a bot and a distinct null element.

```plaintext
class bot =
  fixes bot :: \forall a
  assumes nonEmpty : \exists x. x \neq bot

class null = bot +
  fixes null :: \forall a
  assumes null-is-valid : null \neq bot
```

### 3.1.4. Accommodation of Basic Types to the Abstract Interface

In the following it is shown that the “option-option” type is in fact in the `null` class and that function spaces over these classes again “live” in these classes. This motivates the default construction of the semantic domain for the basic types (`Boolean`, `Integer`, `Real`, ...).

```plaintext
instantiation option :: (type)bot begin
  definition bot-option-def: (bot::\forall a option) \equiv (None::\forall a option)
  instance ⟨proof⟩
end

instantiation option :: (bot)null
```

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null-option-def: (null::'a::bot option) ≡ [ bot ]

instance (proof)
end

fun :: (type,bot) bot
begin
  definition bot-fun-def: bot ≡ (λ x. bot)

  instance (proof)
end

fun :: (type,null) null
begin
  definition null-fun-def: (null::'a ⇒ 'b::null) ≡ (λ x. null)

  instance (proof)
end

A trivial consequence of this adaption of the interface is that abstract and concrete
versions of null are the same on base types (as could be expected).

3.1.5. The Semantic Space of OCL Types: Valuations

Valuations are now functions from a state pair (built upon data universe 'A) to an
arbitrary null-type (i.e., containing at least a distinguished null and invalid element).

type-synonym ('A,'α) val = 'A st ⇒ 'α::null

The definitions for the constants and operations based on valuations will be geared
towards a format that Isabelle can check to be a “conservative” (i.e., logically safe)
axiomatic definition. By introducing an explicit interpretation function (which happens
to be defined just as the identity since we are using a shallow embedding of OCL into
HOL), all these definitions can be rewritten into the conventional semantic textbook
format as follows:

definition Sem :: 'a ⇒ 'a (I[])
where I[x] ≡ x

As a consequence of semantic domain definition, any OCL type will have the two
semantic constants invalid (for exceptional, aborted computation) and null:

definition invalid :: ('A,'α::bot) val
where invalid ≡ λ τ. bot

This conservative Isabelle definition of the polymorphic constant invalid is equivalent
with the textbook definition:

lemma textbook-invalid: I[invalid]τ = bot
Note that the definition:

\[
\text{definition} \ \text{null} :: (\forall \alpha :: \text{null}) \ \text{val} \\
\text{where} \quad \text{null} \equiv \lambda \tau. \ \text{null}
\]

is not necessary since we defined the entire function space over null types again as null-types; the crucial definition is \( \text{null} \equiv \lambda x. \ \text{null} \). Thus, the polymorphic constant \( \text{null} \) is simply the result of a general type class construction. Nevertheless, we can derive the semantic textbook definition for the OCL null constant based on the abstract null:

\[
\text{lemma} \ \text{textbook-null-fun}: \ I[\text{null}:(\forall \alpha :: \text{null}) \ \text{val}] \ \tau = (\text{null}:(\forall \alpha :: \text{null})
\]

\[
\text{3.2. Definition of the Boolean Type}
\]

The semantic domain of the (basic) boolean type is now defined as the Standard: the space of valuation to \( \text{bool} \ \text{option} \ \text{option} \):

\[
\text{type-synonym} \ (\forall \alpha) \ \text{Boolean} = (\forall \alpha, \text{bool} \ \text{option} \ \text{option}) \ \text{val}
\]

\[
\text{3.2.1. Basic Constants}
\]

\[
\text{lemma} \ \text{bot-Boolean-def} : (\text{bot}:(\forall \alpha) \ \text{Boolean}) = (\lambda \tau. \ \bot)
\]

\[
\text{lemma} \ \text{null-Boolean-def} : (\text{null}:(\forall \alpha) \ \text{Boolean}) = (\lambda \tau. \ \bot)
\]

\[
\text{definition} \ \text{true} :: (\forall \alpha) \ \text{Boolean} \\
\text{where} \quad \text{true} \equiv \lambda \tau. \ [\![\text{True}]\!]
\]

\[
\text{definition} \ \text{false} :: (\forall \alpha) \ \text{Boolean} \\
\text{where} \quad \text{false} \equiv \lambda \tau. \ [\![\text{False}]\!]
\]

\[
\text{lemma} \ \text{bool-split}: X \ \tau = \ \text{invalid} \ \tau \ \lor \ X \ \tau = \ \text{null} \ \tau \ \lor \\
X \ \tau = \ \text{true} \ \tau \ \lor \ X \ \tau = \ \text{false} \ \tau
\]

\[
\text{lemma} \ [\text{simp}]: \ \text{false} \ (a, b) = [\![\text{False}]\!]
\]

\[
\text{lemma} \ [\text{simp}]: \ \text{true} \ (a, b) = [\![\text{True}]\!]
\]

\[
\text{lemma} \ \text{textbook-true}: \ I[\text{true}] \ \tau = [\![\text{True}]\!]
\]
**lemma textbook-false:** $I[\text{false}] \tau = \lfloor \lfloor \text{False} \rfloor \rfloor$

\begin{tabular}{ll}

<table>
<thead>
<tr>
<th>Name</th>
<th>Theorem</th>
</tr>
</thead>
<tbody>
<tr>
<td>textbook-invalid</td>
<td>$I[\text{invalid}] \ ?\tau = \text{OCL-core.bot-class.bot}$</td>
</tr>
<tr>
<td>textbook-null-fun</td>
<td>$I[\text{null}] \ ?\tau = \text{null}$</td>
</tr>
<tr>
<td>textbook-true</td>
<td>$I[\text{true}] \ ?\tau = \lfloor \lfloor \text{True} \rfloor \rfloor$</td>
</tr>
<tr>
<td>textbook-false</td>
<td>$I[\text{false}] \ ?\tau = \lfloor \lfloor \text{False} \rfloor \rfloor$</td>
</tr>
</tbody>
</table>

Table 3.1.: Basic semantic constant definitions of the logic (except null)

### 3.2.2. Validity and Definedness

However, this has also the consequence that core concepts like definedness, validness and even cp have to be redefined on this type class:

**definition** valid :: \('A, 'a::null)val => ('A)Boolean ('v - [100]100)

where \( v X \equiv \lambda \tau . \text{if } X \tau = \text{bot } \tau \text{ then false } \tau \text{ else true } \tau \)

**lemma** valid1 [simp]: \( v \text{ invalid = false} \)

\( \langle \text{proof} \rangle \)

**lemma** valid2 [simp]: \( v \text{ null = true} \)

\( \langle \text{proof} \rangle \)

**lemma** valid3 [simp]: \( v \text{ true = true} \)

\( \langle \text{proof} \rangle \)

**lemma** valid4 [simp]: \( v \text{ false = true} \)

\( \langle \text{proof} \rangle \)

**lemma** cp-valid: \( (v X) \tau = (v (\lambda \cdot X \tau)) \tau \)

\( \langle \text{proof} \rangle \)

**definition** defined :: \('A, 'a::null)val => ('A)Boolean (\delta - [100]100)

where \( \delta X \equiv \lambda \tau . \text{if } X \tau = \text{bot } \tau \lor X \tau = \text{null } \tau \text{ then false } \tau \text{ else true } \tau \)

The generalized definitions of invalid and definedness have the same properties as the old ones:

**lemma** defined1 [simp]: \( \delta \text{ invalid = false} \)

\( \langle \text{proof} \rangle \)
The definitions above for the constants defined and valid can be rewritten into the conventional semantic “textbook” format as follows:

\[
\text{lemma textbook-defined: } I \delta X \tau = (\text{if } I X \tau = I \text{bot } \tau \lor I X \tau = I \text{null } \tau \text{ then } I \text{false } \tau \text{ else } I \text{true } \tau)
\]

\[
\text{lemma textbook-valid: } I v X \tau = (\text{if } I X \tau = I \text{bot } \tau \text{ then } I \text{false } \tau \text{ else } I \text{true } \tau)
\]

Table 3.2 and Table 3.3 summarize the results of this section.

<table>
<thead>
<tr>
<th>Name</th>
<th>Theorem</th>
</tr>
</thead>
<tbody>
<tr>
<td>textbook-defined</td>
<td>[ I[\delta X] \tau = (\text{if } I[X] \tau = I[\text{bot}] \tau \lor I[X] \tau = I[\text{null}] \tau \text{ then } I[\text{false}] \tau \text{ else } I[\text{true}] \tau)]</td>
</tr>
<tr>
<td>textbook-valid</td>
<td>[ I[v X] \tau = (\text{if } I[X] \tau = I[\text{OCL-core.bot-class}] \tau \lor I[X] \tau = I[\text{OCL-core.bot-class}] \tau \text{ then } I[\text{false}] \tau \text{ else } I[\text{true}] \tau)]</td>
</tr>
</tbody>
</table>

Table 3.2.: Basic predicate definitions of the logic.
Table 3.3.: Laws of the basic predicates of the logic.

<table>
<thead>
<tr>
<th>Name</th>
<th>Theorem</th>
</tr>
</thead>
<tbody>
<tr>
<td>defined1</td>
<td>( \delta \text{ invalid} = false )</td>
</tr>
<tr>
<td>defined2</td>
<td>( \delta \text{ null} = false )</td>
</tr>
<tr>
<td>defined3</td>
<td>( \delta \text{ true} = true )</td>
</tr>
<tr>
<td>defined4</td>
<td>( \delta \text{ false} = true )</td>
</tr>
<tr>
<td>defined5</td>
<td>( \delta \delta ?X = true )</td>
</tr>
<tr>
<td>defined6</td>
<td>( \delta \upsilon ?X = true )</td>
</tr>
</tbody>
</table>

3.3. The Equalities of OCL

The OCL contains a particular version of equality, written in Standard documents \( _=_ \) and \( _\neq_ \) for its negation, which is referred as weak referential equality hereafter and for which we use the symbol \( \equiv \) throughout the formal part of this document. Its semantics is motivated by the desire of fast execution, and similarity to languages like Java and C, but does not satisfy the needs of logical reasoning over OCL expressions and specifications. We therefore introduce a second equality, referred as strong equality or logical equality and written \( \equiv \) which is not present in the current standard but was discussed in prior texts on OCL like the Amsterdam Manifesto \([19]\) and was identified as desirable extension of OCL in the Aachen Meeting \([15]\) in the future 2.5 OCL Standard. The purpose of strong equality is to define and reason over OCL. It is therefore a natural task in Featherweight OCL to formally investigate the somewhat quite complex relationship between these two.

Strong equality has two motivations: a pragmatic one and a fundamental one.

1. The pragmatic reason is fairly simple: users of object-oriented languages want something like a “shallow object value equality”. You will want to say \( a.\text{boss} \equiv b.\text{boss}@pre \) instead of
   \[
   a.\text{boss} \equiv b.\text{boss}@pre \text{ and (« just the pointers are equal! »)}
   a.\text{boss}.\text{name} \equiv b.\text{boss}@pre.\text{name}@pre \text{ and}
   a.\text{boss}.\text{age} \equiv b.\text{boss}@pre.\text{age}@pre
   \]

Breaking a shallow-object equality down to referential equality of attributes is cumbersome, error-prone, and makes specifications difficult to extend (add for example an attribute sex to your class, and check in your OCL specification everywhere that you did it right with your simulation of strong equality). Therefore, languages like Java offer facilities to handle two different equalities, and it is problematic even in an execution oriented specification language to ignore shallow object equality because it is so common in the code.

2. The fundamental reason goes as follows: whatever you do to reason consistently over a language, you need the concept of equality: you need to know what expressions can be replaced by others because they mean the same thing. People call
this also “Leibniz Equality” because this philosopher brought this principle first explicitly to paper and shed some light over it. It is the theoretic foundation of what you do in an optimizing compiler: you replace expressions by equal ones, which you hope are easier to evaluate. In a typed language, strong equality exists uniformly over all types, it is “polymorphic” \( \_ = \_ :: \alpha \ast \alpha \rightarrow bool \)—this is the way that equality is defined in HOL itself. We can express Leibniz principle as one logical rule of surprising simplicity and beauty:

$$s = t \implies P(s) = P(t)$$  \hspace{1cm} (3.1)

“Whenever we know, that \( s \) is equal to \( t \), we can replace the sub-expression \( s \) in a term \( P \) by \( t \) and we have that the replacement is equal to the original.”

While weak referential equality is defined to be strict in the OCL standard, we will define strong equality as non-strict. It is quite nasty (but not impossible) to define the logical equality in a strict way (the substitutivity rule above would look more complex), however, whenever references were used, strong equality is needed since references refer to particular states (pre or post), and that they mean the same thing can therefore not be taken for granted.

### 3.3.1. Definition

The strict equality on basic types (actually on all types) must be exceptionally defined on \textit{null}—otherwise the entire concept of null in the language does not make much sense. This is an important exception from the general rule that null arguments—especially if passed as “self”-argument—lead to invalid results.

We define strong equality extremely generic, even for types that contain a \textit{null} or \perp element. Strong equality is simply polymorphic in Featherweight OCL, i.e., is defined identical for all types in OCL and HOL.

\begin{align*}
\text{definition} & \quad \text{StrongEq} :: \forall \alpha. \alpha \times \alpha \rightarrow \text{Boolean} \quad \text{infixl} \equiv 30 \\
\text{where} & \quad X \triangleq Y \equiv \lambda \tau. \lfloor \lfloor X \tau = Y \tau \rfloor \rfloor
\end{align*}

From this follow already elementary properties like:

\begin{align*}
\text{lemma} \quad & \quad \text{simp,code-unfold}: \ (\text{true} \triangleq \text{false}) = \text{false} \\
\text{lemma} \quad & \quad \text{simp,code-unfold}: \ (\text{false} \triangleq \text{true}) = \text{false}
\end{align*}

\begin{proof}
\end{proof}

In contrast, referential equality behaves differently for all types—on value types, it is basically strong equality for defined values, but on object types it will compare references—we introduce it as an \textit{overloaded} concept and will handle it for each type instance individually.

\begin{align*}
\text{consts} & \quad \text{StrictRefEq} :: \forall \alpha. \alpha \rightarrow \alpha \rightarrow \text{Boolean} \quad \text{infixl} \equiv 30
\end{align*}
Here is a first instance of a definition of weak equality—for the special case of the type \( \forall \) Boolean, it is just the strict extension of the logical equality:

\[
defs \quad \text{StrictRefEq}_{\text{Boolean}}[\text{code-unfold}] : \\
(x : (\forall \text{Boolean}) \mapsto y \equiv \lambda \tau. \; \text{if } (v \ x) \tau = \text{true} \land (v \ y) \tau = \text{true} \tau \; \text{then } (x \triangleq y) \tau \; \text{else invalid } \tau
\]

which implies elementary properties like:

\[
\text{lemma } \text{[simp,code-unfold]} : (\text{true} \triangleq \text{false}) = \text{false} \\
\langle \text{proof} \rangle
\]

\[
\text{lemma } \text{[simp,code-unfold]} : (\text{false} \triangleq \text{true}) = \text{false} \\
\langle \text{proof} \rangle
\]

\[
\text{lemma } \text{[simp,code-unfold]} : (\text{invalid} \triangleq \text{false}) = \text{invalid} \\
\langle \text{proof} \rangle
\]

\[
\text{lemma } \text{[simp,code-unfold]} : (\text{true} \triangleq \text{true}) = \text{true} \\
\langle \text{proof} \rangle
\]

Thus, the weak equality is not reflexive.

\[
\text{lemma } \text{null-non-false } \text{[simp,code-unfold]} : (\text{null} \triangleq \text{false}) = \text{false} \\
\langle \text{proof} \rangle
\]

\[
\text{lemma } \text{null-non-true } \text{[simp,code-unfold]} : (\text{null} \triangleq \text{true}) = \text{false} \\
\langle \text{proof} \rangle
\]

\[
\text{lemma } \text{false-non-null } \text{[simp,code-unfold]} : (\text{false} \triangleq \text{null}) = \text{false} \\
\langle \text{proof} \rangle
\]

\[
\text{lemma } \text{true-non-null } \text{[simp,code-unfold]} : (\text{true} \triangleq \text{null}) = \text{false} \\
\langle \text{proof} \rangle
\]

3.3.2. Fundamental Predicates on Strong Equality

Equality reasoning in OCL is not humpty dumpty. While strong equality is clearly an equivalence:

\[
\text{lemma } \text{StrongEq-refl } \text{[simp]} : (X \triangleq X) = \text{true} \\
\langle \text{proof} \rangle
\]

\[
\text{lemma } \text{StrongEq-sym} : (X \triangleq Y) = (Y \triangleq X) \\
\langle \text{proof} \rangle
\]
lemma \textit{StrongEq-trans-strong} [simp]:
\begin{align*}
\text{assumes } & A: \ (X \triangleq Y) = \text{true} \\
\text{and } & B: \ (Y \triangleq Z) = \text{true} \\
\text{shows } & (X \triangleq Z) = \text{true}
\end{align*}
\begin{proof}

it is only in a limited sense a congruence, at least from the point of view of this
semantic theory. The point is that it is only a congruence on OCL expressions, not
arbitrary HOL expressions (with which we can mix Featherweight OCL expressions). A
semantic—not syntactic—characterization of OCL expressions is that they are \textit{context-
passing} or \textit{context-invariant}, i.e., the context of an entire OCL expression, i.e. the pre
and post state it refers to, is passed constantly and unmodified to the sub-expressions,
i.e., all sub-expressions inside an OCL expression refer to the same context. Expressed
formally, this boils down to:

\begin{align*}
\text{lemma } \textit{StrongEq-subst} & : \\
\text{assumes } & cp: \ \forall X. \ P(X)\tau = P(\lambda - X \tau)\tau \\
\text{and } & eq: \ (X \triangleq Y)\tau = \text{true} \tau \\
\text{shows } & (P X \triangleq P Y)\tau = \text{true} \tau
\end{align*}
\begin{proof}

\begin{align*}
\text{lemma } \textit{defined7[simp]} & : \delta (X \triangleq Y) = \text{true} \\
\text{lemma } \textit{valid7[simp]} & : v (X \triangleq Y) = \text{true} \\
\text{lemma } \textit{cp-StrongEq} & : (X \triangleq Y) \tau = ((\lambda - X \tau) \triangleq (\lambda - Y \tau)) \tau
\end{align*}
\begin{proof}

3.4. Logical Connectives and their Universal Properties

It is a design goal to give OCL a semantics that is as closely as possible to a “logical
system” in a known sense; a specification logic where the logical connectives can not be
understood other that having the truth-table aside when reading fails its purpose in our
view.

Practically, this means that we want to give a definition to the core operations to be
as close as possible to the lattice laws; this makes also powerful symbolic normalization
of OCL specifications possible as a pre-requisite for automated theorem provers. For
example, it is still possible to compute without any definedness and validity reasoning the
DNF of an OCL specification; be it for test-case generations or for a smooth transition to
a two-valued representation of the specification amenable to fast standard SMT-solvers,
for example.

Thus, our representation of the OCL is merely a 4-valued Kleene-Logics with \textit{invalid}
as least, \textit{null} as middle and \textit{true} resp. \textit{false} as unrelated top-elements.

\textbf{definition} \textit{OclNot} :: (\forall)\textbf{Boolean} \Rightarrow (\forall)\textbf{Boolean} (not)
where \( \text{not } X \equiv \lambda \tau \ . \ \text{case } X \ \tau \ \text{of} \)
\[
\bot \Rightarrow \bot \\
| \bot \Rightarrow \bot \\
| [\tau] \Rightarrow [\neg \tau]
\]

with term "not" we can express the notation:

**syntax**

\( \text{notequal} :: (\forall) \text{Boolean} \Rightarrow (\forall) \text{Boolean} \Rightarrow (\forall) \text{Boolean} \) (infix \(<\sim\) 40)

**translations**

\( a <> b ::= \text{CONST OclNot}(a \equiv b) \)

**lemma** \( \text{cp-OclNot}: (\text{not } X)\tau = (\text{not } (\lambda . X \ \tau)) \tau \)

**lemma** \( \text{OclNot1}[\text{simp}]: \text{not invalid} = \text{invalid} \)

**lemma** \( \text{OclNot2}[\text{simp}]: \text{not null} = \text{null} \)

**lemma** \( \text{OclNot3}[\text{simp}]: \text{not true} = \text{false} \)

**lemma** \( \text{OclNot4}[\text{simp}]: \text{not false} = \text{true} \)

**lemma** \( \text{OclNot-not}[\text{simp}]: \text{not } (\text{not } X) = X \)

**lemma** \( \text{OclNot-inject}: \wedge x y. \text{not } x = \text{not } y \implies x = y \)

**definition** \( \text{OclAnd :: } [(\forall) \text{Boolean}, (\forall) \text{Boolean}] \Rightarrow (\forall) \text{Boolean} \) (infixl and 30)

where \( X \text{ and } Y \equiv (\lambda \tau . \ \text{case } X \ \tau \ \text{of} \)
\[
[\text{False}] \Rightarrow [\text{False}] \\
| \bot \Rightarrow (\text{case } Y \ \tau \ \text{of} \)
\[
[\text{False}] \Rightarrow [\text{False}] \\
| \bot \Rightarrow [\bot] \\
| [\downarrow] \Rightarrow (\text{case } Y \ \tau \ \text{of} \)
\[
[\text{False}] \Rightarrow [\text{False}] \\
| \bot \Rightarrow [\bot] \\
| [\text{True}] \Rightarrow Y \ \tau)
\]

Note that \( \text{not} \) is \( \text{not} \) defined as a strict function; proximity to lattice laws implies that we \textit{need} a definition of \( \text{not} \) that satisfies \( \text{not}(\text{not}(x))=x \).
In textbook notation, the logical core constructs \textit{not} and \textit{op and} were represented as follows:

**Lemma textbook-OclNot:**

\[
I[\text{not}(X)] \tau = \begin{cases} 
\bot \Rightarrow \bot 
\end{cases}
\]

\[
| \ [ \bot ] \Rightarrow [ \bot ] 
| \ [ [ x ] ] \Rightarrow [ [ \neg x ] ]
\]

\[\langle \text{proof} \rangle\]

**Lemma textbook-OclAnd:**

\[
I[X \text{ and } Y] \tau = (\text{case } I[X] \tau \text{ of} \\
\bot \Rightarrow (\text{case } I[Y] \tau \text{ of} \\
\bot \Rightarrow \bot 
| \ [ \bot ] \Rightarrow \bot 
| \ [ [ \text{True} ] ] \Rightarrow \bot 
| \ [ [ \text{False} ] ] \Rightarrow [ [ \text{False} ] ])
| \ [ [ \bot ] ] \Rightarrow (\text{case } I[Y] \tau \text{ of} \\
\bot \Rightarrow \bot 
| \ [ \bot ] \Rightarrow [ \bot ] 
| \ [ [ \text{True} ] ] \Rightarrow [ \bot ] 
| \ [ [ \text{False} ] ] \Rightarrow [ [ \text{False} ] ])
| \ [ [ \text{True} ] ] \Rightarrow (\text{case } I[Y] \tau \text{ of} \\
\bot \Rightarrow \bot 
| \ [ [ \bot ] ] \Rightarrow [ [ \bot ] ] 
| \ [ [ \text{False} ] ] \Rightarrow [ [ \text{False} ] ])
| \ [ [ \text{False} ] ] \Rightarrow [ [ \text{False} ] ])
\]

\[\langle \text{proof} \rangle\]

**Definition OclOr:**

\[
\text{OclOr} :: \forall A \, \text{Boolean}, \forall A \, \text{Boolean} \Rightarrow A \, \text{Boolean}
\]

\[\text{infixl or 25}\]

\[\text{where } X \text{ or } Y \equiv \text{not (not } X \text{ and not } Y)\]

**Definition OclImplies:**

\[
\text{OclImplies} :: \forall A \, \text{Boolean}, \forall A \, \text{Boolean} \Rightarrow A \, \text{Boolean}
\]

\[\text{infixl implies 25}\]

\[\text{where } X \text{ implies } Y \equiv \text{not } X \text{ or } Y\]

**Lemma cp-OclAnd:**

\[
(X \text{ and } Y) \tau = (\lambda \cdot X \tau) \text{ and } (\lambda \cdot Y \tau) \tau
\]

\[\langle \text{proof} \rangle\]

**Lemma cp-OclOr:**

\[
(X :: A \, \text{Boolean}) \text{ or } Y \tau = (\lambda \cdot X \tau) \text{ or } (\lambda \cdot Y \tau) \tau
\]

\[\langle \text{proof} \rangle\]

**Lemma cp-OclImplies:**

\[
(X \text{ implies } Y) \tau = (\lambda \cdot X \tau) \text{ implies } (\lambda \cdot Y \tau) \tau
\]

\[\langle \text{proof} \rangle\]

**Lemma OclAnd1[simp]:** (invalid and true) = invalid

\[\langle \text{proof} \rangle\]

**Lemma OclAnd2[simp]:** (invalid and false) = false

\[\langle \text{proof} \rangle\]

**Lemma OclAnd3[simp]:** (invalid and null) = invalid

\[\langle \text{proof} \rangle\]
lemma OclAnd4[simp]: (invalid and invalid) = invalid
⟨proof⟩

lemma OclAnd5[simp]: (null and true) = null
⟨proof⟩

lemma OclAnd6[simp]: (null and false) = false
⟨proof⟩

lemma OclAnd7[simp]: (null and null) = null
⟨proof⟩

lemma OclAnd8[simp]: (null and invalid) = invalid
⟨proof⟩

lemma OclAnd9[simp]: (false and true) = false
⟨proof⟩

lemma OclAnd10[simp]: (false and false) = false
⟨proof⟩

lemma OclAnd11[simp]: (false and null) = false
⟨proof⟩

lemma OclAnd12[simp]: (false and invalid) = false
⟨proof⟩

lemma OclAnd13[simp]: (true and true) = true
⟨proof⟩

lemma OclAnd14[simp]: (true and false) = false
⟨proof⟩

lemma OclAnd15[simp]: (true and null) = null
⟨proof⟩

lemma OclAnd16[simp]: (true and invalid) = invalid
⟨proof⟩

lemma OclAnd-idem[simp]: (X and X) = X
⟨proof⟩

lemma OclAnd-commute: (X and Y) = (Y and X)
⟨proof⟩

lemma OclAnd-false1[simp]: (false and X) = false
⟨proof⟩

lemma OclAnd-false2[simp]: (X and false) = false
⟨proof⟩

lemma OclAnd-true1[simp]: (true and X) = X
⟨proof⟩

lemma OclAnd-true2[simp]: (X and true) = X
⟨proof⟩
lemma OclAnd-bot1 [simp]: \( \forall \tau. X \tau \neq \text{false} \implies (\text{bot and } X) \tau = \text{bot} \tau \)

(\langle proof \rangle)

lemma OclAnd-bot2 [simp]: \( \forall \tau. X \tau \neq \text{false} \implies (X \text{ and bot}) \tau = \text{bot} \tau \)

(\langle proof \rangle)

lemma OclAnd-null1 [simp]: \( \forall \tau. X \tau \neq \text{false} \implies X \tau \neq \text{bot} \implies (\text{null and } X) \tau = \text{null} \tau \)

(\langle proof \rangle)

lemma OclAnd-null2 [simp]: \( \forall \tau. X \tau \neq \text{false} \implies X \tau \neq \text{bot} \implies (X \text{ and null}) \tau = \text{null} \tau \)

(\langle proof \rangle)

lemma OclAnd-assoc: \((X \text{ and } (Y \text{ and } Z)) = (X \text{ and } Y \text{ and } Z)\)

(\langle proof \rangle)

lemma OclOr1 [simp]: \((\text{invalid or true}) = \text{true}\)

(\langle proof \rangle)

lemma OclOr2 [simp]: \((\text{invalid or false}) = \text{invalid}\)

(\langle proof \rangle)

lemma OclOr3 [simp]: \((\text{invalid or null}) = \text{invalid}\)

(\langle proof \rangle)

lemma OclOr4 [simp]: \((\text{invalid or invalid}) = \text{invalid}\)

(\langle proof \rangle)

lemma OclOr5 [simp]: \((\text{null or true}) = \text{true}\)

(\langle proof \rangle)

lemma OclOr6 [simp]: \((\text{null or false}) = \text{null}\)

(\langle proof \rangle)

lemma OclOr7 [simp]: \((\text{null or null}) = \text{null}\)

(\langle proof \rangle)

lemma OclOr8 [simp]: \((\text{null or invalid}) = \text{invalid}\)

(\langle proof \rangle)

lemma OclOr-idem [simp]: \((X \text{ or } X) = X\)

(\langle proof \rangle)

lemma OclOr-commute: \((X \text{ or } Y) = (Y \text{ or } X)\)

(\langle proof \rangle)

lemma OclOr-false1 [simp]: \((\text{false or } Y) = Y\)

(\langle proof \rangle)

lemma OclOr-false2 [simp]: \((Y \text{ or false}) = Y\)

(\langle proof \rangle)

lemma OclOr-true1 [simp]: \((\text{true or } Y) = \text{true}\)

(\langle proof \rangle)
lemma OclOr-true2: \((Y \lor \text{true})\) = \text{true}
\langle proof \rangle

lemma OclOr-bot1 [simp]: \(\forall \tau. X \tau \neq \text{true} \implies (\text{bot} \lor X) \tau = \text{bot} \tau\)
\langle proof \rangle

lemma OclOr-bot2 [simp]: \(\forall \tau. X \tau \neq \text{true} \implies (X \lor \text{bot}) \tau = \text{bot} \tau\)
\langle proof \rangle

lemma OclOr-null1 [simp]: \(\forall \tau. X \tau \neq \text{true} \implies X \tau \neq \text{bot} \implies (\text{null} \lor X) \tau = \text{null} \tau\)
\langle proof \rangle

lemma OclOr-null2 [simp]: \(\forall \tau. X \tau \neq \text{true} \implies X \tau \neq \text{bot} \implies (X \lor \text{null}) \tau = \text{null} \tau\)
\langle proof \rangle

lemma OclOr-assoc: \((X \lor (Y \lor Z)) = (X \lor Y \lor Z)\)
\langle proof \rangle

lemma OclImplies-true: \((X \implies \text{true}) = \text{true}\)
\langle proof \rangle

lemma deMorgan1: \(\neg (X \land Y) = (\neg X \lor \neg Y)\)
\langle proof \rangle

lemma deMorgan2: \(\neg (X \lor Y) = (\neg X \land \neg Y)\)
\langle proof \rangle

3.5. A Standard Logical Calculus for OCL

definition OclValid :: \(\langle\forall A\rangle st, \langle\forall A\rangle \text{Boolean} \Rightarrow \text{bool} ((1 \cdot - ) = (- )) 50\)
where \(\tau \models P \equiv ((P \tau) = \text{true} \tau)\)

value \(\tau \models \text{true} <> \text{false}\)
value \(\tau \models \text{false} <> \text{true}\)

3.5.1. Global vs. Local Judgements

lemma transform1: \(P = \text{true} \implies \tau \models P\)
\langle proof \rangle

lemma transform1-rev: \(\forall \tau. \tau \models P \implies P = \text{true}\)
\langle proof \rangle

lemma transform2: \((P = Q) \implies ((\tau \models P) = (\tau \models Q))\)
\langle proof \rangle

lemma transform2-rev: \(\forall \tau. (\tau \models \delta P) \land (\tau \models \delta Q) \land (\tau \models P) = (\tau \models Q) \implies P = Q\)
\langle proof \rangle

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However, certain properties (like transitivity) can not be transformed from the global level to the local one, they have to be re-proven on the local level.

**lemma**
**assumes** $H : P = true \implies Q = true$
**shows** $\tau \models P \implies \tau \models Q$

### 3.5.2. Local Validity and Meta-logic

**lemma** foundation1 ![simp]: $\tau \models true$

**lemma** foundation2 ![simp]: $\neg(\tau \models false)$

**lemma** foundation3 ![simp]: $\neg(\tau \models invalid)$

**lemma** foundation4 ![simp]: $\neg(\tau \models null)$

**lemma** bool-split-local ![simp]:
$$(\tau \models (x \triangleq invalid)) \lor (\tau \models (x \triangleq null)) \lor (\tau \models (x \triangleq true)) \lor (\tau \models (x \triangleq false))$$

**lemma** def-split-local:
$$(\tau \models \delta x) = ((\neg(\tau \models (x \triangleq invalid))) \land (\neg(\tau \models (x \triangleq null))))$$

**lemma** foundation5:
$$\tau \models (P \land Q) \implies (\tau \models P) \land (\tau \models Q)$$

**lemma** foundation6:
$$\tau \models P \implies \tau \models \delta P$$

**lemma** foundation7 ![simp]:
$$(\tau \models not (\delta x)) = (\neg (\tau \models \delta x))$$

**lemma** foundation7′ ![simp]:
$$(\tau \models not (\upsilon x)) = (\neg (\tau \models \upsilon x))$$

Key theorem for the $\delta$-closure: either an expression is defined, or it can be replaced
(substituted via StrongEq-L-subst2; see below) by invalid or null. Strictness-reduction rules will usually reduce these substituted terms drastically.

**Lemma foundation8:**
$(\tau \models \delta \ x) \lor (\tau \models (x \triangleq \text{invalid})) \lor (\tau \models (x \triangleq \text{null}))$

**Proof**

**Lemma foundation9:**
$\tau \models \delta \ x \Rightarrow (\tau \models \neg \ x) = (\neg (\tau \models x))$

**Proof**

**Lemma foundation10:**
$\tau \models \delta \ x \Rightarrow \tau \models \delta \ y \Rightarrow (\tau \models (x \land y)) = ((\tau \models x) \land (\tau \models y))$

**Proof**

**Lemma foundation11:**
$\tau \models \delta \ x \Rightarrow \tau \models \delta \ y \Rightarrow (\tau \models (x \lor y)) = ((\tau \models x) \lor (\tau \models y))$

**Proof**

**Lemma foundation12:**
$\tau \models \delta \ x \Rightarrow \tau \models \delta \ y \Rightarrow (\tau \models (x \implies y)) = ((\tau \models x) \rightarrow (\tau \models y))$

**Proof**

**Lemma foundation13:**
$(\tau \models A \triangleq \text{true}) = (\tau \models A)$

**Proof**

**Lemma foundation14:**
$(\tau \models A \triangleq \text{false}) = (\tau \models \neg A)$

**Proof**

**Lemma foundation15:**
$(\tau \models A \triangleq \text{invalid}) = (\tau \models \neg (\forall A))$

**Proof**

**Lemma foundation16:**
$\tau \models (\delta \ X) = (X \tau \neq \text{bot} \land X \tau \neq \text{null})$

**Proof**

**Lemma foundation16':**
$(\tau \models (\delta \ X)) = (X \tau \neq \text{invalid} \land X \tau \neq \text{null} \tau)$

**Proof**

**Lemmas foundation17 = foundation16[THEN iffD1,standard]**

**Lemmas foundation17' = foundation16'[THEN iffD1,standard]**

**Lemma foundation18:**
$\tau \models (\forall \ X) = (X \tau \neq \text{invalid} \tau)$
lemma foundation18': \( \tau \models (\nu X) = (X \tau \neq \bot) \)

lemmas foundation19 = foundation18 THEN iffD1, standard

lemma foundation20: \( \tau \models (\delta X) \Longrightarrow \tau \models v X \)

lemma foundation21: (not A \( \triangleq \) not B) = (A \( \triangleq \) B)

lemma foundation22: (\( \tau \models (X \triangleq Y) \)) = (\( X \tau = Y \tau \))

lemma foundation23: (\( \tau \models P \)) = (\( \tau \models (\lambda \cdot. P \tau) \))

lemmas cp-validity = foundation23

lemma foundation24: (\( \tau \models (not(X \triangleq Y)) \)) = (\( X \tau \neq Y \tau \))

lemma defined-not-I : \( \tau \models \delta (x) \Longrightarrow \tau \models \delta (not x) \)

lemma valid-not-I : \( \tau \models v (x) \Longrightarrow \tau \models v (not x) \)

lemma defined-and-I : \( \tau \models \delta (x) \Longrightarrow \tau \models \delta (y) \Longrightarrow \tau \models \delta (x and y) \)

lemma valid-and-I : \( \tau \models v (x) \Longrightarrow \tau \models v (y) \Longrightarrow \tau \models v (x and y) \)

3.5.3. Local Judgements and Strong Equality

lemma StrongEq-L-refl: \( \tau \models (x \triangleq x) \)

lemma StrongEq-L-sym: \( \tau \models (x \triangleq y) \Longrightarrow \tau \models (y \triangleq x) \)

lemma StrongEq-L-trans: \( \tau \models (x \triangleq y) \Longrightarrow \tau \models (y \triangleq z) \Longrightarrow \tau \models (x \triangleq z) \)
In order to establish substitutivity (which does not hold in general HOL formulas) we introduce the following predicate that allows for a calculus of the necessary side-conditions.

**definition** \( cp : ((\forall', a') \text{ val} \Rightarrow (\forall', \beta') \text{ val}) \Rightarrow \text{ bool} \)

**where** \( cp \ P \equiv (\exists \ f. \ \forall \ X \ \tau. \ P \ X \ \tau = f (X \ \tau) \ \tau) \)

The rule of substitutivity in Featherweight OCL holds only for context-passing expressions, i.e. those that pass the context \( \tau \) without changing it. Fortunately, all operators of the OCL language satisfy this property (but not all HOL operators).

**lemma** \( \text{StrongEq-L-subst1} : \ \forall \ \tau. \ cp \ P = \Rightarrow \ \tau \models (x \triangleq y) = \Rightarrow \ \tau \models (P \ x \triangleq P \ y) \) (\proof\)

**lemma** \( \text{StrongEq-L-subst2} : \ \forall \ \tau. \ cp \ P = \Rightarrow \ \tau \models (x \triangleq y) = \Rightarrow \ \tau \models (P \ x) = \Rightarrow \ \tau \models (P \ y) \) (\proof\)

**lemma** \( \text{StrongEq-L-subst2-rev} : \ \tau \models y \triangleq x = \Rightarrow \ \tau \models cp \ P = \Rightarrow \ \tau \models P \ x = \Rightarrow \ \tau \models P \ y \) (\proof\)

**lemma** \( \text{StrongEq-L-subst3} : \ \text{assumes} \ cp : cp \ P \ \text{and} \ eq : \ \tau \models x \triangleq y \ \text{shows} \ \tau \models (P \ x) = (\tau \models P \ y) \) (\proof\)

**lemma** \( \text{cpI1} : \ (\forall \ X \ \tau. \ f X \ \tau = f (\lambda. \ X \ \tau) \ \tau) = \Rightarrow \ cp \ P = \Rightarrow \ cp (\lambda X. \ f (P \ X)) \) (\proof\)

**lemma** \( \text{cpI2} : \ (\forall \ X \ Y \ \tau. \ f X \ Y \ \tau = f (\lambda. \ X \ \tau) (\lambda. \ Y \ \tau) \ \tau) = \Rightarrow \ cp \ P = \Rightarrow \ cp \ Q = \Rightarrow \ cp (\lambda X. \ f (P \ X) (Q \ X)) \) (\proof\)

**lemma** \( \text{cpI3} : \ (\forall \ X \ Y \ Z \ \tau. \ f X \ Y \ Z \ \tau = f (\lambda. \ X \ \tau) (\lambda. \ Y \ \tau) (\lambda. \ Z \ \tau) \ \tau) = \Rightarrow \ cp \ P = \Rightarrow \ cp \ Q = \Rightarrow \ cp \ R = \Rightarrow \ cp (\lambda X. \ f (P \ X) (Q \ X) (R \ X)) \) (\proof\)

**lemma** \( \text{cpI4} : \ (\forall \ W \ X \ Y \ Z \ \tau. \ f W \ X \ Y \ Z \ \tau = f (\lambda. \ W \ \tau) (\lambda. \ X \ \tau) (\lambda. \ Y \ \tau) (\lambda. \ Z \ \tau) \ \tau) = \Rightarrow \ cp \ P = \Rightarrow \ cp \ Q = \Rightarrow \ cp \ R = \Rightarrow \ cp \ S = \Rightarrow \ cp (\lambda X. \ f (P \ X) (Q \ X) (R \ X) (S \ X)) \) (\proof\)

**lemma** \( \text{cp-const} : \ cp (\lambda \ c) \) (\proof\)
lemma \( \text{cp-id} : \quad \text{cp}(\lambda X. X) \)
(proof)

lemmas \( \text{cp-intro}[\text{intro}, \text{simp}, \text{code-unfold}] = \)
cp-const
cp-id
cp-defined[THEN allI[THEN allI[THEN cpI1], of defined]]
cp-valid[THEN allI[THEN allI[THEN cpI1], of valid]]
cp-OclNot[THEN allI[THEN allI[THEN cpI1], of not]]
cp-OclAnd[THEN allI[THEN allI[THEN cpI1], of op and]]
cp-OclOr[THEN allI[THEN allI[THEN cpI1], of op or]]
cp-OclImplies[THEN allI[THEN allI[THEN cpI1], of op implies]]
cp-StrongEq[THEN allI[THEN allI[THEN cpI1], of StrongEq]]

3.5.4. Laws to Establish Definedness (\(\delta\)-closure)

For the logical connectives, we have — beyond \(\tau \models \hat{P} \Rightarrow \delta \hat{P}\) — the following facts:

lemma \(\text{OclNot-defargs}\):
\(\tau \models (not \, P) \Rightarrow \tau \models \delta \, P\)
(proof)

lemma \(\text{OclNot-contrapos-nn}\):
assumes \(\tau \models \delta \, A\)
assumes \(\tau \models not \, B\)
shows \(\tau \models A \Rightarrow \tau \models B\)
(proof)

So far, we have only one strict Boolean predicate (\(\delta\)-family): the strict equality.

3.6. Miscellaneous

3.6.1. OCL’s if then else endif

definition \(\text{OclIf} :: \[(\pi) \, \text{Boolean} \, , \,(\pi, \alpha::\text{null}) \, \text{val} , \,(\pi, \alpha) \, \text{val}] \Rightarrow (\pi, \alpha) \, \text{val}\)
(if (-) then (-) else (-) endif \(10,10,10\)\)\(50\)
where (if \(C\) then \(B_1\) else \(B_2\) endif) = (\(\lambda \tau. \text{if} (\delta \, C) \, \tau = \text{true} \, \tau\)
then (if \((C \, \tau) = \text{true} \, \tau\)
then \(B_1 \, \tau\)
else \(B_2 \, \tau\)
else invalid \(\tau\))

lemma \(\text{cp-OclIf}):(\text{if} \, C \, \text{then} \, B_1 \, \text{else} \, B_2 \, \text{endif}) \, \tau = \)
(if \((\lambda \cdot. \, C \, \tau) \, \text{then} \,(\lambda \cdot. \, B_1 \, \tau) \, \text{else} \,(\lambda \cdot. \, B_2 \, \tau) \, \text{endif}) \, \tau)
lemmas cp-intro' [intro, simp, code-unfold] =
  cp-intro
  cp-OclIf[THEN allI[THEN allI[THEN allI[THEN allI[THEN cpI?]]]], of OclIf]]

lemma OclIf-invalid [simp]: (if invalid then B₁ else B₂ endif) = invalid
  ⟨proof⟩

lemma OclIf-null [simp]: (if null then B₁ else B₂ endif) = invalid
  ⟨proof⟩

lemma OclIf-true [simp]: (if true then B₁ else B₂ endif) = B₁
  ⟨proof⟩

lemma OclIf-true' [simp]: τ ⊢ P ⇒ (if P then B₁ else B₂ endif)τ = B₁ τ
  ⟨proof⟩

lemma OclIf-false [simp]: (if false then B₁ else B₂ endif) = B₂
  ⟨proof⟩

lemma OclIf-false' [simp]: τ ⊢ not P ⇒ (if P then B₁ else B₂ endif)τ = B₂ τ
  ⟨proof⟩

lemma OclIf-idem1 [simp]: (if δ X then A else A endif) = A
  ⟨proof⟩

lemma OclIf-idem2 [simp]: (if ν X then A else A endif) = A
  ⟨proof⟩

lemma OclNot-if [simp]:
  not (if P then C else E endif) = (if P then not C else not E endif)
  ⟨proof⟩

3.6.2. A Side-calculus for (Boolean) Constant Terms

definition const X ≡ ∀ τ τ'. X τ = X τ'

lemma const-charn: const X ⊢ X τ = X τ'
  ⟨proof⟩

lemma const-subst:
  assumes const-X: const X
   and const-Y: const Y
   and eq : X τ = Y τ
   and cp-P: cp P
   and pp : P Y τ = P Y τ'

⟨proof⟩
shows \( P \times \tau = P \times \tau' \)

\langle proof \rangle

**lemma** const-imply2 :
**assumes** \( \forall \tau_1 \tau_2. \ P \ \tau_1 = P \ \tau_2 \implies Q \ \tau_1 = Q \ \tau_2 \)
**shows** const \( P \implies \) const \( Q \)
\langle proof \rangle

**lemma** const-imply3 :
**assumes** \( \forall \tau_1 \tau_2. \ P \ \tau_1 = P \ \tau_2 \implies Q \ \tau_1 = Q \ \tau_2 \implies R \ \tau_1 = R \ \tau_2 \)
**shows** const \( P \implies \) const \( Q \implies \) const \( R \)
\langle proof \rangle

**lemma** const-imply4 :
**assumes** \( \forall \tau_1 \tau_2. \ P \ \tau_1 = P \ \tau_2 \implies Q \ \tau_1 = Q \ \tau_2 \implies R \ \tau_1 = R \ \tau_2 \implies S \ \tau_1 = S \ \tau_2 \)
**shows** const \( P \implies \) const \( Q \implies \) const \( R \implies \) const \( S \)
\langle proof \rangle

**lemma** const-lam : const \( (\lambda \cdot \ e) \)
\langle proof \rangle

**lemma** const-true : const true
\langle proof \rangle

**lemma** const-false : const false
\langle proof \rangle

**lemma** const-null : const null
\langle proof \rangle

**lemma** const-invalid : const invalid
\langle proof \rangle

**lemma** const-bot : const bot
\langle proof \rangle

**lemma** const-defined :
**assumes** const \( X \)
**shows** const \( (\delta \ X) \)
\langle proof \rangle

**lemma** const-valid :
**assumes** const \( X \)
**shows** const \( (\upsilon \ X) \)
\langle proof \rangle
lemma const-OclValid1:
assumes const $x$
shows $(\tau \models \delta x) = (\tau' \models \delta x)$
(proof)

lemma const-OclValid2:
assumes const $x$
shows $(\tau \models v x) = (\tau' \models v x)$
(proof)

lemma const-OclAnd:
assumes const $X$
assumes const $X'$
shows const $(X \text{ and } X')$
(proof)

lemma const-OclNot:
assumes const $X$
shows const $(\neg X)$
(proof)

lemma const-OclOr:
assumes const $X$
assumes const $X'$
shows const $(X \text{ or } X')$
(proof)

lemma const-OclImplies:
assumes const $X$
assumes const $X'$
shows const $(X \text{ implies } X')$
(proof)

lemma const-StrongEq:
assumes const $X$
assumes const $X'$
shows const $(X \equal{} X')$
(proof)

lemma const-OclIf:
assumes const $B$
and const $C1$
and const $C2$
shows const $(\text{if } B \text{ then } C1 \text{ else } C2 \text{ endif})$
(proof)
lemmas \( \text{const-ss} = \text{const-bot} \ \text{const-null} \ \text{const-invalid} \ \text{const-false} \ \text{const-true} \ \text{const-lam} \ \text{const-defined} \ \text{const-valid} \ \text{const-StrongEq} \ \text{const-OclNot} \ \text{const-OclAnd} \ \text{const-OclOr} \ \text{const-OclImplies} \ \text{const-OclIf} \)
4. Formalization II: Library Definitions

theory OCL-lib
imports OCL-core
begin

The structure of this chapter roughly follows the structure of Chapter 10 of the OCL
standard [33], which introduces the OCL Library.

4.1. Basic Types: Void and Integer

4.1.1. The Construction of the Void Type
type-synonym (′A) Void = (′A, unit option) val

This minimal OCL type contains only two elements: invalid and null. Void could
initially be defined as unit option option, however the cardinal of this type is more than
two, so it would have the cost to consider Some None and Some (Some ()) seemingly
everywhere.

4.1.2. The Construction of the Integer Type

Since Integer is again a basic type, we define its semantic domain as the valuations over
int option option.
type-synonym (′A) Integer = (′A, int option option) val

Although the remaining part of this library reasons about integers abstractly, we
provide here as example some convenient shortcuts.
definition OclInt0 ::(′A) Integer (0)
where 0 = (λ . . [0::int])
definition OclInt1 ::(′A) Integer (1)
where 1 = (λ . . [1::int])
definition OclInt2 ::(′A) Integer (2)
where 2 = (λ . . [2::int])
definition OclInt3 ::(′A) Integer (3)
where 3 = (λ . . [3::int])
definition OclInt4 ::(′A) Integer (4)
where 4 = (λ . . [4::int])
definition \( \text{OclInt5} :: (\text{hdr} \\text{Integer} (5)) \)
where \( 5 = (\lambda \cdot \lfloor \lfloor 5 :: \text{int} \rfloor \rfloor) \)

definition \( \text{OclInt6} :: (\text{hdr} \\text{Integer} (6)) \)
where \( 6 = (\lambda \cdot \lfloor \lfloor 6 :: \text{int} \rfloor \rfloor) \)

definition \( \text{OclInt7} :: (\text{hdr} \\text{Integer} (7)) \)
where \( 7 = (\lambda \cdot \lfloor \lfloor 7 :: \text{int} \rfloor \rfloor) \)

definition \( \text{OclInt8} :: (\text{hdr} \\text{Integer} (8)) \)
where \( 8 = (\lambda \cdot \lfloor \lfloor 8 :: \text{int} \rfloor \rfloor) \)

definition \( \text{OclInt9} :: (\text{hdr} \\text{Integer} (9)) \)
where \( 9 = (\lambda \cdot \lfloor \lfloor 9 :: \text{int} \rfloor \rfloor) \)

definition \( \text{OclInt10} :: (\text{hdr} \\text{Integer} (10)) \)
where \( 10 = (\lambda \cdot \lfloor \lfloor 10 :: \text{int} \rfloor \rfloor) \)

4.1.3. Validity and Definedness Properties

Lemma \( \delta \) (null :: (\text{hdr} \text{Integer}) = false (proof)

Lemma \( \nu \) (null :: (\text{hdr} \text{Integer}) = true (proof)

Lemma \( \delta \) (\lambda \cdot \lfloor \lfloor n \rfloor \rfloor = true (proof)

Lemma \( \nu \) (\lambda \cdot \lfloor \lfloor n \rfloor \rfloor = true (proof)

4.1.4. Arithmetical Operations on Integer

Definition

Here is a common case of a built-in operation on built-in types. Note that the arguments must be both defined (non-null, non-bot).
Note that we can not follow the lexis of the OCL Standard for Isabelle technical reasons; these operators are heavily overloaded in the HOL library that a further overloading would lead to heavy technical buzz in this document.

definition OclAdd\_Integer :: (′A)\,Integer ⇒ (′A)\,Integer ⇒ (′A)\,Integer (infix ‘+ 40) where
x ‘+ y ≡ λ τ. if (δ x) τ = true τ ∧ (δ y) τ = true τ then ⌊⌊⌈⌈ x τ⌉⌉ + ⌈⌈ y τ⌉⌉⌋⌋ else invalid τ

definition OclLess\_Integer :: (′A)\,Integer ⇒ (′A)\,Integer ⇒ (′A)\,Boolean (infix ‘< 40) where
x ‘< y ≡ λ τ. if (δ x) τ = true τ ∧ (δ y) τ = true τ then ⌊⌊⌈⌈ x τ⌉⌉ < ⌈⌈ y τ⌉⌉⌋⌋ else invalid τ

definition OclLe\_Integer :: (′A)\,Integer ⇒ (′A)\,Integer ⇒ (′A)\,Boolean (infix ‘≤ 40) where
x ‘≤ y ≡ λ τ. if (δ x) τ = true τ ∧ (δ y) τ = true τ then ⌊⌊⌈⌈ x τ⌉⌉ ≤ ⌈⌈ y τ⌉⌉⌋⌋ else invalid τ

Basic Properties

lemma OclAdd\_Integer\,-commute: (X ‘+ Y) = (Y ‘+ X) ⟨proof⟩

Execution with Invalid or Null or Zero as Argument

lemma OclAdd\_Integer\,-strict1[simp,code-unfold] : (x ‘+ invalid) = invalid ⟨proof⟩

lemma OclAdd\_Integer\,-strict2[simp,code-unfold] : (invalid ‘+ x) = invalid ⟨proof⟩

lemma [simp,code-unfold] : (x ‘+ null) = invalid ⟨proof⟩

lemma [simp,code-unfold] : (null ‘+ x) = invalid ⟨proof⟩

lemma OclAdd\_Integer\,-zero1[simp,code-unfold] :
(x ‘+ 0) = (if u x and not (δ x) then invalid else x endif) ⟨proof⟩

lemma OclAdd\_Integer\,-zero2[simp,code-unfold] :
(0 ‘+ x) = (if u x and not (δ x) then invalid else x endif) ⟨proof⟩

Context Passing

lemma cp-OclAdd\_Integer::(X ‘+ Y) τ = ((λ X τ) ‘+ (λ · Y τ)) τ ⟨proof⟩
lemma $cp$-$Ocl$Less$_{Integer}$: $(X \prec Y) \tau = ((\lambda . X \tau) \prec (\lambda . Y \tau)) \tau$

lemma $cp$-$Ocl$Le$_{Integer}$: $(X \preceq Y) \tau = ((\lambda . X \tau) \preceq (\lambda . Y \tau)) \tau$

Test Statements

Here follows a list of code-examples, that explain the meanings of the above definitions by compilation to code and execution to $True$.

value $\tau \models (9 \prec 10)$
value $\tau \models ((4 + 4) \preceq 10)$
value $\neg(\tau \models ((4 + (4 + 4)) \prec 10))$
value $\tau \models \text{not} (v (null + 1))$

4.2. Fundamental Predicates on Basic Types: Strict Equality

4.2.1. Definition

The last basic operation belonging to the fundamental infrastructure of a value-type in OCL is the weak equality, which is defined similar to the $\mathfrak{A}$ Boolean-case as strict extension of the strong equality:

defs
$\text{StrictRefEq}_{Integer}$[code-unfold]:
$(x::(\mathfrak{A})Integer) \equiv y \equiv \lambda \tau. \text{if } (\nu x) \tau = \text{true } \tau \land (\nu y) \tau = \text{true } \tau$
then $(x \equiv y) \tau$
else invalid $\tau$

value $\tau \models 1 <> 2$
value $\tau \models 2 <> 1$
value $\tau \models 2 \equiv 2$

4.2.2. Logic and Algebraic Layer on Basic Types

Validity and Definedness Properties (I)

lemma $\text{StrictRefEq}_{Boolean}$-defined-args-valid:
$\tau \models \delta((x::(\mathfrak{A})Boolean) \equiv y)) = ((\tau \models (v x)) \land (\tau \models (v y)))$

lemma $\text{StrictRefEq}_{Integer}$-defined-args-valid:
$\tau \models \delta((x::(\mathfrak{A})Integer) \equiv y)) = ((\tau \models (v x)) \land (\tau \models (v y)))$

Validity and Definedness Properties (II)

lemma $\text{StrictRefEq}_{Boolean}$-defargs:
$\tau \models ((x::(\mathfrak{A})Boolean) \equiv y) \implies (\tau \models (v x)) \land (\tau \models (v y))$
lemma StrictRefEq\textsubscript{Integer}\textsubscript{-defargs}:
\[ \tau \models ((x::(A)\text{Integer}) \doteq y) \Longrightarrow (\tau \models (v \ x)) \land (\tau \models (v \ y)) \]

Validity and Definedness Properties (III) Miscellaneous

lemma StrictRefEq\textsubscript{Boolean}\textsubscript{-strict}'' : \[ \delta ((x::(A)\text{Boolean}) \doteq y) = (v(x) \land v(y)) \]

lemma StrictRefEq\textsubscript{Integer}\textsubscript{-strict}'' : \[ \delta ((x::(A)\text{Integer}) \doteq y) = (v(x) \land v(y)) \]

lemma StrictRefEq\textsubscript{Integer}\textsubscript{-strict}:
\begin{align*}
\text{assumes} & \quad A: v (x::(A)\text{Integer}) = \text{true} \\
\text{and} & \quad B: v \ y = \text{true} \\
\text{shows} & \quad v (x \doteq y) = \text{true}
\end{align*}

lemma StrictRefEq\textsubscript{Integer}\textsubscript{-strict}':
\begin{align*}
\text{assumes} & \quad A: v (((x::(A)\text{Integer})) \doteq y) = \text{true} \\
\text{shows} & \quad v \ x = \text{true} \land v \ y = \text{true}
\end{align*}

Reflexivity

lemma StrictRefEq\textsubscript{Boolean}\textsubscript{-refl}[simp,code-unfold] : \[ ((x::(A)\text{Boolean}) \doteq x) = (\text{if } (v \ x) \text{ then true else invalid endif}) \]

lemma StrictRefEq\textsubscript{Integer}\textsubscript{-refl}[simp,code-unfold] : \[ ((x::(A)\text{Integer}) \doteq x) = (\text{if } (v \ x) \text{ then true else invalid endif}) \]

Execution with Invalid or Null as Argument

lemma StrictRefEq\textsubscript{Boolean}\textsubscript{-strict1}[simp,code-unfold] : \[ ((x::(A)\text{Boolean}) \doteq \text{invalid}) = \text{invalid} \]

lemma StrictRefEq\textsubscript{Boolean}\textsubscript{-strict2}[simp,code-unfold] : \[ (\text{invalid} \doteq (x::(A)\text{Boolean})) = \text{invalid} \]

lemma StrictRefEq\textsubscript{Integer}\textsubscript{-strict1}[simp,code-unfold] : \[ ((x::(A)\text{Integer}) \doteq \text{invalid}) = \text{invalid} \]

lemma StrictRefEq\textsubscript{Integer}\textsubscript{-strict2}[simp,code-unfold] : \[ (\text{invalid} \doteq (x::(A)\text{Integer})) = \text{invalid} \]
lemma integer-non-null [simp]: \((\lambda.- [[n]]) \equiv (null::(\mathcal{A})\text{Integer})\) = false

lemma null-non-integer [simp]: \(((null::(\mathcal{A})\text{Integer}) \equiv (\lambda.- [[n]]))\) = false

lemma OclInt0-non-null [simp,code-unfold]: \(((0 \div null) = false)\) = false (proof)
lemma null-non-OclInt0 [simp,code-unfold]: \((null = (0))\) = false (proof)
lemma OclInt1-non-null [simp,code-unfold]: \(((1 \div null) = false)\) = false (proof)
lemma null-non-OclInt1 [simp,code-unfold]: \((null = (1))\) = false (proof)
lemma OclInt2-non-null [simp,code-unfold]: \(((2 \div null) = false)\) = false (proof)
lemma null-non-OclInt2 [simp,code-unfold]: \((null = (2))\) = false (proof)
lemma OclInt6-non-null [simp,code-unfold]: \(((6 \div null) = false)\) = false (proof)
lemma null-non-OclInt6 [simp,code-unfold]: \((null = (6))\) = false (proof)
lemma OclInt8-non-null [simp,code-unfold]: \(((8 \div null) = false)\) = false (proof)
lemma null-non-OclInt8 [simp,code-unfold]: \((null = (8))\) = false (proof)
lemma OclInt9-non-null [simp,code-unfold]: \(((9 \div null) = false)\) = false (proof)
lemma null-non-OclInt9 [simp,code-unfold]: \((null = (9))\) = false (proof)

Const

lemma [simp,code-unfold]: const\((0)\) (proof)
lemma [simp,code-unfold]: const\((1)\) (proof)
lemma [simp,code-unfold]: const\((2)\) (proof)
lemma [simp,code-unfold]: const\((6)\) (proof)
lemma [simp,code-unfold]: const\((8)\) (proof)
lemma [simp,code-unfold]: const\((9)\) (proof)

Behavior vs StrongEq

lemma StrictRefEqBoolean-vs-StrongEq:
\(\tau \models (v \ x) \implies \tau \models (v \ y) \implies (\tau \models (((\mathcal{A})\text{Boolean}) \equiv y) \triangleq (x \triangleq y)))\)

lemma StrictRefEqInteger-vs-StrongEq:
\(\tau \models (v \ x) \implies \tau \models (v \ y) \implies (\tau \models (((\mathcal{A})\text{Integer}) \equiv y) \triangleq (x \triangleq y)))\)

Context Passing

lemma cp-StrictRefEqBoolean:
\(((X::(\mathcal{A})\text{Boolean}) \equiv Y) \tau = ((\lambda.- X \tau) \equiv (\lambda.- Y \tau))\) \tau

lemma cp-StrictRefEqInteger:
\(((X::(\mathcal{A})\text{Integer}) \equiv Y) \tau = ((\lambda.- X \tau) \equiv (\lambda.- Y \tau))\) \tau

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lemmas $cp\text{-intro}[\text{intro}, \text{simp}, \text{code-unfold}] =$
$\text{cp}\text{-intro'}$
$\text{cp-StrictRefEq}_{\text{Boolean}}[\text{THEN allI}[\text{THEN allI}[\text{THEN allI}[\text{THEN cpI2}]], \text{of StrictRefEq}]]$
$\text{cp-StrictRefEq}_{\text{Integer}}[\text{THEN allI}[\text{THEN allI}[\text{THEN allI}[\text{THEN cpI2}]], \text{of StrictRefEq}]]$
$\text{cp-OclAdd}_{\text{Integer}}[\text{THEN allI}[\text{THEN allI}[\text{THEN allI}[\text{THEN cpI2}]], \text{of OclAdd}_{\text{Integer}}]]$
$\text{cp-OclLess}_{\text{Integer}}[\text{THEN allI}[\text{THEN allI}[\text{THEN allI}[\text{THEN cpI2}]], \text{of OclLess}_{\text{Integer}}]]$
$\text{cp-OclLe}_{\text{Integer}}[\text{THEN allI}[\text{THEN allI}[\text{THEN allI}[\text{THEN cpI2}]], \text{of OclLe}_{\text{Integer}}]]$

### 4.2.3. Test Statements on Basic Types.

Here follows a list of code-examples, that explain the meanings of the above definitions by compilation to code and execution to $\text{True}$.

Elementary computations on Booleans

- $\text{value } \tau \models v(\text{true})$
- $\text{value } \tau \models \delta(\text{false})$
- $\text{value } \neg(\tau \models \delta(\text{null}))$
- $\text{value } \neg(\tau \models \delta(\text{invalid}))$
- $\text{value } \tau \models v((\text{null}::(\mathbb{A})\text{Boolean}))$
- $\text{value } \neg(\tau \models v(\text{invalid}))$
- $\text{value } \tau \models (\text{true} \text{ and } \text{true})$
- $\text{value } \tau \models ((\text{null} \text{ or } \text{null}) \triangleq \text{null})$
- $\text{value } \tau \models ((\text{null} \triangleq \text{null}) \triangleq \text{null})$
- $\text{value } \tau \models ((\text{true} \triangleq \text{false}) \triangleq \text{false})$
- $\text{value } \tau \models ((\text{invalid} \triangleq \text{false}) \triangleq \text{false})$
- $\text{value } \tau \models ((\text{invalid} \triangleq \text{false}) \triangleq \text{invalid})$

Elementary computations on Integer

- $\text{value } \tau \models v \, 4$
- $\text{value } \tau \models \delta \, 4$
- $\text{value } \tau \models v \, (\text{null}::(\mathbb{A})\text{Integer})$
- $\text{value } \tau \models (\text{invalid} \triangleq \text{invalid})$
- $\text{value } \tau \models (\text{null} \triangleq \text{null})$
- $\text{value } \tau \models (4 \triangleq 4)$
- $\text{value } \neg(\tau \models (9 \triangleq 10))$
- $\text{value } \neg(\tau \models (\text{invalid} \triangleq 10))$
- $\text{value } \neg(\tau \models (\text{null} \triangleq 10))$
- $\text{value } \neg(\tau \models (\text{invalid} \triangleq (\text{invalid}::(\mathbb{A})\text{Integer})))$
- $\text{value } \neg(\tau \models v \, (\text{invalid} \triangleq (\text{invalid}::(\mathbb{A})\text{Integer})))$
- $\text{value } \neg(\tau \models v \, (\text{invalid} \triangleq (<>) \, (\text{invalid}::(\mathbb{A})\text{Integer})))$
- $\text{value } \neg(\tau \models v \, (\text{invalid} \triangleq (<>) \, (\text{invalid}::(\mathbb{A})\text{Integer})))$
- $\text{value } \tau \models (\text{null} \triangleq (\text{null}::(\mathbb{A})\text{Integer}))$
- $\text{value } \tau \models (\text{null} \triangleq (\text{null}::(\mathbb{A})\text{Integer}))$
- $\text{value } \tau \models (4 \triangleq 4)$
- $\text{value } \neg(\tau \models (4 \triangleq 4))$
- $\text{value } \neg(\tau \models (4 \triangleq 10))$
value $\tau \models (4 \not< 10)$  
value $\neg(\tau \models (0 \not< \mathtt{null}))$  
value $\neg(\tau \models (\delta (0 \not< \mathtt{null})))$

4.3. Complex Types: The Set-Collection Type (I) Core

4.3.1. The Construction of the Set Type

no-notation None ($\bot$)

notation bot ($\bot$)

For the semantic construction of the collection types, we have two goals:

1. we want the types to be fully abstract, i.e., the type should not contain junk-elements that are not representable by OCL expressions, and

2. we want a possibility to nest collection types (so, we want the potential to talking about $\mathtt{Set(\mathtt{Set(\mathtt{Sequences(\mathtt{Pairs(\text{X}, \text{Y})})})})}$).

The former principle rules out the option to define $\alpha$ $\mathtt{Set}$ just by $\alpha$ Set, $(\alpha \text{ option option})$ set val. This would allow sets to contain junk elements such as $\{\bot\}$ which we need to identify with undefinedness itself. Abandoning fully abstractness of rules would later on produce all sorts of problems when quantifying over the elements of a type. However, if we build an own type, then it must conform to our abstract interface in order to have nested types: arguments of type-constructors must conform to our abstract interface, and the result type too.

The core of an own type construction is done via a type definition which provides the raw-type $\alpha$ Set-0. It is shown that this type “fits” indeed into the abstract type interface discussed in the previous section.

typedef $\alpha$ Set-0 = \{X::($\alpha$::null) set option option.  
X = bot $\lor$ X = null $\lor$ ($\forall x \in \lceil\lceil X\rceil\rceil$. $x \not= bot$)\}  
\langle proof \rangle

instantiation Set-0 :: (null)bot
begin
  definition bot-Set-0-def: (bot::($\alpha$::null) Set-0) $\equiv$ Abs-Set-0 None
  instance (proof)
end

instantiation Set-0 :: (null)null
begin
  definition null-Set-0-def: (null::($\alpha$::null) Set-0) $\equiv$ Abs-Set-0 \{ None \}
  instance (proof)
end

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... and lifting this type to the format of a valuation gives us:

type-synonym \((\alpha, \alpha) \text{Set} = (\alpha, \alpha \text{Set-0})\) val

4.3.2. Validity and Definedness Properties

Every element in a defined set is valid.

lemma Set-inv-lemma: \(\tau \models (\delta X) \implies \forall x \in [[\text{Rep-Set-0} (X \tau)]] \cdot x \neq \bot\)

lemma Set-inv-lemma':
  assumes x-def : \(\tau \models \delta X\)
  and e-mem : \(e \in [[\text{Rep-Set-0} (X \tau)]]\)
  shows \(\tau \models v (\lambda -. e)\)

lemma abs-rep-simp':
  assumes S-all-def : \((\tau :: \alpha) \models \delta S\)
  shows Abs-Set-0 \(\llbracket [[\text{Rep-Set-0} (S \tau)]] \rrbracket = S \tau\)

lemma S-lift':
  assumes S-all-def : \((\tau :: \alpha \text{st}) \models \delta S\)
  shows \(\exists S'. (\lambda a (\cdot :: \alpha \text{st}). a) \cdot [[\text{Rep-Set-0} (S \tau)]] = (\lambda a (\cdot :: \alpha \text{st}). [a])' S'\)

lemma invalid-set-OclNot-defined [simp,code-unfold]:\(\delta(\text{invalid} :: (\alpha, \alpha::\text{null}) \text{Set}) = \text{false}\)

lemma null-set-OclNot-defined [simp,code-unfold]:\(\delta(\text{null} :: (\alpha, \alpha::\text{null}) \text{Set}) = \text{false}\)

lemma invalid-set-valid [simp,code-unfold]:\(v(\text{invalid} :: (\alpha, \alpha::\text{null}) \text{Set}) = \text{false}\)

lemma null-set-valid [simp,code-unfold]:\(v(\text{null} :: (\alpha, \alpha::\text{null}) \text{Set}) = \text{true}\)

... which means that we can have a type \((\alpha, (\alpha, (\alpha) \text{Set}) \text{Set})\) \text{Set} corresponding exactly to \(\text{Set(\text{Set(\text{Set}}(\alpha \text{Integer})) \text{Set})}\) in OCL notation. Note that the parameter \(\alpha\) still refers to the object universe; making the OCL semantics entirely parametric in the object universe makes it possible to study (and prove) its properties independently from a concrete class diagram.

4.3.3. Constants on Sets

definition mtSet:=(\alpha, \alpha::null) \text{Set} \llbracket\{\}\rrbracket)

where Set{} \(\equiv (\lambda \tau. \text{Abs-Set-0} \llbracket\{\}\rrbracket\) )
lemma mtSet-defined[simp, code-unfold]: δ(Set{}) = true
⟨proof⟩

lemma mtSet-valid[simp, code-unfold]: v(Set{}) = true
⟨proof⟩

lemma mtSet-rep-set: [[Rep-Set-0 (Set{})]] = {}
⟨proof⟩

Note that the collection types in OCL allow for null to be included; however, there is the null-collection into which inclusion yields invalid.

4.4. Complex Types: The Set-Collection Type (II) Library

This part provides a collection of operators for the Set type.

4.4.1. Computational Operations on Set

Definition

definition OclIncluding :: (\'A, 'α::null) Set, ('\'A, 'α) val] ⇒ ('\'A, 'α) Set
where OclIncluding x y = (λ τ. if (δ x) τ = true τ ∧ (υ y) τ = true τ
then Abs-Set-0 [[ [[Rep-Set-0 (x τ)]] ∪ {y τ} ]]
else ⊥)

notation OclIncluding (−→ including'(·'))

syntax
-OclFinset :: args => ('\'A, 'a::null) Set (Set{·})
translations
Set{x, xs} ::= CONST OclIncluding (Set{x}) x
Set{x} ::= CONST OclIncluding (Set{}) x

definition OclExcluding :: (\'A, 'α::null) Set, ('\'A, 'α) val] ⇒ ('\'A, 'α) Set
where OclExcluding x y = (λ τ. if (δ x) τ = true τ ∧ (υ y) τ = true τ
then Abs-Set-0 [[ [[Rep-Set-0 (x τ)]] − {y τ} ]]
else ⊥)

notation OclExcluding (−→ excluding'(·'))

definition OclIncludes :: (\'A, 'α::null) Set, ('\'A, 'α) val] ⇒ 'A Boolean
where OclIncludes x y = (λ τ. if (δ x) τ = true τ ∧ (υ y) τ = true τ
then [[[y τ] ∈ [[Rep-Set-0 (x τ)]] ]]
else ⊥)

notation OclIncludes (−→ includes'(·'))

definition OclExcludes :: (\'A, 'α::null) Set, ('\'A, 'α) val] ⇒ 'A Boolean
where OclExcludes x y = (not(OclIncludes x y))
The case of the size definition is somewhat special, we admit explicitly in Featherweight OCL the possibility of infinite sets. For the size definition, this requires an extra condition that assures that the cardinality of the set is actually a defined integer.

**Definition**

\[
\text{OclSize} :: (\alpha, \alpha::null) \text{Set} \Rightarrow \alpha \text{ Integer}
\]

where

\[
\text{OclSize} x = (\lambda \tau. \text{if } (\delta x \tau = \text{true} \land \text{finite}([[\text{Rep-Set-0} (x \tau)]]) \text{ then } \lfloor \text{int} (\text{card} [[\text{Rep-Set-0} (x \tau)]) \rfloor \text{ else } \bot)
\]

**Notation**

\[
\text{OclSize} (-\rightarrow \text{size}('))
\]

The following definition follows the requirement of the standard to treat null as neutral element of sets. It is a well-documented exception from the general strictness rule and the rule that the distinguished argument self should be non-null.

**Definition**

\[
\text{OclIsEmpty} :: (\alpha, \alpha::null) \text{Set} \Rightarrow \alpha \text{ Boolean}
\]

where

\[
\text{OclIsEmpty} x = (\upsilon x \text{ and } \neg (\delta x) \text{ or } (\text{OclSize} x = 0))
\]

**Notation**

\[
\text{OclIsEmpty} (-\rightarrow \text{isEmpty}('))
\]

**Definition**

\[
\text{OclNotEmpty} :: (\alpha, \alpha::null) \text{Set} \Rightarrow \alpha \text{ Boolean}
\]

where

\[
\text{OclNotEmpty} x = \neg (\text{OclIsEmpty} x)
\]

**Notation**

\[
\text{OclNotEmpty} (-\rightarrow \text{notEmpty}('))
\]

The definition of OclForall mimics the one of op and: OclForall is not a strict operation.

**Definition**

\[
\text{OclForall} :: [(\alpha, \alpha::null) \text{Set}, (\alpha, \alpha) \text{val}] \Rightarrow \alpha \text{ Boolean}
\]

where

\[
\text{OclForall} S P = (\lambda \tau. \text{if } (\delta S \tau = \text{true} \text{ then } \text{if } (\exists x \in [[\text{Rep-Set-0} (S \tau)]]. P(\lambda - x \tau = \text{false}) \text{ then false } \tau \text{ else if } (\exists x \in [[\text{Rep-Set-0} (S \tau)]]. P(\lambda - x \tau = x \tau) \text{ then } \bot \tau \text{ else if } (\exists x \in [[\text{Rep-Set-0} (S \tau)]]. P(\lambda - x \tau = \text{null} \tau) \text{ then null } \tau \text{ else true } \tau \text{ else } \bot)
\]

**Syntax**

\[
\text{-OclForall :: [(\alpha, \alpha::null) \text{Set}, id, (\alpha)\text{Boolean}] \Rightarrow (\alpha) \text{ Boolean} \quad ((\cdot)\rightarrow \text{forAll}'(\cdot')})
\]

**Translations**

\[
X \rightarrow \text{forAll}(x \mid P) == \text{CONST} \; \text{OclForall} \; X \; (\%x. \; P)
\]

Like OclForall, OclExists is also not strict.
definition **OclExists** :: \((\exists \, x \colon \alpha \text{::null}) \text{Set} . (\exists \, \alpha \text{val}\Rightarrow(\exists \, ) \text{Boolean}) \Rightarrow \alpha \text{ Boolean}\)

where **OclExists** \(S \, P \equiv \text{not}(\text{OclForall} \, S \, (\lambda \, x . \text{not} \, (P \, X)))\)

**syntax**

- **OclExist** :: \([(\exists \, \alpha \text{::null}) \text{Set} . \text{id} . (\exists \, \alpha \text{Boolean}) \Rightarrow \alpha \text{ Boolean} \quad (\lambda \to \exists \text{'}|\text{'}\text{')}\]

**translations**

\(X \to \exists \text{x}(x \mid P) \equiv \text{CONST} \, \text{OclExists} \, X \, (\%x. \, P)\)

**definition** **OclIterate** :: \([(\exists \, \alpha \text{::null}) \text{Set} . (\exists \, \beta \text{::null}) \text{val}, \quad (\exists \, \alpha \text{val}\Rightarrow(\exists \, \beta \text{val}\Rightarrow(\exists \, \beta \text{val} \Rightarrow (\exists \, \beta) \text{val})) \Rightarrow (\exists \, \beta) \text{val}}\]

where **OclIterate** \(S \, A \, F = (\lambda \, \tau . \text{if} \, (\delta \, S) \, \tau = \text{true} \wedge (\nu \, A) \, \tau = \text{true} \wedge \text{finite}[\text{Rep-Set-0} \, (S \, \tau)]\]

\(\text{then} \, (\text{Finite-Set.fold} \, (F) \, (A) \, ((\lambda \, \alpha . \, \alpha) \cdot [\text{Rep-Set-0} \, (S \, \tau)]))) \, \tau \quad \text{else} \, \bot\)

**syntax**

- **OclIterate** :: \([(\exists \, \alpha \text{::null}) \text{Set} . \text{id} . \text{id} . \alpha, \, \beta] \Rightarrow (\exists \, \gamma) \text{val} \quad (\lambda \to \text{iterate}(\cdot;\,\text{'}|\text{'}\text{')}\)

**translations**

\(X \to \text{iterate}(a; \, x \mid A \mid P) \equiv \text{CONST} \, \text{OclIterate} \, X \, A \, (\%a. \, (\%x. \, P))\)

**definition** **OclSelect** :: \([(\exists \, \alpha \text{::null}) \text{Set} . (\exists \, \alpha) \text{val}\Rightarrow(\exists \, \alpha) \text{Boolean}] \Rightarrow (\exists \, \alpha) \text{Set}\)

where **OclSelect** \(S \, P = (\lambda \, \tau . \text{if} \, (\delta \, S) \, \tau = \text{true} \, \tau \\text{then} \, \text{if} \, (\exists \, x \in [\text{Rep-Set-0} \, (S \, \tau)] \ldots \text{else} \, \text{Abs-Set-0} \, [\{x \in [\text{Rep-Set-0} \, (S \, \tau)] \ldots \text{P} \, (\lambda . \, x) \, \tau \neq \text{false} \, \tau)] \ldots \text{else} \, \bot}\)

**syntax**

- **OclSelect** :: \([(\exists \, \alpha \text{::null}) \text{Set} . \text{id} . (\exists \, \alpha) \text{Boolean}] \Rightarrow \alpha \text{ Boolean} \quad (\lambda \to \text{select}(\cdot;\,\text{'}|\text{'}\text{')}\)

**translations**

\(X \to \text{select}(x \mid P) \equiv \text{CONST} \, \text{OclSelect} \, X \, (\%x. \, P)\)

**definition** **OclReject** :: \([(\exists \, \alpha \text{::null}) \text{Set} . (\exists \, \alpha) \text{val}\Rightarrow(\exists \, \alpha) \text{Boolean}] \Rightarrow (\exists \, \alpha \text{::null}) \text{Set}\)

where **OclReject** \(S \, P = \text{OclSelect} \, S \, (\lambda o \, P)\)

**syntax**

- **OclReject** :: \([(\exists \, \alpha \text{::null}) \text{Set} . (\exists \, \alpha) \text{Boolean}] \Rightarrow \alpha \text{ Boolean} \quad (\lambda \to \text{reject}(\cdot;\,\text{'}|\text{'}\text{')}\)

**translations**

\(X \to \text{reject}(x \mid P) \equiv \text{CONST} \, \text{OclReject} \, X \, (\%x. \, P)\)

**Definition (futur operators)**

**consts**

\(\text{OclCount} :: [(\exists \, \alpha) \text{::null}) \text{Set} . (\exists \, \alpha) \text{Set}] \Rightarrow \alpha \text{ Integer}\)

\(\text{OclSum} :: (\exists \, \alpha \text{::null}) \text{Set} \Rightarrow \alpha \text{ Integer}\)

\(\text{OclIncludesAll} :: [(\exists \, \alpha \text{::null}) \text{Set} . (\exists \, \alpha) \text{Set}] \Rightarrow \alpha \text{ Boolean}\)

\(\text{OclExcludesAll} :: [(\exists \, \alpha \text{::null}) \text{Set} . (\exists \, \alpha) \text{Set}] \Rightarrow \alpha \text{ Boolean}\)

\(\text{OclComplement} :: (\exists \, \alpha \text{::null}) \text{Set} \Rightarrow (\exists \, \alpha) \text{Set}\)

\(\text{OclUnion} :: [(\exists \, \alpha \text{::null}) \text{Set} . (\exists \, \alpha) \text{Set}] \Rightarrow (\exists \, \alpha) \text{Set}\)

\(\text{OclIntersection} :: [(\exists \, \alpha \text{::null}) \text{Set} . (\exists \, \alpha) \text{Set}] \Rightarrow (\exists \, \alpha) \text{Set}\)

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notation
  OclCount (-\rightarrow count'(\cdot))
notation
  OclSum (-\rightarrow\sum'(\cdot))
notation
  OclIncludesAll (-\rightarrow includesAll'(\cdot))
notation
  OclExcludesAll (-\rightarrow excludesAll'(\cdot))
notation
  OclComplement (-\rightarrow complement'(\cdot))
notation
  OclUnion (-\rightarrow union'(\cdot))
notation
  OclIntersection(-\rightarrow intersection'(\cdot))

4.4.2. Validity and Definedness Properties

OclIncluding

lemma OclIncluding-defined-args-valid:
(\tau \models \delta(X \rightarrow including(x))) = ((\tau \models (\delta X)) \land (\tau \models (v \ x)))
⟨proof⟩

lemma OclIncluding-valid-args-valid:
(\tau \models v(X \rightarrow including(x))) = ((\tau \models (\delta X)) \land (\tau \models (v \ x)))
⟨proof⟩

lemma OclIncluding-defined-args-valid'[simp,code-unfold]:
\delta(X \rightarrow including(x)) = ((\delta X) \land (v \ x))
⟨proof⟩

lemma OclIncluding-valid-args-valid'[simp,code-unfold]:
v(X \rightarrow including(x)) = ((\delta X) \land (v \ x))
⟨proof⟩

OclExcluding

lemma OclExcluding-defined-args-valid:
(\tau \models \delta(X \rightarrow excluding(x))) = ((\tau \models (\delta X)) \land (\tau \models (v \ x)))
⟨proof⟩

lemma OclExcluding-valid-args-valid:
(\tau \models v(X \rightarrow excluding(x))) = ((\tau \models (\delta X)) \land (\tau \models (v \ x)))
⟨proof⟩

lemma OclExcluding-valid-args-valid'[simp,code-unfold]:
\(\delta(X \rightarrow excluding(x)) = ((\delta X) \text{ and } (v \ x))\)

\[\text{proof}\]

\textbf{lemma OclExcluding-valid-args-valid''\[simp,code-unfold\]:}
\(v(X \rightarrow excluding(x)) = ((\delta X) \text{ and } (v \ x))\)

\[\text{proof}\]

\textbf{OclIncludes}

\textbf{lemma OclIncludes-defined-args-valid:}
\((\tau \models \delta(X \rightarrow includes(x))) = ((\tau \models (\delta X)) \land (\tau \models (v \ x)))\)

\[\text{proof}\]

\textbf{lemma OclIncludes-valid-args-valid\[simp,code-unfold\]:}
\(\delta(X \rightarrow includes(x)) = ((\delta X) \text{ and } (v \ x))\)

\[\text{proof}\]

\textbf{lemma OclIncludes-valid-args-valid''\[simp,code-unfold\]:}
\(v(X \rightarrow includes(x)) = ((\delta X) \text{ and } (v \ x))\)

\[\text{proof}\]

\textbf{OclExcludes}

\textbf{lemma OclExcludes-defined-args-valid:}
\((\tau \models \delta(X \rightarrow excludes(x))) = ((\tau \models (\delta X)) \land (\tau \models (v \ x)))\)

\[\text{proof}\]

\textbf{lemma OclExcludes-valid-args-valid\[simp,code-unfold\]:}
\(\delta(X \rightarrow excludes(x)) = ((\delta X) \text{ and } (v \ x))\)

\[\text{proof}\]

\textbf{lemma OclExcludes-valid-args-valid''\[simp,code-unfold\]:}
\(v(X \rightarrow excludes(x)) = ((\delta X) \text{ and } (v \ x))\)

\[\text{proof}\]

\textbf{OclSize}

\textbf{lemma OclSize-defined-args-valid: }\tau \models \delta (X \rightarrow size()) \implies \tau \models \delta X \\}

\[\text{proof}\]

\textbf{lemma OclSize-infinite:}

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assumes non-finite: \( \tau \models \text{not}(\delta(S \rightarrow \text{size}())) \)
shows (\( \tau \models \text{not}(\delta(S)) \) \lor \neg \text{finite } [[\text{Rep-Set-0} (S \ \tau)]] \)
{proof}

\text{lemma } \tau \models \delta X \implies \neg \text{finite } [[\text{Rep-Set-0} (X \ \tau)]] \implies \neg \tau \models \delta (X \rightarrow \text{size}())
{proof}

\text{lemma } \text{size-defined}: 
assumes X-finite: \( \forall \tau. \text{finite } [[\text{Rep-Set-0} (X \ \tau)]] \)
shows \( \delta (X \rightarrow \text{size}()) = \delta X \)
{proof}

\text{lemma } \text{size-defined}': 
assumes X-finite: \( \text{finite } [[\text{Rep-Set-0} (X \ \tau)]] \)
shows \( \tau \models \delta (X \rightarrow \text{size}()) = (\tau \models \delta X) \)
{proof}

\text{OclIsEmpty}

\text{lemma } \text{OclIsEmpty-defined-args-valid}: \tau \models \delta (X \rightarrow \text{isEmpty}()) \implies \tau \models v X
{proof}

\text{lemma } \tau \models \delta (\text{null} \rightarrow \text{isEmpty}())
{proof}

\text{lemma } \text{OclIsEmpty-infinite}: \tau \models \delta X \implies \neg \text{finite } [[\text{Rep-Set-0} (X \ \tau)]] \implies \neg \tau \models \delta (X \rightarrow \text{isEmpty}())
{proof}

\text{OclNotEmpty}

\text{lemma } \text{OclNotEmpty-defined-args-valid}: \tau \models \delta (X \rightarrow \text{notEmpty}()) \implies \tau \models v X
{proof}

\text{lemma } \tau \models \delta (\text{null} \rightarrow \text{notEmpty}())
{proof}

\text{lemma } \text{OclNotEmpty-infinite}: \tau \models \delta X \implies \neg \text{finite } [[\text{Rep-Set-0} (X \ \tau)]] \implies \neg \tau \models \delta (X \rightarrow \text{notEmpty}())
{proof}

\text{lemma } \text{OclNotEmpty-has-elt}: \tau \models \delta X \implies \\
\tau \models X \rightarrow \text{notEmpty}() \implies \\
\exists e. e \in [[\text{Rep-Set-0} (X \ \tau)]]
{proof}

\text{OclANY}

\text{lemma } \text{OclANY-defined-args-valid}: \tau \models \delta (X \rightarrow \text{any}()) \implies \tau \models \delta X
{proof}
lemma \( \tau \models \delta \; X \implies \tau \models \chi \implies \neg \tau \models \delta \; (X \implies \text{any}) \)

\( \langle \text{proof} \rangle \)

lemma OclANY-valid-args-valid:
\( (\tau \models \chi (X \implies \text{any})) \equiv (\tau \models \chi \; X) \)

\( \langle \text{proof} \rangle \)

lemma OclANY-valid-args-valid''[simp, code-unfold]:
\( \chi (X \implies \text{any}) = (\chi \; X) \)

\( \langle \text{proof} \rangle \)

4.4.3. Execution with Invalid or Null or Infinite Set as Argument

OclIncluding

lemma OclIncluding-invalid[simp, code-unfold]: invalid \( \implies \) including(x) = invalid

\( \langle \text{proof} \rangle \)

lemma OclIncluding-invalid-args[simp, code-unfold]: X \( \implies \) including(invalid) = invalid

\( \langle \text{proof} \rangle \)

lemma OclIncluding-null[simp, code-unfold]: null \( \implies \) including(x) = invalid

\( \langle \text{proof} \rangle \)

OclExcluding

lemma OclExcluding-invalid[simp, code-unfold]: invalid \( \implies \) excludes(x) = invalid

\( \langle \text{proof} \rangle \)

lemma OclExcluding-invalid-args[simp, code-unfold]: X \( \implies \) excludes(invalid) = invalid

\( \langle \text{proof} \rangle \)

lemma OclExcluding-null[simp, code-unfold]: null \( \implies \) excludes(x) = invalid

\( \langle \text{proof} \rangle \)

OclIncludes

lemma OclIncludes-invalid[simp, code-unfold]: invalid \( \implies \) includes(x) = invalid

\( \langle \text{proof} \rangle \)

lemma OclIncludes-invalid-args[simp, code-unfold]: X \( \implies \) includes(invalid) = invalid

\( \langle \text{proof} \rangle \)

lemma OclIncludes-null[simp, code-unfold]: null \( \implies \) includes(x) = invalid

\( \langle \text{proof} \rangle \)

OclExcludes

lemma OclExcludes-invalid[simp, code-unfold]: invalid \( \implies \) excludes(x) = invalid
lemma OclExcludes-invalid-args[simp,code-unfold]: \((X \rightarrow \text{excludes}(\text{invalid})) = \text{invalid}\)

lemma OclExcludes-null[simp,code-unfold]: \((\text{null} \rightarrow \text{excludes}(x)) = \text{invalid}\)

OclSize
lemma OclSize-invalid[simp,code-unfold]: \((\text{invalid} \rightarrow \text{size}()) = \text{invalid}\)

lemma OclSize-null[simp,code-unfold]: \((\text{null} \rightarrow \text{size}()) = \text{invalid}\)

OclIsEmpty
lemma OclIsEmpty-invalid[simp,code-unfold]: \((\text{invalid} \rightarrow \text{isEmpty}()) = \text{invalid}\)

lemma OclIsEmpty-null[simp,code-unfold]: \((\text{null} \rightarrow \text{isEmpty}()) = \text{true}\)

OclNotEmpty
lemma OclNotEmpty-invalid[simp,code-unfold]: \((\text{invalid} \rightarrow \text{notEmpty}()) = \text{invalid}\)

lemma OclNotEmpty-null[simp,code-unfold]: \((\text{null} \rightarrow \text{notEmpty}()) = \text{false}\)

OclANY
lemma OclANY-invalid[simp,code-unfold]: \((\text{invalid} \rightarrow \text{any}()) = \text{invalid}\)

lemma OclANY-null[simp,code-unfold]: \((\text{null} \rightarrow \text{any}()) = \text{null}\)

OclForall
lemma OclForall-invalid[simp,code-unfold]: \((\text{invalid} \rightarrow \text{forAll}(a \mid P a) = \text{invalid}\)

lemma OclForall-null[simp,code-unfold]: \((\text{null} \rightarrow \text{forAll}(a \mid P a) = \text{invalid}\)

OclExists
lemma OclExists-invalid[simp,code-unfold]: \((\text{invalid} \rightarrow \text{exists}(a \mid P a) = \text{invalid}\)

lemma OclExists-null[simp,code-unfold]: \((\text{null} \rightarrow \text{exists}(a \mid P a) = \text{invalid}\)

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\langle\textit{proof}\rangle

\textbf{lemma} \textit{OclExists-null}:[\textit{simp,code-unfold}]:null\rightarrow exists(a \mid P a) = invalid \\
\langle\textit{proof}\rangle

\textbf{OclIterate}

\textbf{lemma} \textit{OclIterate-invalid}:[\textit{simp,code-unfold}]:invalid\rightarrow iterate(a; x = A \mid P a x) = invalid \\
\langle\textit{proof}\rangle

\textbf{lemma} \textit{OclIterate-null}:[\textit{simp,code-unfold}]:null\rightarrow iterate(a; x = A \mid P a x) = invalid \\
\langle\textit{proof}\rangle

\textbf{lemma} \textit{OclIterate-invalid-args}:[\textit{simp,code-unfold}]:S\rightarrow iterate(a; x = invalid \mid P a x) = invalid \\
\langle\textit{proof}\rangle

\textbf{lemma} \textit{OclIterate-infinite}:
\textbf{assumes} non-finite: \tau \models not(\delta(S\rightarrow size(})) \\
\textbf{shows} (OclIterate S A F) \tau = invalid \tau \\
\langle\textit{proof}\rangle

\textbf{OclSelect}

\textbf{lemma} \textit{OclSelect-invalid}:[\textit{simp,code-unfold}]:invalid\rightarrow select(a \mid P a) = invalid \\
\langle\textit{proof}\rangle

\textbf{lemma} \textit{OclSelect-null}:[\textit{simp,code-unfold}]:null\rightarrow select(a \mid P a) = invalid \\
\langle\textit{proof}\rangle

\textbf{OclReject}

\textbf{lemma} \textit{OclReject-invalid}:[\textit{simp,code-unfold}]:invalid\rightarrow reject(a \mid P a) = invalid \\
\langle\textit{proof}\rangle

\textbf{lemma} \textit{OclReject-null}:[\textit{simp,code-unfold}]:null\rightarrow reject(a \mid P a) = invalid \\
\langle\textit{proof}\rangle

\textbf{4.4.4. Context Passing}

\textbf{lemma} \textit{cp-OclIncluding}:
\lambda \rightarrow incluring(x)) \tau = ((\lambda \rightarrow X \tau)\rightarrow incluring(\lambda \rightarrow x \tau)) \tau \\
\langle\textit{proof}\rangle

\textbf{lemma} \textit{cp-OclExcluding}:
lemma cp-OclIncludes:
\((X \rightarrow including(x)) \tau = ((\lambda \cdot X \tau) \rightarrow including(\lambda \cdot x \tau)) \tau\) 
(proof)

lemma cp-OclIncludes1:
\((X \rightarrow includes(x)) \tau = (X \rightarrow includes(\lambda \cdot x \tau)) \tau\) 
(proof)

lemma cp-OclExcludes:
\((X \rightarrow excludes(x)) \tau = ((\lambda \cdot X \tau) \rightarrow excludes(\lambda \cdot x \tau)) \tau\) 
(proof)

lemma cp-OclSize: \(X \rightarrow \text{size}()) \tau = ((\lambda \cdot X \tau) \rightarrow \text{size}()) \tau\) 
(proof)

lemma cp-OclIsEmpty: \(X \rightarrow \text{isEmpty}() \tau = ((\lambda \cdot X \tau) \rightarrow \text{isEmpty}()) \tau\) 
(proof)

lemma cp-OclNotEmpty: \(X \rightarrow \text{notEmpty}() \tau = ((\lambda \cdot X \tau) \rightarrow \text{notEmpty}()) \tau\) 
(proof)

lemma cp-OclANY: \(X \rightarrow \text{any}() \tau = ((\lambda \cdot X \tau) \rightarrow \text{any}()) \tau\) 
(proof)

lemma cp-OclForall:
\((S \rightarrow \text{forAll}(x \mid P x)) \tau = ((\lambda \cdot S \tau) \rightarrow \text{forAll}(x \mid P (\lambda \cdot x \tau))) \tau\) 
(proof)

lemma cp-OclForall1 [simp,introl!]:
\(cp S \Rightarrow cp (\lambda X. ((S X) \rightarrow \text{forAll}(x \mid P x)))\) 
(proof)

lemma cp-OclExists:
\(cp (\lambda X St x. P (\lambda \tau. x) X St) \Rightarrow cp S \Rightarrow cp (\lambda X. (S X) \rightarrow \text{forAll}(x \mid P x X))\) 
(proof)

lemma cp-OclExists:
\((\forall x. cp(P x)) \Rightarrow cp(\lambda X. ((S X) \rightarrow \text{forAll}(x \mid P x X)))\) 
(proof)
\( (S \rightarrow \exists x \mid P x) \tau = ((\lambda x. S \tau) \rightarrow \exists x \mid P (\lambda x. x)) \tau \)

\[ \langle \text{proof} \rangle \]

**Lemma cp-OclExists1** [simp.intros!]:

\( cp S \Rightarrow cp (\lambda X. ((S X) \rightarrow \exists x \mid P x)) \)

\[ \langle \text{proof} \rangle \]

**Lemma cp-OclIterate**: \( (X \rightarrow iterate(a; x = A \mid P a x)) \tau = ((\lambda X \tau) \rightarrow iterate(a; x = A \mid P a x)) \tau \)

\[ \langle \text{proof} \rangle \]

**Lemma cp-OclSelect**: \( (X \rightarrow select(a \mid P a)) \tau = ((\lambda X \tau) \rightarrow select(a \mid P a)) \tau \)

\[ \langle \text{proof} \rangle \]

**Lemma cp-OclReject**: \( (X \rightarrow reject(a \mid P a)) \tau = ((\lambda X \tau) \rightarrow reject(a \mid P a)) \tau \)

\[ \langle \text{proof} \rangle \]

**Lemmas cp-intro''[intro!, simp.code-unfold] =**

\( cp-intro' \)

\( cp-OclIncluding \ [\text{THEN all1[THEN all1[THEN all1[THEN cp12]]], of OclIncluding}] \)

\( cp-OclExcluding \ [\text{THEN all1[THEN all1[THEN all1[THEN cp12]]], of OclExcluding}] \)

\( cp-OclIncludes \ [\text{THEN all1[THEN all1[THEN all1[THEN cp12]]], of OclIncludes}] \)

\( cp-OclExcludes \ [\text{THEN all1[THEN all1[THEN all1[THEN cp12]]], of OclExcludes}] \)

\( cp-OclSize \ [\text{THEN all1[THEN all1[THEN cp11], of OclSize}] \)

\( cp-OclIsEmpty \ [\text{THEN all1[THEN all1[THEN cp11], of OclIsEmpty}] \)

\( cp-OclNotEmpty \ [\text{THEN all1[THEN all1[THEN cp11], of OclNotEmpty}] \)

\( cp-OclANY \ [\text{THEN all1[THEN all1[THEN cp11], of OclANY}] \)

4.4.5. Const

**Lemma const-OclIncluding[simp, code-unfold]** :

**Assumes** const-x : const x

**And** const-S : const S

**Shows** const \( (S \rightarrow including(x)) \)

\[ \langle \text{proof} \rangle \]

4.5. Fundamental Predicates on Set: Strict Equality

4.5.1. Definition

After the part of foundational operations on sets, we detail here equality on sets. Strong equality is inherited from the OCL core, but we have to consider the case of the strict equality. We decide to overload strict equality in the same way we do for other value’s in OCL:

**Def** StrictRefEqSet :
\[(x::(\mathcal{A},'\alpha::null)\text{Set}) \doteq y \equiv \lambda \tau. \begin{cases} \text{if } (v \ x) \ \tau = \text{true} \land (v \ y) \ \tau = \text{true} \\ \text{then } (x \triangleq y) \ \tau \\ \text{else invalid } \tau \end{cases} \]

One might object here that for the case of objects, this is an empty definition. The answer is no, we will restrain later on states and objects such that any object has its oid stored inside the object (so the ref, under which an object can be referenced in the store will represented in the object itself). For such well-formed stores that satisfy this invariant (the WFF-invariant), the referential equality and the strong equality—and therefore the strict equality on sets in the sense above—coincides.

4.5.2. Logic and Algebraic Layer on Set

Reflexivity

To become operational, we derive:

**lemma** StrictRefEqSet-refl[simp,code-unfold]:
\[(x::(\mathcal{A},'\alpha::null)\text{Set}) \doteq x) = (if (v \ x) \ \text{then } \text{true else } \text{invalid endif})\]

**proof**

Symmetry

**lemma** StrictRefEqSet-sym:
\[(x::(\mathcal{A},'\alpha::null)\text{Set}) \doteq y) = (y \doteq x)\]

**proof**

Execution with Invalid or Null as Argument

**lemma** StrictRefEqSet-strict1[simp,code-unfold]:
\[(x::(\mathcal{A},'\alpha::null)\text{Set}) \doteq \text{invalid})= \text{invalid}\]

**proof**

**lemma** StrictRefEqSet-strict2[simp,code-unfold]:
\[(\text{invalid} \doteq (y::(\mathcal{A},'\alpha::null)\text{Set}))= \text{invalid}\]

**proof**

**lemma** StrictRefEqSet-strictEq-valid-args-valid:
\[(\tau \models \delta ((x::(\mathcal{A},'\alpha::null)\text{Set}) \doteq y)) = ((\tau \models (v \ x)) \land (\tau \models v \ y))\]

**proof**

Behavior vs StrongEq

**lemma** StrictRefEqSet-vs-StrongEq:
\[\tau \models v \ x \Rightarrow \tau \models v \ y \Rightarrow (\tau \models ((x::(\mathcal{A},'\alpha::null)\text{Set}) \doteq y) \triangleq (x \triangleq y)))\]

**proof**

Context Passing

**lemma** cp-StrictRefEqSet:(X::(\mathcal{A},'\alpha::null)\text{Set}) \doteq Y) \ \tau = ((\lambda \cdot \ X \ \tau) \doteq (\lambda \cdot \ Y \ \tau)) \ \tau\]

**proof**
 Const

lemma const-StrictRefEqSet :
  assumes const (X :: (-,::null) Set)
  assumes const X'
  shows const (X = X')
⟨proof⟩

4.6. Execution on Set’s Operators (with mtSet and recursive case as arguments)

4.6.1. OclIncluding

lemma OclIncluding-finite-rep-set :
  assumes X-def : τ |= δ X
  and x-val : τ |= v x
  shows finite ![Rep-Set-0 (X ->including(x) τ)] = finite ![Rep-Set-0 (X τ)]
⟨proof⟩

lemma OclIncluding-rep-set :
  assumes S-def : τ |= δ S
  shows ![Rep-Set-0 (S ->including(λ- [[x]]) τ)] = insert [[x]] ![Rep-Set-0 (S τ)]
⟨proof⟩

lemma OclIncluding-notempty-rep-set :
  assumes X-def : τ |= δ X
  and a-val : τ |= v a
  shows ![Rep-Set-0 (X ->including(a) τ)] ≠ {}
⟨proof⟩

lemma OclIncluding-includes :
  assumes τ |= X ->includes(x)
  shows X ->including(x) τ = X τ
⟨proof⟩

4.6.2. OclExcluding

lemma OclExcluding-charn0[simp]:
  assumes val-x:τ |= (v x)
  shows τ |= ((Set{}) ->excluding(x)) ≡ Set{})
⟨proof⟩

lemma OclExcluding-charn0-exec[simp,code-unfold]:
  (Set{}) ->excluding(x)) = (if (v x) then Set{} else invalid endif)
⟨proof⟩

lemma OclExcluding-charn1 :
  assumes def-X:τ |= (δ X)

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and \( \text{val-x}:\tau \models (v \, x) \)
and \( \text{val-y}:\tau \models (v \, y) \)
and \( \text{neq} \cdot \tau \models \neg (x \triangleq y) \)
s-shows \( \tau \models ((X\rightarrow \text{including}(x))\rightarrow \text{excluding}(y)) \triangleq ((X\rightarrow \text{excluding}(y))\rightarrow \text{including}(x)) \)

\langle \text{proof} \rangle

\text{lemma OclExcluding-charn2:}
\text{assumes def-X}:\tau \models (\delta \, X) 
and \( \text{val-x}:\tau \models (v \, x) \)
\text{shows} \( \tau \models (((X\rightarrow \text{including}(x))\rightarrow \text{excluding}(x)) \triangleq (X\rightarrow \text{excluding}(x))) \)
\langle \text{proof} \rangle

One would like a generic theorem of the form:

\text{lemma OclExcluding-charn-exec:}
"(X\rightarrow \text{including}(x::(:\mathcal{A},'a::null)\text{val})\rightarrow \text{excluding}(y)) = 
(\text{if } \delta \, X \text{ then if } x \triangleq y 
then X\rightarrow \text{excluding}(y) 
else X\rightarrow \text{excluding}(y)\rightarrow \text{including}(x) 
endif 
else \text{ invalid endif})"

Unfortunately, this does not hold in general, since referential equality is an overloaded concept and has to be defined for each type individually. Consequently, it is only valid for concrete type instances for Boolean, Integer, and Sets thereof...

The computational law \textit{OclExcluding-charn-exec} becomes generic since it uses strict equality which in itself is generic. It is possible to prove the following generic theorem and instantiate it later (using properties that link the polymorphic logical strong equality with the concrete instance of strict quality).

\text{lemma OclExcluding-charn-exec:}
\text{assumes strict1}: (x \triangleq \text{invalid}) = \text{invalid} 
and \text{strict2}: (\text{invalid} \triangleq y) = \text{invalid} 
and \text{StrictRefEq-valid-args-valid}: \lambda (x::(\mathcal{A},'a::null)\text{val}) \, y \, \tau. 
(\tau \models \delta \, (x \triangleq y) = ((\tau \models (v \, x)) \land (\tau \models v \, y)) 
and \text{cp-StrictRefEq}: \lambda (X::(\mathcal{A},'a::null)\text{val}) \, Y \, \tau. \, (X \triangleq Y) \, \tau = ((\lambda\. \, X \, \tau) \triangleq (\lambda\. \, Y \, \tau)) \, \tau 
and \text{StrictRefEq-vs-StrongEq}: \lambda (x::(\mathcal{A},'a::null)\text{val}) \, \tau. 
\tau \models (x \triangleq y) \Rightarrow (\tau \models (x \triangleq y) \triangleq (x \triangleq y)) 
\text{shows} \, (X\rightarrow \text{including}(x::(\mathcal{A},'a::null)\text{val})\rightarrow \text{excluding}(y)) = 
(\text{if } \delta \, X \text{ then if } x \triangleq y 
then X\rightarrow \text{excluding}(y) 
else X\rightarrow \text{excluding}(y)\rightarrow \text{including}(x) 
endif 
else \text{ invalid endif}) 
\langle \text{proof} \rangle

\text{schematic-lemma OclExcluding-charn-exec}\_\text{Integer}[\text{simp,code-unfold}]: ?X 
\langle \text{proof} \rangle

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schematic-lemma OclExcluding-charn-execute_Boolean [simp, code-unfold]: ?X
⟨proof⟩

schematic-lemma OclExcluding-charn-execute_set [simp, code-unfold]: ?X
⟨proof⟩

lemma OclExcluding-finite-rep-set:
  assumes X-def: τ |= δ X
  and x-val: τ |= v x
  shows finite [[Rep-Set-0 (X -> excluding(x)) τ]] = finite [[Rep-Set-0 (X τ)]]
⟨proof⟩

lemma OclExcluding-rep-set:
  assumes S-def: τ |= δ S
  shows [[Rep-Set-0 (S -> excluding(λ. [[x]]) τ)]] = [[Rep-Set-0 (S τ)]] - {[[[x]]]}
⟨proof⟩

4.6.3. OclIncludes

lemma OclIncludes-charn0 [simp]:
  assumes val-x: τ |= (ψ x)
  shows τ |= not (Set {} -> includes(x))
⟨proof⟩

lemma OclIncludes-charn0' [simp, code-unfold]:
  Set {} -> includes(x) = (if v x then false else invalid endif)
⟨proof⟩

lemma OclIncludes-charn1:
  assumes def-X: τ |= (δ X)
  assumes val-x: τ |= (v x)
  shows τ |= (X -> including(x) -> includes(x))
⟨proof⟩

lemma OclIncludes-charn2:
  assumes def-X: τ |= (δ X)
  and val-x: τ |= (v x)
  and val-y: τ |= (v y)
  and neq : τ |= not (x ≡ y)
  shows τ |= (X -> including(x) -> includes(y)) ≡ (X -> includes(y))
⟨proof⟩

  Here is again a generic theorem similar as above.

lemma OclIncludes-execute-generic:

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assumes strict1: (x \equiv invalid) = invalid
and strict2: (invalid \equiv y) = (invalid
and cp-StrictRefEq: \lambda (X::(\mathbf{\alpha},'a::\mathbf{null})val) Y. (X \equiv Y) \tau = ((\lambda-. X \tau) \equiv (\lambda-. Y \tau)) \tau
and StrictRefEq-vs-StrongEq: \lambda (x::(\mathbf{\alpha},'a::\mathbf{null})val) y. 
\tau \models v x \implies \tau \models v y \implies (\tau \models ((x \equiv y) \triangleq (x \triangleq y)))

shows 
(X \rightarrow including(x::(\mathbf{\alpha},'a::\mathbf{null})val)\rightarrow includes(y)) =
(if \delta X then if x \equiv y then true else X \rightarrow includes(y) endif else invalid endif

\langle proof \rangle

schematic-lemma OclIncludes-execute\_Integer[simp,code-unfold]: ?X
\langle proof \rangle

schematic-lemma OclIncludes-execute\_Boolean[simp,code-unfold]: ?X
\langle proof \rangle

schematic-lemma OclIncludes-execute\_Set[simp,code-unfold]: ?X
\langle proof \rangle

lemma OclIncludes-including-generic :
assumes OclIncludes-execute-generic [simp] :
\lambda X x y. 
(X \rightarrow including(x::(\mathbf{\alpha},'a::\mathbf{null})val)\rightarrow includes(y)) =
(if \delta X then if x \equiv y then true else X \rightarrow includes(y) endif else invalid endif
and StrictRefEq-strict'': \lambda x y. \delta ((x::(\mathbf{\alpha},'a::\mathbf{null})val) \equiv y) = (\nu(x) and \nu(y))
and a-val : \tau \models v a
and x-val : \tau \models v x
and S-incl : \tau \models (S)\rightarrow includes((x::(\mathbf{\alpha},'a::\mathbf{null})val))
shows \tau \models S \rightarrow including((a::(\mathbf{\alpha},'a::\mathbf{null})val))\rightarrow includes(x)
\langle proof \rangle

lemmas OclIncludes-including\_Integer =
OclIncludes-including-generic[\OF OclIncludes-execute\_Integer StrictRefEqInteger-strict'']

4.6.4. OclExcludes

4.6.5. OclSize

lemma [simp,code-unfold]: Set{} \rightarrow size() = 0
\langle proof \rangle

lemma OclSize-including-exec[simp,code-unfold]:
((X \rightarrow including(x)) \rightarrow size()) = (if \delta X and \nu x then 
X \rightarrow size() + if X \rightarrow includes(x) then 0 else 1 endif 
else invalid 
endif)

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4.6.6. OclIsEmpty

lemma [simp, code-unfold]: Set{} -> isEmpty() = true

lemma OclIsEmpty-including [simp]:
assumes X-def: \( \tau \models \delta X \)
    and X-finite: finite [Rep-Set-0 (X \( \tau \)])
    and a-val: \( \tau \models \nu a \)
shows X -> including(a) -> isEmpty() \( \tau \) = false \( \tau \)

4.6.7. OclNotEmpty

lemma [simp, code-unfold]: Set{} -> notEmpty() = false

lemma OclNotEmpty-including [simp, code-unfold]:
assumes X-def: \( \tau \models \delta X \)
    and X-finite: finite [Rep-Set-0 (X \( \tau \)])
    and a-val: \( \tau \models \nu a \)
shows X -> including(a) -> notEmpty() \( \tau \) = true \( \tau \)

4.6.8. OclANY

lemma [simp, code-unfold]: Set{} -> any() = null

lemma OclANY-singleton-exec [simp, code-unfold]:
(Set{} -> including(a)) -> any() = a

4.6.9. OclForall

lemma OclForall-mtSet-exec [simp, code-unfold]:
((Set{}) ->forall(z | P(z))) = true

lemma OclForall-including-exec [simp, code-unfold]:
assumes cp\( \theta \) : cp P
shows ((S -> including(x)) ->forall(z | P(z))) = (if \( \delta S \) and \( \nu x \) then \( P x \) and \( S ->forall(z | P(z)) \)
else invalid)

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4.6.10. OclExists

lemma OclExists-mtSet-exec[simp,code-unfold] :

\((\{\}\rightarrow\exists z. P(z)) = \text{false}\)

(proof)

lemma OclExists-including-exec[simp,code-unfold] :

assumes cp: cp P

shows (\((S\rightarrow\text{including}(x))\rightarrow\exists z. P(z)) = (\text{if } \delta S \text{ and } v x \text{ then } P x \text{ or } (S\rightarrow\exists z. P(z)) \text{ else invalid } \text{ endif})

(proof)

4.6.11. OclIterate

lemma OclIterate-empty[simp,code-unfold] : ((\{\}\rightarrow\text{iterate}(a; x = A \mid P a x)) = A

(proof)

In particular, this does hold for A = null.

lemma OclIterate-including :

assumes S-finite: \(\tau \models \delta(S\rightarrow\text{size}())\)

and \(F\text{-valid-arg}: (v A) \tau = (v (F a A)) \tau\)

and \(F\text{-commute}: \text{comp-fun-commute } F\)

and \(F\text{-cp} : \lambda x y \tau. F x y \tau = F (\lambda - x \tau) y \tau\)

shows (\((S\rightarrow\text{including}(a))\rightarrow\text{iterate}(a; x = A \mid F a x)) \tau = (\((S\rightarrow\text{excluding}(a))\rightarrow\text{iterate}(a; x = F a A \mid F a x)) \tau)

(proof)

4.6.12. OclSelect

lemma OclSelect-mtSet-exec[simp,code-unfold]: OclSelect mtSet P = mtSet

(proof)

definition OclSelect-body :: - \Rightarrow - \Rightarrow - \Rightarrow (\forall a, 'a option option) Set

\equiv (\lambda P x acc. \text{if } P x \neq \text{false then acc else acc}\rightarrow\text{including}(x) \text{ endif})

lemma OclSelect-including-exec[simp,code-unfold]:

assumes P-cp : cp P

shows OclSelect (X\rightarrow\text{including}(y)) P = OclSelect-body P y (OclSelect (X\rightarrow\text{excluding}(y)) P)

(is - = ?select)

(proof)

4.6.13. OclReject

lemma OclReject-mtSet-exec[simp,code-unfold]: OclReject mtSet P = mtSet

(proof)

lemma OclReject-including-exec[simp,code-unfold]:
\text{assumes } P\text{-}cp : cp P
\text{shows } Ocl\text{Reject} (X \to \text{including}(y)) P = Ocl\text{Select-body}(\text{not } o P) y (Ocl\text{Reject} (X \to \text{excluding}(y)) P)

\langle proof \rangle

4.7. Execution on Set’s Operators (higher composition)

4.7.1. Ocl\text{Includes}

\text{lemma} Ocl\text{Includes}-\text{any}[\text{simp, code-unfold}]:
\[ X \to \text{includes}(X \to \text{any}()) = (\text{if } \delta X \text{ then}
  \text{if } \delta (X \to \text{size}()) \text{ then not}(X \to \text{isEmpty}())
  \text{ else } X \to \text{includes}(\text{null}) \text{ endif}
  \text{ else invalid endif}) \]

\langle proof \rangle

4.7.2. Ocl\text{Size}

\text{lemma} [\text{simp, code-unfold}]: \delta (\text{Set} \{} \to \text{size}()) = \text{true} \]
\langle proof \rangle

\text{lemma} [\text{simp, code-unfold}]: \delta ((X \to \text{including}(x)) \to \text{size}()) = (\delta (X \to \text{size}()) \text{ and } v(x)) \]
\langle proof \rangle

\text{lemma} [\text{simp}]:
\text{assumes } X\text{-finite: } \forall \tau. \text{finite } [[\text{Rep-Set-0} (X \tau)]]
\text{shows } \delta ((X \to \text{including}(x)) \to \text{size}()) = (\delta (X \to \text{size}()) \text{ and } v(x)) \]
\langle proof \rangle

4.7.3. Ocl\text{Forall}

\text{lemma} Ocl\text{Forall}-\text{rep-set-false}:
\text{assumes } \tau \models \delta X
\text{shows } (Ocl\text{Forall} X P \tau = \text{false } \tau) = (\exists x \in [[\text{Rep-Set-0} (X \tau)]]. P (\lambda \tau. x) \tau = \text{false } \tau) \]
\langle proof \rangle

\text{lemma} Ocl\text{Forall}-\text{rep-set-true}:
\text{assumes } \tau \models \delta X
\text{shows } (\tau \models \forall \text{OclForall} X P) = (\forall x \in [[\text{Rep-Set-0} (X \tau)]]. \tau \models P (\lambda \tau. x)) \]
\langle proof \rangle

\text{lemma} Ocl\text{Forall}-\text{includes}:
\text{assumes } x\text{-def : } \tau \models \delta x
\text{ and } y\text{-def : } \tau \models \delta y
\text{shows } (\tau \models \forall \text{OclForall} x (Ocl\text{Includes} y)) = ([[\text{Rep-Set-0} (x \tau)]]) \subseteq ([[\text{Rep-Set-0} (y \tau)]]])
lemma OclForall-not-includes:
\[\text{assumes } \tau \models \delta x \quad \text{and} \quad \tau \models \delta y \]
\[\text{shows } (\text{OclForall } x \ (\text{OclIncludes } y) \ \tau = \text{false } \tau) = (\neg [[\text{Rep-Set-0} \ (x \ \tau)] \subseteq [[\text{Rep-Set-0} \ (y \ \tau)]])\]
\[\langle \text{proof} \rangle\]

lemma OclForall-iterate:
\[\text{assumes } S \text{-finite: finite } [[\text{Rep-Set-0} \ (S \ \tau)]]\]
\[\text{shows } S \rightarrow \text{forAll} (x \mid P \ x) \ \tau = (S \rightarrow \text{iterate}(x; \ \text{acc = true} \mid \ \text{acc and } P \ x)) \ \tau\]
\[\langle \text{proof} \rangle\]

lemma OclForall-cong:
\[\text{assumes } \forall x. x \in [[\text{Rep-Set-0} \ (X \ \tau)]] \implies \tau \models P (\lambda \tau. x) \implies \tau \models Q (\lambda \tau. x)\]
\[\text{assumes } P: \tau \models \text{OclForall } X\ P\]
\[\text{assumes } Q: \tau \models \text{OclForall } X\ Q\]
\[\text{shows } \tau \models \text{OclForall } X\ R\]
\[\langle \text{proof} \rangle\]

4.7.4. Strict Equality

lemma StrictRefEqSet-defined:
\[\text{assumes } x \text{-def: } \tau \models \delta x\]
\[\text{assumes } y \text{-def: } \tau \models \delta y\]
\[\text{shows } ((x :: (\emptyset, \alpha :: \text{null})\ \text{Set}) = y) \ \tau =\]
\[\quad (x \rightarrow \text{forAll}(x \mid y \rightarrow \text{includes}(z))) \text{ and } (y \rightarrow \text{forAll}(z \mid x \rightarrow \text{includes}(z))) \ \tau\]
\[\langle \text{proof} \rangle\]

lemma StrictRefEqSet-exec[simp, code-unfold]:
\[(x :: (\emptyset, \alpha :: \text{null})\ \text{Set}) = y\]
\[\text{if } \delta x \ \text{then }\]
\[\quad \text{if } \delta y \]
\[\quad \text{then } ((x \rightarrow \text{forAll}(z \mid y \rightarrow \text{includes}(z))) \text{ and } (y \rightarrow \text{forAll}(z \mid x \rightarrow \text{includes}(z))))\]
\[\quad \text{else if } v y \]
\[\quad \text{then } \text{false} (* x' \rightarrow \text{includes} = \text{null} *)\]
\[\quad \text{else invalid}\]
\[\quad \text{endif}\]
\[\quad \text{endif}\]
\[\text{else if } v x (* \text{null} = ??? *)\]
\[\quad \text{then } \text{if } v y \ \text{then not(\delta y) else invalid}\]
\[\text{else invalid}\]

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lemma StrictRefEqSet-L subst1 : cp P \implies \tau \models v x \implies \tau \models v y \implies \tau \models v P x \implies \tau \models v P y \\
\tau \models (x:(\text{A},'a::null) Set) \uplus y \implies \tau \models (P x:(\text{A},'a::null) Set) \uplus P y 

lemma OclIncluding cong' : 
shows \tau \models \delta s \implies \tau \models \delta t \implies \tau \models v x \\
\tau \models (s:(\text{A},'a::null) Set) \uplus t \implies \tau \models (s-> including(x) \uplus (t-> including(x)))

lemma OclIncluding cong : \bigwedge (s:(\text{A},'a::null) Set) t x y. \tau \models \delta t \implies \tau \models v y \\
\tau \models s \uplus t \implies x = y \implies \tau \models s-> including(x) \uplus (t-> including(y))

lemma const StrictRefEqSet including : const a \implies const S \implies const X \implies 
const (X \uplus S-> including(a))

4.8. Test Statements

lemma syntax test : Set\{2,1\} = (Set\{\}->including(1)->including(2))

lemma short-cut' [simp, code-unfold] : (8 \doteq 6) = false
lemma short-cut'' [simp, code-unfold] : (2 \doteq 1) = false
lemma short-cut''' [simp, code-unfold] : (1 \doteq 2) = false

Elementary computations on Sets.

declare OclSelect body-def [simp]

value \neg \tau \models v(invalid::(\text{A},'a::null) Set))
value \tau \models v(null::(\text{A},'a::null) Set)
\[
\begin{align*}
\text{value } \vdash (\tau \models \delta(\text{null}::(\text{\&}, \text{\&}) \text{Set})) \\
\text{value } \vdash (\tau \models v(\text{Set}())) \\
\text{value } \vdash (\tau \models v(\text{Set}()\{2\}, \text{null})) \\
\text{value } \vdash (\tau \models \delta(\text{Set}()\{2\}, \text{null})) \\
\text{value } \vdash (\tau \models (\text{Set}()\{2\} -> \text{includes}(\text{1}))) \\
\text{value } \vdash (\tau \models (\text{Set}()\{2\} -> \text{includes}(\text{null}))) \\
\text{value } \vdash (\tau \models (\text{Set}()\{2\} -> \text{includes}(\text{null}))) \\
\text{value } \vdash (\tau \models (\text{Set}()\{null, 2\} -> \text{includes}(\text{null}))) \\
\text{value } \vdash ((\text{Set}()\{\} -> \text{forall}(z \mid 0 < z)) \\
\text{value } \vdash ((\text{Set}()\{\} -> \text{forall}(z \mid 0 < z)) \\
\text{value } \vdash ((\text{Set}()\{\} -> \text{exists}(z \mid z < 0))) \\
\text{value } \vdash ((\text{Set}()\{\} -> \text{forall}(z \mid 0 < z))) \\
\text{value } \vdash ((\text{Set}()\{\} -> \text{exists}(z \mid 0 < z))) \\
\text{value } \vdash (\tau \models (\text{Set}()\{\text{null}::\text{\'a Boolean} \} \doteq \text{Set}())) \\
\text{value } \vdash (\tau \models (\text{Set}()\{\text{null}::\text{\'a Integer} \} \doteq \text{Set}())) \\
\text{value } (\tau \models (\text{Set}()\{\text{\&} \mid x \} \doteq \text{Set}()\{\text{\&} \mid x \}))) \\
\text{value } (\tau \models (\text{Set}()\{\text{\&} \mid x \} \doteq \text{Set}()\{\text{\&} \mid x \}))) \\
\text{lemma } (\tau \models (\text{Set}()\{\text{true} \} \doteq \text{Set}()\{\text{false} \})) \text{ (proof) } \\
\text{lemma } (\tau \models (\text{Set}()\{\text{true, true} \} \doteq \text{Set}()\{\text{false} \})) \text{ (proof) } \\
\text{lemma } (\tau \models (\text{Set}()\{\text{2} \} \doteq \text{Set}()\{\text{1} \})) \text{ (proof) } \\
\text{lemma } (\tau \models (\text{Set}()\{\text{2}, \text{null, 2} \} \doteq \text{Set}()\{\text{null, 2} \})) \text{ (proof) } \\
\text{lemma } (\tau \models (\text{Set}()\{\text{1}, \text{null, 2} \} \triangleright \text{Set}()\{\text{null, 2} \})) \text{ (proof) } \\
\text{lemma } (\tau \models (\text{Set}()\{\text{null} \} \triangleright \text{select}(x \mid \text{not } x) \doteq \text{Set}()\{\text{null} \})) \text{ (proof) } \\
\text{lemma } (\tau \models (\text{Set}()\{\text{null} \} \triangleright \text{reject}(x \mid \text{not } x) \doteq \text{Set}()\{\text{null} \})) \text{ (proof) } \\
\text{lemma } \text{const } (\text{Set}()\{\text{null, 2} \}, \text{invalid}) \text{ (proof) } \\
\end{align*}
\]
5. Formalization III: State Operations and Objects

theory OCL-state
imports OCL-lib
begin

5.1. Introduction: States over Typed Object Universes

In the following, we will refine the concepts of a user-defined data-model (implied by a class-diagram) as well as the notion of state used in the previous section to much more detail. Surprisingly, even without a concrete notion of objects and a universe of object representation, the generic infrastructure of state-related operations is fairly rich.

5.1.1. Recall: The Generic Structure of States

Recall the foundational concept of an object id (oid), which is just some infinite set.

type-synonym oid = nat

Further, recall that states are pair of a partial map from oid's to elements of an object universe ′a—the heap—and a map to relations of objects. The relations were encoded as lists of pairs to leave the possibility to have Bags, OrderedSets or Sequences as association ends.

This leads to the definitions:

record ′a state =
  heap :: "oid → ′a”
  assocs2 :: "oid → (oid × oid) list ”
  assocs3 :: "oid → (oid × oid × oid) list ”

type-synonym ′a st = ”′a state × ′a state”

Now we refine our state-interface. In certain contexts, we will require that the elements of the object universe have a particular structure; more precisely, we will require that there is a function that reconstructs the oid of an object in the state (we will settle the question how to define this function later).

class object = fixes oid-of :: ′a ⇒ oid
Thus, if needed, we can constrain the object universe to objects by adding the following type class constraint:

\[
\text{typ } \alpha :: \text{object}
\]

The major instance needed are instances constructed over options: once an object, options of objects are also objects.

\[
\text{instantiation } \text{option} :: (\text{object})\text{object}
\]

\[
\begin{align*}
\text{definition } \text{oid-of-option-def}: & \quad \text{oid-of } x = \text{oid-of } (\text{the } x) \\
\text{instance } & \quad (\text{proof})
\end{align*}
\]

5.2. Fundamental Predicates on Object: Strict Equality

Definition

Generic referential equality - to be used for instantiations with concrete object types ...

\[
\begin{align*}
\text{definition } \text{StrictRefEqObject} :: (\alpha, \text{a}::\{\text{object}, \text{null}\})\text{val} & \Rightarrow (\alpha, \text{a})\text{val} \Rightarrow (\alpha)\text{Boolean} \\
\text{where } & \quad \text{StrictRefEqObject} x y \\
& \quad \equiv \lambda \tau. \text{if } (v \ x \ \tau) \ \tau = \text{true } \land (v \ y \ \tau) = \text{true } \tau \\
& \quad \text{then if } x \ \tau = \text{null } \lor y \ \tau = \text{null } \\
& \quad \text{then } \lfloor \lfloor x \ \tau = \text{null } \land y \ \tau = \text{null } \rfloor \rfloor \\
& \quad \text{else } \lfloor (\text{oid-of } (x \ \tau)) = (\text{oid-of } (y \ \tau)) \rfloor \rfloor \\
& \quad \text{else invalid } \tau
\end{align*}
\]

5.2.1. Logic and Algebraic Layer on Object

Validity and Definedness Properties

We derive the usual laws on definedness for (generic) object equality:

\begin{itemize}
  \item \text{lemma } \text{StrictRefEqObject-defargs:}
  \item \text{assumes } x-val : \tau \models v \ x
  \item \text{shows } \tau \models \text{StrictRefEqObject} x x
\end{itemize}

\begin{proof}
\end{proof}

Symmetry

\begin{itemize}
  \item \text{lemma } \text{StrictRefEqObject-sym :}
  \item \text{assumes } x-val : \tau \models v \ x
  \item \text{shows } \tau \models \text{StrictRefEqObject} x x
\end{itemize}

\begin{proof}
\end{proof}

Execution with Invalid or Null as Argument

\begin{itemize}
  \item \text{lemma } \text{StrictRefEqObject-strict1 [simp, code-unfold] :}
  \item \text{(StrictRefEqObject} x \text{invalid}) = \text{invalid}
\end{itemize}

\begin{proof}
\end{proof}

\begin{itemize}
  \item \text{lemma } \text{StrictRefEqObject-strict2 [simp, code-unfold] :}
\end{itemize}

\begin{proof}
\end{proof}
\[(\text{StrictRefEq}_{\text{Object}} \ invalid \ x) = \ invalid\]

\[\langle \text{proof} \rangle\]

**Context Passing**

**lemma cp-StrictRefEq_{Object}:**
\[(\text{StrictRefEq}_{\text{Object}} \ x \ y \ \tau) = (\text{StrictRefEq}_{\text{Object}} (\lambda \cdot \ x \ \tau) (\lambda \cdot \ y \ \tau)) \ \tau\]

\[\langle \text{proof} \rangle\]

**lemmas cp-intro''[intro!, simp, code-unfold] =**
\[cp-\text{StrictRefEq}_{\text{Object}}[\text{THEN} \ \text{allI}[\text{THEN} \ \text{allI}[\text{THEN} \ \text{allI}[\text{THEN} \ \text{cpI2}]],
\text{of} \ \text{StrictRefEq}_{\text{Object}}]\]

**Behavior vs StrongEq**

It remains to clarify the role of the state invariant \(\text{inv}_\sigma(\sigma)\) mentioned above that states the condition that there is a “one-to-one” correspondence between object representations and oid’s: \(\forall \text{oid} \in \text{dom} \ \sigma. \ \text{oid} = \text{OidOf} \ \lceil \sigma(\text{oid}) \rceil\). This condition is also mentioned in [33, Annex A] and goes back to Richters [35]; however, we state this condition as an invariant on states rather than a global axiom. It can, therefore, not be taken for granted that an oid makes sense both in pre- and post-states of OCL expressions.

We capture this invariant in the predicate \(WFF:\)

**definition WFF :: (\exists::\text{object})st \Rightarrow \text{bool}**

**where WFF \(\tau = ((\forall x \in \text{ran}(\text{heap}(\text{fst} \ \tau))). \ [\text{heap}(\text{fst} \ \tau) \ (\text{oid-of} \ x) = x] \land
(\forall x \in \text{ran}(\text{heap}(\text{snd} \ \tau))). \ [\text{heap}(\text{snd} \ \tau) \ (\text{oid-of} \ x) = x)]\)

It turns out that \(WFF\) is a key-concept for linking strict referential equality to logical equality: in well-formed states (i.e. those states where the self (oid-of) field contains the pointer to which the object is associated to in the state), referential equality coincides with logical equality.

We turn now to the generic definition of referential equality on objects: Equality on objects in a state is reduced to equality on the references to these objects. As in HOL-OCL [6, 8], we will store the reference of an object inside the object in a (ghost) field. By establishing certain invariants (“consistent state”), it can be assured that there is a “one-to-one-correspondence” of objects to their references—and therefore the definition below behaves as we expect.

Generic Referential Equality enjoys the usual properties: (quasi) reflexivity, symmetry, transitivity, substitutivity for defined values. For type-technical reasons, for each concrete object type, the equality \(\doteq\) is defined by generic referential equality.

**theorem StrictRefEq_{Object}-vs-StrongEq:**

**assumes WFF: WFF \(\tau\)**

**and valid-x: \(\tau \models (v \ x)\)**

**and valid-y: \(\tau \models (v \ y)\)**

**and x-present-pre: \(x \ \tau \in \text{ran}(\text{heap}(\text{fst} \ \tau))\)**

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and \( y \text{-present-pre}: y \tau \in \text{ran} (\text{heap}(\text{fst}\ \tau)) \)
and \( x \text{-present-post}: x \tau \in \text{ran} (\text{heap}(\text{snd}\ \tau)) \)
and \( y \text{-present-post}: y \tau \in \text{ran} (\text{heap}(\text{snd}\ \tau)) \)

shows \( (\tau \mid= (\text{StrictRefEq}_\text{Object} \ x \ y)) = (\tau \mid= (x \triangleq y)) \)

\( \langle \text{proof} \rangle \)

\text{theorem} \text{StrictRefEq}_\text{Object}-vs-StrongEq' :
\text{assumes} \ WFF: WFF \ \tau 
and valid-x: \tau \mid= (v \ (x :: (\mathbb{A}::\text{object},'\alpha::\text{null,object})\text{val})) 
and valid-y: \tau \mid= (v \ y) 
and oid-preserve: \bigwedge x. x \in \text{ran} (\text{heap}(\text{fst}\ \tau)) \lor x \in \text{ran} (\text{heap}(\text{snd}\ \tau)) \implies H x \neq \bot \implies \text{oid-of} (H x) = \text{oid-of} x 
and xy-together: x \tau \in H \ ' \text{ran} (\text{heap}(\text{fst}\ \tau)) \land y \tau \in H \ ' \text{ran} (\text{heap}(\text{fst}\ \tau)) \lor x \tau \in H \ ' \text{ran} (\text{heap}(\text{snd}\ \tau)) \land y \tau \in H \ ' \text{ran} (\text{heap}(\text{snd}\ \tau)) 

shows \( (\tau \mid= (\text{StrictRefEq}_\text{Object} \ x \ y)) = (\tau \mid= (x \triangleq y)) \)

\( \langle \text{proof} \rangle \)

So, if two object descriptions live in the same state (both pre or post), the referential equality on objects implies in a WFF state the logical equality.

5.3. Operations on Object

5.3.1. Initial States (for testing and code generation)

\text{definition} \tau_0 :: (\mathbb{A})\text{st} 
where \ \tau_0 \equiv (\langle \text{heap}=\text{Map}\.\text{empty}, \text{assocs}_2=\text{Map}\.\text{empty}, \text{assocs}_3=\text{Map}\.\text{empty} \rangle, 
(\langle \text{heap}=\text{Map}\.\text{empty}, \text{assocs}_2=\text{Map}\.\text{empty}, \text{assocs}_3=\text{Map}\.\text{empty} \rangle)

5.3.2. OclAllInstances

To denote OCL types occurring in OCL expressions syntactically—as, for example, as “argument” of \text{oclAllInstances}()—we use the inverses of the injection functions into the object universes; we show that this is a sufficient “characterization.”

\text{definition} OclAllInstances-generic :: (\mathbb{A}::\text{object}) \text{st} \Rightarrow \mathbb{A} \text{state} \Rightarrow (\mathbb{A}::\text{object} \rightarrow '\alpha) \Rightarrow 
(\mathbb{A}, '\alpha \text{ option option}) \text{Set} 
where OclAllInstances-generic fst-snd H = 
(\lambda \tau. \text{Abs-Set-0} [[ \text{Some } 'i (\langle \text{H } ' \text{ran} (\text{heap} (\text{fst-snd} \ \tau)) \rangle ) - \{ \text{None } \} ]]]

\text{lemma} OclAllInstances-generic-defined: \tau \mid= \delta (\text{OclAllInstances-generic pre-post } H) 
\langle \text{proof} \rangle 

\text{lemma} OclAllInstances-generic-init-empty: 
\text{assumes} \ [\text{simp}]: \bigwedge x. \text{pre-post} (x, x) = x 
\text{shows} \tau_0 \mid= \text{OclAllInstances-generic pre-post } H \triangleq \text{Set}\{\}
\langle \text{proof} \rangle
lemma represented-generic-objects-nonnull:
assumes $A: \tau \models ((\text{OclAllInstances-generic pre-post } (H::(\forall::\text{object} \rightarrow 'a))) \rightarrow \text{includes}(x))$
shows $\tau \models \text{not}(x \triangleq \text{null})$
⟨proof⟩

lemma represented-generic-objects-defined:
assumes $A: \tau \models ((\text{OclAllInstances-generic pre-post } (H::(\forall::\text{object} \rightarrow 'a))) \rightarrow \text{includes}(x))$
shows $\tau \models \delta (\text{OclAllInstances-generic pre-post } H) \land \tau \models \delta x$
⟨proof⟩

One way to establish the actual presence of an object representation in a state is:

lemma represented-generic-objects-in-state:
assumes $A: \tau \models (\text{OclAllInstances-generic pre-post } H) \rightarrow \text{includes}(x)$
shows $x \tau \in (\text{Some o } H) \leftarrow \text{ran}(\text{heap(pre-post } \tau))$
⟨proof⟩

lemma state-update-vs-allInstances-generic-empty:
assumes $\text{simp}: \forall a, \text{pre-post (mk a)} = a$
shows $(\text{mk (\text{heap=empty, assocs}_2=\text{A, assocs}_3=\text{B})} \models \text{OclAllInstances-generic pre-post Type} \leftarrow \text{Set({})})$
⟨proof⟩

Here comes a couple of operational rules that allow to infer the value of oclAllInstances from the context $\tau$. These rules are a special-case in the sense that they are the only rules that relate statements with different $\tau$’s. For that reason, new concepts like “constant contexts P” are necessary (for which we do not elaborate an own theory for reasons of space limitations; in examples, we will prove resulting constraints straightforward by hand).

lemma state-update-vs-allInstances-generic-including:
assumes $\text{simp}: \forall a, \text{pre-post (mk a)} = a$
assumes $\forall x. \sigma' \text{ oid = Some } x \Rightarrow x = \text{Object}$
and $\text{Type Object} \neq \text{None}$
shows $(\text{OclAllInstances-generic pre-post Type})$
$(\text{mk (\text{heap=\sigma' (oid\rightarrow Object), assocs}_2=\text{A, assocs}_3=\text{B})})$
$= ((\text{OclAllInstances-generic pre-post Type}) \rightarrow \text{including}(\lambda -. || \text{ drop (Type Object) } ||))$
$(\text{mk (\text{heap=\sigma',assocs}_2=\text{A, assocs}_3=\text{B})})$
⟨proof⟩

lemma state-update-vs-allInstances-generic-including:
assumes $\text{simp}: \forall a, \text{pre-post (mk a)} = a$
assumes $\forall x. \sigma' \text{ oid = Some } x \Rightarrow x = \text{Object}$
and $\text{Type Object} \neq \text{None}$
shows $(\text{OclAllInstances-generic pre-post Type})$
$(\text{mk (\text{heap=\sigma' (oid\rightarrow Object), assocs}_2=\text{A, assocs}_3=\text{B})})$
$= "$
\[
((\lambda\cdot (OclAllInstances-generic\ pre-post\ Type)
\quad (\text{mk}\ \langle\text{heap}\!\! =\!\! \sigma',\ assocs_2\!\! =\!\! A,\ assocs_3\!\! =\!\! B\rangle\rangle)\rightarrow \text{including}(\lambda\cdot \lfloor\lfloor \text{drop}\ (\text{Type}\ Object)\rfloor\rfloor))
\quad (\text{mk}\ \langle\text{heap}\!\! =\!\! \sigma'(\text{oid} \rightarrow \text{Object}),\ assocs_2\!\! =\!\! A,\ assocs_3\!\! =\!\! B\rangle\rangle)
\langle\text{proof}\rangle
\]

**Lemma** state-update-vs-allInstances-generic-noincluding':

**Assumes** [simp]: \(\forall a.\ \text{pre-post}\ (\text{mk}\ a) = a\)

**Assumes** \(\forall x.\ \sigma\'\ \text{oid} = \text{Some}\ x \implies x = \text{Object}\)

and **Type Object = None**

**Shows** (OclAllInstances-generic pre-post Type)

\[
(\text{mk}\ \langle\text{heap}\!\! =\!\! \sigma'(\text{oid} \rightarrow \text{Object}),\ assocs_2\!\! =\!\! A,\ assocs_3\!\! =\!\! B\rangle\rangle)
= (OclAllInstances-generic pre-post Type)
\]

\[
(\text{mk}\ \langle\text{heap}\!\! =\!\! \sigma',\ assocs_2\!\! =\!\! A,\ assocs_3\!\! =\!\! B\rangle\rangle)
\langle\text{proof}\rangle
\]

**Theorem** state-update-vs-allInstances-generic-ntc:

**Assumes** [simp]: \(\forall a.\ \text{pre-post}\ (\text{mk}\ a) = a\)

**Assumes** oid-def: \(\text{oid} \notin \text{dom}\ \sigma'\)

and **non-type-conform**: Type Object = None

and **cp-ctxt**: \(\text{cp} P\)

and **const-ctxt**: \(\forall X.\ \text{const}\ X \implies \text{const}\ (P\ X)\)

**Shows** (OclAllInstances-generic pre-post Type)

\[
(\text{mk}\ \langle\text{heap}\!\! =\!\! \sigma'(\text{oid} \rightarrow \text{Object}),\ assocs_2\!\! =\!\! A,\ assocs_3\!\! =\!\! B\rangle\rangle) \models P
\]

(\text{OclAllInstances-generic pre-post Type})

\[
(\text{mk}\ \langle\text{heap}\!\! =\!\! \sigma',\ assocs_2\!\! =\!\! A,\ assocs_3\!\! =\!\! B\rangle\rangle) \models P
\langle\text{proof}\rangle
\]

**Theorem** state-update-vs-allInstances-generic-tc:

**Assumes** [simp]: \(\forall a.\ \text{pre-post}\ (\text{mk}\ a) = a\)

**Assumes** oid-def: \(\text{oid} \notin \text{dom}\ \sigma'\)

and **type-conform**: Type Object \(\neq\) None

and **cp-ctxt**: \(\text{cp} P\)

and **const-ctxt**: \(\forall X.\ \text{const}\ X \implies \text{const}\ (P\ X)\)

**Shows** (OclAllInstances-generic pre-post Type)

\[
(\text{mk}\ \langle\text{heap}\!\! =\!\! \sigma'(\text{oid} \rightarrow \text{Object}),\ assocs_2\!\! =\!\! A,\ assocs_3\!\! =\!\! B\rangle\rangle) \models P
\]

(\text{OclAllInstances-generic pre-post Type})

\[
(\text{mk}\ \langle\text{heap}\!\! =\!\! \sigma',\ assocs_2\!\! =\!\! A,\ assocs_3\!\! =\!\! B\rangle\rangle) \models P
\]

\[
(\text{mk}\ \langle\text{heap}\!\! =\!\! \sigma',\ assocs_2\!\! =\!\! A,\ assocs_3\!\! =\!\! B\rangle\rangle) \models P
\langle\text{proof}\rangle
\]

\[
\text{is}\ (\gamma\!\! |\!\! \varphi) = (\gamma'\!\! |\!\! \varphi')
\]

\[
\langle\text{proof}\rangle
\]

**Declare** OclAllInstances-generic-def [simp]
OclAllInstances (@post)

definition OclAllInstances-at-post :: (\texttt{A} :: object \to 'a) \Rightarrow (\texttt{A}, 'a option option) Set
\texttt{- .allInstances}'()

where OclAllInstances-at-post = OclAllInstances-generic snd

lemma OclAllInstances-at-post-defined: \tau \models H .allInstances()
⟨proof⟩

lemma \tau_0 \models H .allInstances() \triangleq Set{}
⟨proof⟩

lemma represented-at-post-objects-nil:
assumes A: \tau \models ((H::('A::object \to 'a)).allInstances()) \to> includes(x)
shows \tau \models not(x \triangleq \texttt{null})
⟨proof⟩

lemma represented-at-post-objects-defined:
assumes A: \tau \models ((H::('A::object \to 'a)).allInstances()) \to> includes(x)
shows \tau \models \delta (H .allInstances()) \land \tau \models \delta x
⟨proof⟩

One way to establish the actual presence of an object representation in a state is:

lemma
assumes A: \tau \models H .allInstances() \to> includes(x)
shows x \tau \in (\texttt{Some o H}) .' ran (heap snd \tau))
⟨proof⟩

lemma state-update-vs-allInstances-at-post-empty:
shows (\sigma, (heap=empty, assoc_2=A, assoc_3=B)) \models Type .allInstances() \triangleq Set{}
⟨proof⟩

Here comes a couple of operational rules that allow to infer the value of oclAllInstances from the context \tau. These rules are a special-case in the sense that they are the only rules that relate statements with different \tau's. For that reason, new concepts like “constant contexts P” are necessary (for which we do not elaborate an own theory for reasons of space limitations; in examples, we will prove resulting constraints straight forward by hand).

lemma state-update-vs-allInstances-at-post-including':
assumes \land x. \sigma' oid = Some x \Rightarrow x = Object
and Type Object \neq \texttt{None}
shows (Type .allInstances())
\sigma, (heap=\sigma'(oid\to Object), assoc_2=A, assoc_3=B)
= ((Type .allInstances()) \to> including(\lambda -. [[ drop (Type Object) ]]))
\sigma, (heap=\sigma',assoc_2=A, assoc_3=B))
⟨proof⟩
lemma state-update-vs-allInstances-at-post-including:
assumes $\forall x. \sigma’ \text{ oid} = \text{ Some } x \implies x = \text{ Object}$
and $\text{ Type } \text{ Object} \neq \text{ None}$
shows $(\text{ Type } \text{. allInstances()}, (\sigma, (\|heap=\sigma’(\text{ oid}\rightarrow\text{ Object}), \text{ associ}_{2}=A, \text{ associ}_{3}=B|}))$
  $= ((\lambda-.(\text{ Type } \text{. allInstances()}, (\sigma, (\|heap=\sigma’, \text{ associ}_{2}=A, \text{ associ}_{3}=B|))) \implies \text{ including}(\lambda-.[.| \text{ drop } (\text{ Type } \text{ Object})] |))$ $\langle proof \rangle$

lemma state-update-vs-allInstances-at-post-noincluding’:
assumes $\forall x. \sigma’ \text{ oid} = \text{ Some } x \implies x = \text{ Object}$
and $\text{ Type } \text{ Object} = \text{ None}$
shows $(\text{ Type } \text{. allInstances()}, (\sigma, (\|heap=\sigma’(\text{ oid}\rightarrow\text{ Object}), \text{ associ}_{2}=A, \text{ associ}_{3}=B|)))$
  $= (\text{ Type } \text{. allInstances()}, (\sigma, (\|heap=\sigma’, \text{ associ}_{2}=A, \text{ associ}_{3}=B|)))$
$\langle proof \rangle$

theorem state-update-vs-allInstances-at-post-ntc:
assumes oid-def: $\text{ oid} \notin \text{ dom } \sigma’$
and $\text{ non-type-conform: } \text{ Type } \text{ Object} = \text{ None}$
and $\text{ cp-ctxt: } \text{ cp } P$
and $\text{ const-ctxt: } \forall X. \text{ const } X \implies \text{ const } (P\ X)$
shows $((\sigma, (\|heap=\sigma’(\text{ oid}\rightarrow\text{ Object}), \text{ associ}_{2}=A, \text{ associ}_{3}=B|)) \vdash (P(\text{ Type } \text{. allInstances()})))$
  $\implies ((\sigma, (\|heap=\sigma’, \text{ associ}_{2}=A, \text{ associ}_{3}=B|)) \vdash (P(\text{ Type } \text{. allInstances()})))$
$\langle proof \rangle$

theorem state-update-vs-allInstances-at-post-tc:
assumes oid-def: $\text{ oid} \notin \text{ dom } \sigma’$
and $\text{ type-conform: } \text{ Type } \text{ Object} \neq \text{ None}$
and $\text{ cp-ctxt: } \text{ cp } P$
and $\text{ const-ctxt: } \forall X. \text{ const } X \implies \text{ const } (P\ X)$
shows $((\sigma, (\|heap=\sigma’(\text{ oid}\rightarrow\text{ Object}), \text{ associ}_{2}=A, \text{ associ}_{3}=B|)) \vdash (P(\text{ Type } \text{. allInstances()})))$
  $\implies ((\sigma, (\|heap=\sigma’, \text{ associ}_{2}=A, \text{ associ}_{3}=B|)) \vdash (P(\text{ Type } \text{. allInstances()}))$
  $\implies (\text{ including}(\lambda-.[.| \langle Type Object\rangle])))$
$\langle proof \rangle$

OclAllInstances (@pre)

definition OclAllInstances-at-pre :: (\$A:: object \rightarrow 'a\$) \Rightarrow (\$A, 'a option option\$) Set
  (- .allInstances@pre'(\$)
where \( \text{OclAllInstances-at-pre} = \text{OclAllInstances-generic} \)

**Lemma**: \( \text{OclAllInstances-at-pre-defined}: \tau \models \delta (H\ .allInstances@pre()) \)

**Proof**

**Lemma**: \( \tau_0 \models H\ .allInstances@pre() \triangleq \text{Set}{} \)

**Lemma** \( \text{represented-at-pre-objects-nonnul}: \)

**Assumptions**: \( A: \tau \models ((\forall x: \text{object} \rightarrow \alpha).\ allInstances@pre()) \rightarrow \text{includes}(x) \)

**Shows**: \( \tau \models \text{not}(x \triangleq \text{null}) \)

**Proof**

**Lemma** \( \text{represented-at-pre-objects-defined}: \)

**Assumptions**: \( A: \tau \models ((\forall x: \text{object} \rightarrow \alpha).\ allInstances@pre()) \rightarrow \text{includes}(x) \)

**Shows**: \( \tau \models \delta (H\ .allInstances@pre()) \land \tau \models \delta x \)

**Proof**

One way to establish the actual presence of an object representation in a state is:

**Lemma** \( \text{state-update-vs-allInstances-at-pre-empty}: \)

**Assumes**: \( \lambda x. \sigma' \cdot \text{oid} = \text{Some} x \Rightarrow x = \text{Object} \)

**And**: \( \text{Type Object} \neq \text{None} \)

**Shows**: \( (\text{heap}=\sigma'(\text{oid} \rightarrow \text{Object}), \text{assoc}_2=A, \text{assoc}_3=B), \sigma) \models \text{Type} \ .\ allInstances@pre() \triangleq \text{Set}{} \)

**Proof**

Here comes a couple of operational rules that allow to infer the value of \( \text{oclAllInstances@pre} \) from the context \( \tau \). These rules are a special-case in the sense that they are the only rules that relate statements with different \( \tau \)'s. For that reason, new concepts like “constant contexts P” are necessary (for which we do not elaborate an own theory for reasons of space limitations; in examples, we will prove resulting constraints straight forward by hand).

**Lemma** \( \text{state-update-vs-allInstances-at-pre-including}: \)

**Assumptions**: \( \lambda x. \sigma' \cdot \text{oid} = \text{Some} x \Rightarrow x = \text{Object} \)

**Shows**: \( (\text{heap}=\sigma'(\text{oid} \rightarrow \text{Object}), \text{assoc}_2=A, \text{assoc}_3=B), \sigma) \models (\text{Type} \ .\ allInstances@pre())->\text{including}(\lambda \cdot (\text{\vert \vert \ drop (\text{Type Object} \ {\mid})})) \)

**Proof**

**Lemma** \( \text{state-update-vs-allInstances-at-pre-including}: \)

**Assumptions**: \( \lambda x. \sigma' \cdot \text{oid} = \text{Some} x \Rightarrow x = \text{Object} \)
\[\text{and Type Object} \neq \text{None}\]
\[
\text{shows } (\text{Type}. \text{allInstances}@\text{pre}())
\]
\[= (\lambda. (\text{Type}. \text{allInstances}@\text{pre}())
\]
\[= (\text{Non-type-conform state-update-vs-allInstances-at-pre-ntc theorem }\langle \text{state-update-vs-allInstances-at-pre-noincluding lemma }\rangle \langle \text{state-update-vs-allInstances-at-pre-tc theorem }\rangle)
\]
\[
\text{proof}
\]

\[\text{lemma state-update-vs-allInstances-at-pre-noincluding'}:\]
\[
\text{assumes } x. \sigma' \text{ oid = Some } x \implies x = \text{Object}
\]
\[\text{and Type Object} = \text{None}\]
\[\text{shows } (\text{Type}. \text{allInstances}@\text{pre}())
\]
\[= (\text{Non-type-conform state-update-vs-allInstances-at-pre-tc theorem }\langle \text{state-update-vs-allInstances-at-pre-tc theorem }\rangle)
\]
\[
\text{proof}
\]

\[\text{theorem state-update-vs-allInstances-at-pre-ntc:}
\]
\[\text{assumes oid-def : oid} \notin \text{dom } \sigma'
\]
\[\text{and non-type-conform : Type Object} = \text{None}
\]
\[\text{and cp-ctxt : } \text{cp } P
\]
\[\text{and const-ctxt : } \bigwedge X. \text{const } X \implies \text{const } (P X)
\]
\[
\text{shows } ((\text{heap}=\sigma'(\text{oid} \rightarrow \text{Object}, \text{assocs}_2=A, \text{assocs}_3=B], \sigma) \models (P(\text{Type}. \text{allInstances}@\text{pre}())))
\]
\[
\text{proof}
\]

\[\text{theorem state-update-vs-allInstances-at-pre-tc:}
\]
\[\text{assumes oid-def : oid} \notin \text{dom } \sigma'
\]
\[\text{and type-conform : Type Object} \neq \text{None}
\]
\[\text{and cp-ctxt : } \text{cp } P
\]
\[\text{and const-ctxt : } \bigwedge X. \text{const } X \implies \text{const } (P X)
\]
\[
\text{shows } ((\text{heap}=\sigma'(\text{oid} \rightarrow \text{Object}, \text{assocs}_2=A, \text{assocs}_3=B], \sigma) \models (P(\text{Type}. \text{allInstances}@\text{pre}())))
\]
\[
\text{proof}
\]

\[\text{post or } \text{pre}
\]

\[\text{theorem StrictRefEqObject-vs-StrongEq'}:\]
\[\text{assumes WFF: WFF } \tau
\]
\[\text{and valid-x: } \tau \models (v (x :: (\text{\_::object,}\alpha::\text{object option option})\text{val}))
\]
\[\text{and valid-y: } \tau \models (v y)
\]

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and oid-preserve: $\forall x. x \in \text{ran}(\text{heap}(\text{fst} \tau)) \lor x \in \text{ran}(\text{heap}(\text{snd} \tau)) \implies$ $\text{oid-of}(H x) = \text{oid-of} x$

and xy-together: $\tau \models ((H . \text{allInstances}()) \rightarrow \text{includes}(x) \text{ and } H . \text{allInstances}()) \rightarrow \text{includes}(y))$

or $\quad (H . \text{allInstances}@\text{pre}()) \rightarrow \text{includes}(x) \text{ and } H . \text{allInstances}@\text{pre}()) \rightarrow \text{includes}(y))$

shows $(\tau \models \text{StrictRefEq(Object, x y)}) = (\tau \models (x \equiv y))$

(proof)

5.3.3. OclIsNew, OclIsDeleted, OclIsMaintained, OclIsAbsent

definition OclIsNew:: (A, 'α::{null,object}) val ⇒ ('A)Boolean (\text{(-).oclIsNew}(''))

where $X . \text{oclIsNew}() \equiv (\lambda \tau . \text{if } (\delta X) \tau = \text{true} \tau$

then $\text{[oid-of } (X \tau) \notin \text{dom(heap(fst} \tau)) \land$

$\text{oid-of } (X \tau) \in \text{dom(heap(snd} \tau))\text{]}$ $\text{else invalid } \tau)$

definition OclIsDeleted:: (A, 'α::{null,object}) val ⇒ ('A)Boolean (\text{(-).oclIsDeleted}(''))

where $X . \text{oclIsDeleted}() \equiv (\lambda \tau . \text{if } (\delta X) \tau = \text{true} \tau$

then $\text{[oid-of } (X \tau) \in \text{dom(heap(fst} \tau)) \land$

$\text{oid-of } (X \tau) \in \text{dom(heap(snd} \tau))\text{]}$ $\text{else invalid } \tau)$

definition OclIsMaintained:: (A, 'α::{null,object}) val ⇒ ('A)Boolean (\text{(-).oclIsMaintained}(''))

where $X . \text{oclIsMaintained}() \equiv (\lambda \tau . \text{if } (\delta X) \tau = \text{true} \tau$

then $\text{[oid-of } (X \tau) \in \text{dom(heap(fst} \tau)) \land$

$\text{oid-of } (X \tau) \in \text{dom(heap(snd} \tau))\text{]}$ $\text{else invalid } \tau)$

definition OclIsAbsent:: (A, 'α::{null,object}) val ⇒ ('A)Boolean (\text{(-).oclIsAbsent}(''))

where $X . \text{oclIsAbsent}() \equiv (\lambda \tau . \text{if } (\delta X) \tau = \text{true} \tau$

then $\text{[oid-of } (X \tau) \notin \text{dom(heap(fst} \tau)) \land$

$\text{oid-of } (X \tau) \notin \text{dom(heap(snd} \tau))\text{]}$ $\text{else invalid } \tau)$

lemma state-split : $\tau \models \delta X \implies$

$\quad \tau \models (X . \text{oclIsNew}()) \lor \tau \models (X . \text{oclIsDeleted}()) \lor$

$\quad \tau \models (X . \text{oclIsMaintained}()) \lor \tau \models (X . \text{oclIsAbsent}())$

(proof)

lemma notNew-vs-others : $\tau \models \delta X \implies$

$\quad (\neg \tau \models (X . \text{oclIsNew}())) = (\tau \models (X . \text{oclIsDeleted}()) \lor$

$\quad \tau \models (X . \text{oclIsMaintained}()) \lor \tau \models (X . \text{oclIsAbsent}())$

(proof)
5.3.4. OclIsModifiedOnly

Definition

The following predicate—which is not part of the OCL standard—provides a simple, but powerful means to describe framing conditions. For any formal approach, be it animation of OCL contracts, test-case generation or die-hard theorem proving, the specification of the part of a system transition that does not change is of primordial importance. The following operator establishes the equality between old and new objects in the state (provided that they exist in both states), with the exception of those objects.

\[
\text{OclIsModifiedOnly} :: (\forall A :: \text{object}, \alpha :: \{\text{null}, \text{object}\}) \Rightarrow \forall A \text{ Boolean}
\]

\[
\text{OclIsModifiedOnly}(\sigma) \equiv (\lambda (\sigma, \sigma') . \text{let } X' = (\text{oid-of } X(\sigma, \sigma')) - X';
\]

\[
S = ((\text{dom heap } \sigma) \cap \text{dom heap } \sigma') - X';
\]

\[
in \text{if } (\delta X) (\sigma, \sigma') = \text{true } (\sigma, \sigma') \land (\forall x \in X(\sigma, \sigma')). x \neq \text{null}
\]

\[
\text{then } \text{valid } (\sigma, \sigma') \text{ } \text{else } \text{invalid } (\sigma, \sigma')
\]

Execution with Invalid or Null or Null Element as Argument

\[
\text{invalid} \Rightarrow \text{OclIsModifiedOnly}() = \text{invalid}
\]

\[
\text{null} \Rightarrow \text{OclIsModifiedOnly}() = \text{invalid}
\]

\[
\text{Context Passing}
\]

\[
\text{cp-OclIsModifiedOnly} : X \Rightarrow \text{OclIsModifiedOnly}() \Rightarrow (\lambda X. X \Rightarrow \text{OclIsModifiedOnly}())
\]

5.3.5. OclSelf

The following predicate—which is not part of the OCL standard—explicitly retrieves in the pre or post state the original OCL expression given as argument.

\[
\text{OclSelf} x H \text{ fst-snd} = (\lambda \tau. \text{ if } (\delta x) \tau = \text{true } \tau
\]

\[
\text{then if oid-of } (x \tau) \in \text{dom(heap fst } \tau)) \land \text{oid-of } (x \tau) \in \text{dom(heap } \text{snd } \tau)
\]

\[
\text{then } H [\text{heap(fst-snd } \tau)](\text{oid-of } (x \tau))]
\]

\[
\text{else invalid } \tau
\]

\[
\text{else invalid } \tau
\]

\[
\text{OclSelf-at-pre} :: (\forall A :: \text{object}, \alpha :: \{\text{null}, \text{object}\}) \Rightarrow
\]

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\[
\forall \alpha \Rightarrow (\exists \alpha::null,object)\) val \ ((-@pre(-))
\]

where \(x @@pre H = OclSelf x H \) fst

definition \(OclSelf-at-post :: (\exists \alpha::null,object)\) val \Rightarrow
\[
(\exists \alpha \Rightarrow (\exists \alpha::null,object))\) val \ ((-@post(-))
\]

where \(x @@post H = OclSelf x H \) snd

5.3.6. Framing Theorem

lemma all-oid-diff:
assumes def-x : \(\tau |\delta x\)
assumes def-X : \(\tau |\delta X\)
assumes def-X' : \(\forall x. x \in [[\text{Rep-Set-0}(X \tau)]] \implies x \neq \text{null}\)
defines \(P \equiv (\lambda a. \text{not} (\text{StrictRefEq} Object x a))\)
shows \((\tau |\delta X \rightarrow \text{forAll}(a | P a)) = (\text{oid-of} (x \tau) \notin \text{oid-of} '([[\text{Rep-Set-0}(X \tau)]])\)
\)

theorem framing:
assumes modifiesclause: \(\tau |\text{excluding}(x)) \rightarrow \text{oclIsModifiedOnly}()\)
andoid-is-typerepr : \(\tau |\text{forAll}(a | \text{not} (x \equiv a))\)
shows \((\tau |\text{forAll}(a | P a))\)
\)

As corollary, the framing property can be expressed with only the strong equality as comparison operator.

theorem framing':
assumes wff : \(WFF \tau\)
assumes modifiesclause: \(\tau |\text{excluding}(x)) \rightarrow \text{oclIsModifiedOnly}()\)
andoid-is-typerepr : \(\tau |\text{forAll}(a | \text{not} (x \equiv a))\)
andoid-preserver : \(\forall x. x \in \text{ran}(\text{heap}(\tau)) \lor x \in \text{ran}(\text{heap}(\text{snd} \tau)) \implies \text{oid-of} (H x) = \text{oid-of} x\)
andzy-together: \(\tau |\text{forAll} y | (H . \text{allInstances}()) \rightarrow \text{includes}(x) \land H . \text{allInstances}() \rightarrow \text{includes}(y)) \lor (H . \text{allInstances}@pre() \rightarrow \text{includes}(x) \land H . \text{allInstances}@pre() \rightarrow \text{includes}(y))\)
shows \((\tau |\text{forAll} P \equiv (x @@pre P))\)
\)

5.3.7. Miscellaneous

lemma pre-post-new: \(\tau |(x . \text{oclIsNew}()) \implies \neg (\tau |\nu(x @@pre H1)) \land \neg (\tau |\nu(x @@post H2))\)
\)

lemma pre-post-old: \(\tau |(x . \text{oclIsDeleted}()) \implies \neg (\tau |\nu(x @@pre H1)) \land \neg (\tau |\nu(x @@post H2))\)
\)

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\textbf{lemma} pre-post-absent: \( \tau \models (x . \text{oclIsAbsent}()) \implies \neg (\tau \models \nu(x @\text{pre } H1)) \land \neg (\tau \models \nu(x @\text{post } H2)) \)

\langle proof \rangle

\textbf{lemma} pre-post-maintained: \( (\tau \models \nu(x @\text{pre } H1) \lor \tau \models \nu(x @\text{post } H2)) \implies \tau \models (x . \text{oclIsMaintained}()) \)

\langle proof \rangle

\textbf{lemma} pre-post-maintained':
\( \tau \models (x . \text{oclIsMaintained}()) \implies (\tau \models \nu(x @\text{pre } (\text{Some } o H1)) \land \tau \models \nu(x @\text{post } (\text{Some } o H2))) \)

\langle proof \rangle

\textbf{lemma} framing-same-state: \( (\sigma, \sigma) \models (x @\text{pre } H \triangleq (x @\text{post } H)) \)

\langle proof \rangle

\textit{end}

\textbf{theory} OCL-tools
\textbf{imports} OCL-core
\textbf{begin}

\textbf{end}

\textbf{theory} OCL-main
\textbf{imports} OCL-lib OCL-state OCL-tools
\textbf{begin}

\textbf{end}
Part III.

Examples
6. The Employee Analysis Model

6.1. The Employee Analysis Model (UML)

theory Employee-AnalysisModel-UMLPart
imports ..:/OCL-main
begin

6.1.1. Introduction

For certain concepts like classes and class-types, only a generic definition for its resulting semantics can be given. Generic means, there is a function outside HOL that “compiles” a concrete, closed-world class diagram into a “theory” of this data model, consisting of a bunch of definitions for classes, accessors, method, casts, and tests for actual types, as well as proofs for the fundamental properties of these operations in this concrete data model.

Such generic function or “compiler” can be implemented in Isabelle on the ML level. This has been done, for a semantics following the open-world assumption, for UML 2.0 in [4, 7]. In this paper, we follow another approach for UML 2.4: we define the concepts of the compilation informally, and present a concrete example which is verified in Isabelle/HOL.

Outlining the Example

We are presenting here an “analysis-model” of the (slightly modified) example Figure 7.3, page 20 of the OCL standard [33]. Here, analysis model means that associations were really represented as relation on objects on the state—as is intended by the standard—rather by pointers between objects as is done in our “design model” (see Section 7.1). To be precise, this theory contains the formalization of the data-part covered by the UML class model (see Figure 6.1):

This means that the association (attached to the association class EmployeeRanking) with the association ends boss and employees is implemented by the attribute boss and the operation employees (to be discussed in the OCL part captured by the subsequent theory).
6.1.2. Example Data-Universe and its Infrastructure

Ideally, the following is generated automatically from a UML class model.

Our data universe consists in the concrete class diagram just of node’s, and implicitly of the class object. Each class implies the existence of a class type defined for the corresponding object representations as follows:

\[
\text{datatype } \text{type} \text{Person} = \text{mk} \text{Person} \text{ oid} \\
\quad \quad \quad \quad \quad \quad \text{int option}
\]

\[
\text{datatype } \text{typeOclAny} = \text{mk} \text{OclAny} \text{ oid} \\
\quad \quad \quad \quad \quad \quad \text{(int option) option}
\]

Now, we construct a concrete “universe of OclAny types” by injection into a sum type containing the class types. This type of OclAny will be used as instance for all respective type-variables.

\[
\text{datatype } \mathcal{A} = \text{in} \text{Person} \text{type} \text{Person} \mid \text{in} \text{OclAny} \text{type} \text{OclAny}
\]

Having fixed the object universe, we can introduce type synonyms that exactly correspond to OCL types. Again, we exploit that our representation of OCL is a “shallow embedding” with a one-to-one correspondance of OCL-types to types of the meta-language HOL.

\[
\begin{align*}
\text{type-synonym } \text{Boolean} & = \mathcal{A} \text{ Boolean} \\
\text{type-synonym } \text{Integer} & = \mathcal{A} \text{ Integer} \\
\text{type-synonym } \text{Void} & = \mathcal{A} \text{ Void} \\
\text{type-synonym } \text{OclAny} & = (\mathcal{A}, \text{typeOclAny option option} \text{ val}) \\
\text{type-synonym } \text{Person} & = (\mathcal{A}, \text{typePerson option option} \text{ val}) \\
\text{type-synonym } \text{Set-Integer} & = (\mathcal{A}, \text{int option option} \text{ Set}) \\
\text{type-synonym } \text{Set-Person} & = (\mathcal{A}, \text{typePerson option option} \text{ Set})
\end{align*}
\]

Just a little check:

\[
\text{typ } \text{Boolean}
\]

To reuse key-elements of the library like referential equality, we have to show that the
object universe belongs to the type class "oclany," i.e., each class type has to provide a function \(oid-of\) yielding the object id (oid) of the object.

**instantiation** type \(\text{Person}::\) object
begin
  definition \(oid-of\text{-type}\text{Person-def}:\) \(oid-of x = (case x of mk\text{Person} oid - \Rightarrow oid)\)
  instance (proof)
end

**instantiation** type \(\text{OclAny}::\) object
begin
  definition \(oid-of\text{-type}\text{OclAny-def}:\) \(oid-of x = (case x of mk\text{OclAny} oid - \Rightarrow oid)\)
  instance (proof)
end

**instantiation** \(\exists::\) object
begin
  definition \(oid-of\cdot\exists-def:\) \(oid-of x = (case x of\)

  \(\text{in}\text{Person} person \Rightarrow oid-of person\)

  | \(\text{in}\text{OclAny} oclany \Rightarrow oid-of oclany)\)

  instance (proof)
end

### 6.1.3. Instantiation of the Generic Strict Equality

We instantiate the referential equality on \(\text{Person}\) and \(\text{OclAny}\)

**defs**(overloaded) \(\text{StrictRefEqObject}\cdot\text{Person} :: (x::\text{Person}) \doteq y \equiv \text{StrictRefEqObject} x y\)

**defs**(overloaded) \(\text{StrictRefEqObject}\cdot\text{OclAny} :: (x::\text{OclAny}) \doteq y \equiv \text{StrictRefEqObject} x y\)

**lemmas**

\(\text{cp-StrictRefEqObject}[of x::\text{Person} y::\text{Person} \tau,\)

simplified \(\text{StrictRefEqObject}\cdot\text{Person}[\text{symmetric}]\]

\(\text{cp-intro}(9)\)

\([of P::\text{Person} \Rightarrow \text{PersonQ}::\text{Person} \Rightarrow \text{Person},\]

simplified \(\text{StrictRefEqObject}\cdot\text{Person}[\text{symmetric}]\]

\(\text{StrictRefEqObject-def} :: [of x::\text{Person} y::\text{Person},\]

simplified \(\text{StrictRefEqObject}\cdot\text{Person}[\text{symmetric}]\]

\(\text{StrictRefEqObject-defargs} :: [of x::\text{Person} y::\text{Person},\]

simplified \(\text{StrictRefEqObject}\cdot\text{Person}[\text{symmetric}]\]

\(\text{StrictRefEqObject-strict1} :: [of x::\text{Person},\]

simplified \(\text{StrictRefEqObject}\cdot\text{Person}[\text{symmetric}]\]

\(\text{StrictRefEqObject-strict2} :: [of x::\text{Person},\]

simplified \(\text{StrictRefEqObject}\cdot\text{Person}[\text{symmetric}]\]

For each Class \(C\), we will have a casting operation \(.\text{oclAsType}(C)\), a test on the actual type \(.\text{oclIsTypeOf}(C)\) as well as its relaxed form \(.\text{oclIsKindOf}(C)\) (corresponding exactly to Java’s \text{instanceof}-operator.

Thus, since we have two class-types in our concrete class hierarchy, we have two op-
6.1.4. OclAsType

Definition

**consts** OclAsType\_OclAny :: ´a \to OclAny ((\_).oclAsType'(OclAny'))

**consts** OclAsType\_Person :: ´a \to Person ((\_).oclAsType'(Person'))

**definition** OclAsType\_OclAny-\_\_A = (\_u. [case u of in\_OclAny a \to a

| in\_Person (mk\_Person oid a) \to mk\_OclAny oid [a]])

**lemma** OclAsType\_OclAny-\_\_A-some: OclAsType\_OclAny\_\_A x \neq None

⟨proof⟩

defs (overloaded) OclAsType\_OclAny\_\_OclAny:

(X::OclAny) .oclAsType(OclAny) \equiv X

defs (overloaded) OclAsType\_OclAny\_\_Person:

(X::Person) .oclAsType(OclAny) \equiv

(\lambda \_ \_ \_e. case X \_ \_e of

| _ e \to invalid \_ e

| [\_] \to null \_ e

| [mk\_Person oid a ] \to [mk\_OclAny oid [a]]

**definition** OclAsType\_Person\_\_A = (\_u. case u of in\_Person p \to [p]

| in\_OclAny (mk\_OclAny oid [a]) \to mk\_Person oid a

| _ \to None)

defs (overloaded) OclAsType\_Person\_\_OclAny:

(X::OclAny) .oclAsType(Person) \equiv

(\lambda \_ \_ \_e. case X \_ \_e of

| _ e \to invalid \_ e

| [\_] \to null \_ e

| [mk\_OclAny oid _ ] \to invalid \_ e (* down-cast exception *)

| [mk\_OclAny oid [a] ] \to [mk\_Person oid a]])

defs (overloaded) OclAsType\_Person\_\_Person:

(X::Person) .oclAsType(Person) \equiv X

**lemmas** [simp] =

OclAsType\_OclAny\_\_OclAny

OclAsType\_Person\_\_Person

Context Passing

**lemma** cp-OclAsType\_OclAny\_\_Person\_\_Person: cp P \implies cp(\_X. (P (X::Person)::Person) .oclAsType(OclAny)))

⟨proof⟩

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lemma \text{cp-OclAsType}_\text{OclAny-OclAny-OclAny}: \text{cp} \ P \implies \text{cp}(\lambda X. (P (X::\text{OclAny})::\text{OclAny}))
\text{oclAsType}(\text{OclAny}))
\langle \text{proof} \rangle

\text{lemma } \text{cp-OclAsType}_\text{Person-Person-Person}: \text{cp} \ P \implies \text{cp}(\lambda X. (P (X::\text{Person})::\text{Person}))
\text{oclAsType}(\text{Person}))
\langle \text{proof} \rangle

\text{lemma } \text{cp-OclAsType}_\text{Person-OclAny-OclAny}: \text{cp} \ P \implies \text{cp}(\lambda X. (P (X::\text{OclAny})::\text{OclAny}))
\text{oclAsType}(\text{Person}))
\langle \text{proof} \rangle

\text{lemma } \text{cp-OclAsType}_\text{OclAny-Person-OclAny}: \text{cp} \ P \implies \text{cp}(\lambda X. (P (X::\text{Person})::\text{OclAny}))
\text{oclAsType}(\text{Person}))
\langle \text{proof} \rangle

\text{lemma } \text{cp-OclAsType}_\text{OclAny-Person-Person}: \text{cp} \ P \implies \text{cp}(\lambda X. (P (X::\text{OclAny})::\text{OclAny}))
\text{oclAsType}(\text{Person}))
\langle \text{proof} \rangle

\text{lemma } \text{cp-OclAsType}_\text{Person-OclAny-Person}: \text{cp} \ P \implies \text{cp}(\lambda X. (P (X::\text{Person})::\text{Person}))
\text{oclAsType}(\text{Person}))
\langle \text{proof} \rangle

\text{lemmas}[\text{simp}]=
\text{cp-OclAsType}_\text{OclAny-Person-Person}
\text{cp-OclAsType}_\text{OclAny-OclAny-OclAny}
\text{cp-OclAsType}_\text{Person-Person-Person}
\text{cp-OclAsType}_\text{Person-OclAny-OclAny}
\text{cp-OclAsType}_\text{OclAny-Person-OclAny}
\text{cp-OclAsType}_\text{OclAny-OclAny-Person}
\text{cp-OclAsType}_\text{Person-Person-OclAny-Person}

\text{Execution with Invalid or Null as Argument}

\text{lemma } \text{OclAsType}_\text{OclAny-OclAny-strict} : (\text{invalid}::\text{OclAny}) . \text{oclAsType}(\text{OclAny}) = \text{invalid}
\langle \text{proof} \rangle

\text{lemma } \text{OclAsType}_\text{OclAny-OclAny-nullstrict} : (\text{null}::\text{OclAny}) . \text{oclAsType}(\text{OclAny}) = \text{null}
\langle \text{proof} \rangle

\text{lemma } \text{OclAsType}_\text{OclAny-Person-strict[simp]} : (\text{invalid}::\text{Person}) . \text{oclAsType}(\text{OclAny}) = \text{invalid}
\langle \text{proof} \rangle

\text{lemma } \text{OclAsType}_\text{OclAny-Person-nullstrict[simp]} : (\text{null}::\text{Person}) . \text{oclAsType}(\text{OclAny}) = \text{null}
\langle \text{proof} \rangle

\text{lemma } \text{OclAsType}_\text{Person-OclAny-strict[simp]} : (\text{invalid}::\text{OclAny}) . \text{oclAsType}(\text{Person}) = \text{invalid}
lemma OclAsType_Person-OclAny-nullstrict[simp] : (null::OclAny) .oclAsType(Person) = null

lemma OclAsType_Person-Person-strict : (invalid::Person) .oclAsType(Person) = invalid

lemma OclAsType_Person-Person-nullstrict : (null::Person) .oclAsType(Person) = null

6.1.5. OclIsTypeOf

Definition

consts OclIsTypeOf_OclAny :: 'α ⇒ Boolean ((OclAny)' .oclIsTypeOf 'OclAny'))
consts OclIsTypeOf_Person :: 'α ⇒ Boolean ((OclAny)' .oclIsTypeOf 'Person'))

defs (overloaded) OclIsTypeOf_OclAny-OclAny:
  (X::OclAny) .oclIsTypeOf(OclAny) ≡
  (λτ. case X τ of
    ⊥ ⇒ invalid τ
    | [⊥] ⇒ true τ (* invalid ?? *)
    | [mkOclAny oid ⊥ ] ⇒ true τ
    | [mkOclAny oid [ ] ] ⇒ false τ)

defs (overloaded) OclIsTypeOf_OclAny-Person:
  (X::Person) .oclIsTypeOf(OclAny) ≡
  (λτ. case X τ of
    ⊥ ⇒ invalid τ
    | [⊥] ⇒ true τ (* invalid ?? *)
    | [ ] ⇒ false τ)

defs (overloaded) OclIsTypeOf_Person-OclAny:
  (X::OclAny) .oclIsTypeOf(Person) ≡
  (λτ. case X τ of
    ⊥ ⇒ invalid τ
    | [ ] ⇒ true τ
    | [mkOclAny oid ⊥ ] ⇒ false τ
    | [mkOclAny oid [ ] ] ⇒ true τ)

defs (overloaded) OclIsTypeOf_Person-Person:
  (X::Person) .oclIsTypeOf(Person) ≡
  (λτ. case X τ of
    ⊥ ⇒ invalid τ
    | [ ] ⇒ true τ)
Context Passing

**lemma** \( cp - OclIsTypeOf \) _Person-Person_
\[ cp(\lambda X. (P(X::Person)::Person).oclIsTypeOf(OclAny)) \]
⟨proof⟩

**lemma** \( cp - OclIsTypeOf \) _OclAny-OclAny_
\[ cp(\lambda X. (P(X::OclAny)::OclAny).oclIsTypeOf(OclAny)) \]
⟨proof⟩

**lemma** \( cp - OclIsTypeOf \) _Person-Person_
\[ cp(\lambda X. (P(X::Person)::Person).oclIsTypeOf(Person)) \]
⟨proof⟩

**lemma** \( cp - OclIsTypeOf \) _OclAny-OclAny_
\[ cp(\lambda X. (P(X::OclAny)::OclAny).oclIsTypeOf(OclAny)) \]
⟨proof⟩

**lemma** \( cp - OclIsTypeOf \) _OclAny-Person_
\[ cp(\lambda X. (P(X::OclAny)::Person).oclIsTypeOf(Person)) \]
⟨proof⟩

**lemma** \( cp - OclIsTypeOf \) _Person-Person_
\[ cp(\lambda X. (P(X::Person)::Person).oclIsTypeOf(Person)) \]
⟨proof⟩

**lemma** \( cp - OclIsTypeOf \) _OclAny-Person_
\[ cp(\lambda X. (P(X::OclAny)::Person).oclIsTypeOf(Person)) \]
⟨proof⟩

**lemma** \( cp - OclIsTypeOf \) _Person-Person_
\[ cp(\lambda X. (P(X::Person)::Person).oclIsTypeOf(Person)) \]
⟨proof⟩

**lemmas** \([simp]=
\]

**Execution with Invalid or Null as Argument**

**lemma** \( OclIsTypeOf \) _OclAny-OclAny-strict1[simp]:
\[ \text{invalid::OclAny).oclIsTypeOf(OclAny) = invalid \]
⟨proof⟩

**lemma** \( OclIsTypeOf \) _OclAny-OclAny-strict2[simp]:
\[ \text{null::OclAny).oclIsTypeOf(OclAny) = true \]
⟨proof⟩

**lemma** \( OclIsTypeOf \) _OclAny-Person-strict1[simp]:
\[ \text{invalid::Person).oclIsTypeOf(OclAny) = invalid \]
\textbf{Up Down Casting}

\textbf{lemma} actualType-larger-staticType:
assumes isdef: \( \tau \models (\delta X) \)
shows \( \tau \models (X :: \text{Person}) \).oclIsTypeOf(OclAny) \( \triangleq \) false
\langle proof \rangle

\textbf{lemma} down-cast-type:
assumes isOclAny: \( \tau \models (X :: \text{OclAny}) \).oclIsTypeOf(OclAny)
and non-null: \( \tau \models (\delta X) \)
shows \( \tau \models (X \ . \ oclAsType(\text{Person})) \triangleq \text{invalid} \)
\langle proof \rangle

\textbf{lemma} down-cast-type':
assumes isOclAny: \( \tau \models (X :: \text{OclAny}) \).oclIsTypeOf(OclAny)
and non-null: \( \tau \models (\delta X) \)
shows \( \tau \models \text{not} \ (v \ (X \ . \ oclAsType(\text{Person}))) \)
\langle proof \rangle

\textbf{lemma} up-down-cast :
assumes isdef: \( \tau \models (\delta X) \)
shows \( \tau \models ((X :: \text{Person}) \ . \ oclAsType(\text{OclAny}) \ . \ oclAsType(\text{Person}) \triangleq X) \)
\langle proof \rangle

\textbf{lemma} up-down-cast-Person-OclAny-Person [simp]:
shows \( ((X :: \text{Person}) \ . \ oclAsType(\text{OclAny}) \ . \ oclAsType(\text{Person}) = X) \)
\langle proof \rangle

\textbf{lemma} up-down-cast-Person-OclAny-Person': assumes \( \tau \models v \ X \)
shows \( \tau \models (((X :: \text{Person}) \ . \ oclAsType(\text{OclAny}) \ . \ oclAsType(\text{Person})) \triangleq X) \)
\langle proof \rangle
lemma up-down-cast-Person-OclAny-Person′′: assumes τ |= v (X :: Person) shows τ |= (X .oclIsTypeOf(Person) implies (X .oclAsType(OclAny) .oclAsType(Person)) ≠ X)
⟨proof⟩

6.1.6. OclIsKindOf

Definition

consts OclIsKindOfOclAny :: 'α ⇒ Boolean ((·).oclIsKindOf'(OclAny'))
consts OclIsKindOfPerson :: 'α ⇒ Boolean ((·).oclIsKindOf'(Person'))

defs (overloaded) OclIsKindOfOclAny-OclAny:
(X::OclAny) .oclIsKindOf(OclAny) ≡
(λτ. case X τ of
  ⊥ ⇒ invalid τ
  | - ⇒ true τ)
defs (overloaded) OclIsKindOfOclAny-Person:
(X::Person) .oclIsKindOf(OclAny) ≡
(λτ. case X τ of
  ⊥ ⇒ invalid τ
  | - ⇒ true τ)
defs (overloaded) OclIsKindOfPerson-OclAny:
(X::OclAny) .oclIsKindOf(Person) ≡
(λτ. case X τ of
  ⊥ ⇒ invalid τ
  | [⊥] ⇒ true τ
  | [[mkOclAny oid ⊥]] ⇒ false τ
  | [[mkOclAny oid [-]]] ⇒ true τ)
defs (overloaded) OclIsKindOfPerson-Person:
(X::Person) .oclIsKindOf(Person) ≡
(λτ. case X τ of
  ⊥ ⇒ invalid τ
  | - ⇒ true τ)

Context Passing

lemma cp-OclIsKindOfOclAny-Person-Person: cp P ⇒
cp(λX.(P(X::Person)::Person).oclIsKindOf(OclAny))
⟨proof⟩
lemma cp-OclIsKindOfOclAny-OclAny-OclAny: cp P ⇒
cp(λX.(P(X::OclAny)::OclAny).oclIsKindOf(OclAny))
⟨proof⟩
lemma cp-OclIsKindOfPerson-Person-Person: cp P ⇒
cp(λX.(P(X::Person)::Person).oclIsKindOf(Person))
⟨proof⟩
lemma \( cp\text{-}\text{OclIsKindOf\_Person\_OclAny\_OclAny} \): 
\[
\begin{align*}
\text{cp}(\lambda X. (P(X::\text{OclAny})::\text{OclAny}).\text{oclIsKindOf}(\text{Person})) & \\
\text{cp} & \Rightarrow & P
\end{align*}
\]

\begin{proof}
\end{proof}

lemma \( cp\text{-}\text{OclIsKindOf\_OclAny\_Person\_OclAny} \): 
\[
\begin{align*}
\text{cp}(\lambda X. (P(X::\text{Person})::\text{OclAny}).\text{oclIsKindOf}(\text{OclAny})) & \\
\text{cp} & \Rightarrow & P
\end{align*}
\]

\begin{proof}
\end{proof}

lemma \( cp\text{-}\text{OclIsKindOf\_OclAny\_OclAny\_Person} \): 
\[
\begin{align*}
\text{cp}(\lambda X. (P(X::\text{OclAny})::\text{OclAny}).\text{oclIsKindOf}(\text{OclAny})) & \\
\text{cp} & \Rightarrow & P
\end{align*}
\]

\begin{proof}
\end{proof}

lemma \( cp\text{-}\text{OclIsKindOf\_Person\_OclAny\_Person} \): 
\[
\begin{align*}
\text{cp}(\lambda X. (P(X::\text{Person})::\text{OclAny}).\text{oclIsKindOf}(\text{Person})) & \\
\text{cp} & \Rightarrow & P
\end{align*}
\]

\begin{proof}
\end{proof}

lemma \( \text{OclIsKindOf\_OclAny\_OclAny}\text{-strict1} \): 
\[
\begin{align*}
\text{invalid}::\text{OclAny}.\text{oclIsKindOf}(\text{OclAny}) & = \\
\text{invalid} & \Rightarrow & \text{invalid}
\end{align*}
\]

\begin{proof}
\end{proof}

lemma \( \text{OclIsKindOf\_OclAny\_OclAny}\text{-strict2} \): 
\[
\begin{align*}
\text{null}::\text{OclAny}.\text{oclIsKindOf}(\text{OclAny}) & = \\
\text{true} & \Rightarrow & \text{true}
\end{align*}
\]

\begin{proof}
\end{proof}

lemma \( \text{OclIsKindOf\_OclAny\_Person}\text{-strict1} \): 
\[
\begin{align*}
\text{invalid}::\text{Person}.\text{oclIsKindOf}(\text{OclAny}) & = \\
\text{invalid} & \Rightarrow & \text{invalid}
\end{align*}
\]

\begin{proof}
\end{proof}

lemma \( \text{OclIsKindOf\_OclAny\_Person}\text{-strict2} \): 
\[
\begin{align*}
\text{null}::\text{Person}.\text{oclIsKindOf}(\text{OclAny}) & = \\
\text{true} & \Rightarrow & \text{true}
\end{align*}
\]

\begin{proof}
\end{proof}

lemma \( \text{OclIsKindOf\_Person\_OclAny\_strict1} \): 
\[
\begin{align*}
\text{invalid}::\text{OclAny}.\text{oclIsKindOf}(\text{Person}) & = \\
\text{invalid} & \Rightarrow & \text{invalid}
\end{align*}
\]

\begin{proof}
\end{proof}
lemma $OclIsKindOf_{Person-OclAny}$-strict2[simp]: $(\text{null}::\text{OclAny}) . oclIsKindOf(\text{Person}) = \text{true}$
(proof)

lemma $OclIsKindOf_{Person-Person}$-strict1[simp]: $(\text{invalid}::\text{Person}) . oclIsKindOf(\text{Person}) = \text{invalid}$
(proof)

lemma $OclIsKindOf_{Person-Person}$-strict2[simp]: $(\text{null}::\text{Person}) . oclIsKindOf(\text{Person}) = \text{true}$
(proof)

Up Down Casting

lemma actualKind-larger-staticKind:
assumes isdef: $\tau \models (\delta \ X)$
shows $\tau \models (X::\text{Person}) . oclIsKindOf(\text{OclAny}) \triangleq \text{true}$
(proof)

lemma down-cast-kind:
assumes isOclAny: $\neg \tau \models (X::\text{OclAny}) . oclIsKindOf(\text{Person})$
and non-null: $\tau \models (\delta \ X)$
shows $\tau \models (X . oclAsType(\text{Person})) \triangleq \text{invalid}$
(proof)

6.1.7. OclAllInstances

To denote OCL-types occuring in OCL expressions syntactically—as, for example, as “argument” of oclAllInstances()—we use the inverses of the injection functions into the object universes; we show that this is sufficient “characterization.”

definition $\text{Person} \equiv \text{OclAsType}_{\text{Person}}$-\texttt{A}$
definition $\text{OclAny} \equiv \text{OclAsType}_{\text{OclAny}}$-\texttt{A}$
lemmas [simp] = $\text{Person-def} \ \text{OclAny-def}$

lemma OclAllInstances-genericOclAny-exec: $\text{OclAllInstances-generic pre-post OclAny} = (\lambda \tau . \text{Abs-Set-0} \ [\ [ \text{Some } \cdot \text{OclAny} \cdot \text{ran} \ (\text{heap} \ (\text{pre-post} \ \tau)) ]])$
(proof)

lemma OclAllInstances-at-postOclAny-exec: $\text{OclAny . allInstances}() = (\lambda \tau . \text{Abs-Set-0} \ [\ [ \text{Some } \cdot \text{OclAny} \cdot \text{ran} \ (\text{heap} \ (\text{snd} \ \tau)) ]])$
(proof)

lemma OclAllInstances-at-preOclAny-exec: $\text{OclAny . allInstances}@\text{pre}() = (\lambda \tau . \text{Abs-Set-0} \ [\ [ \text{Some } \cdot \text{OclAny} \cdot \text{ran} \ (\text{heap} \ (\text{fst} \ \tau)) ]])$
(proof)

OclIsTypeOf

lemma OclAny-allInstances-generic-oclIsTypeOfOclAny1:
assumes [simp]: $\forall x . \text{pre-post} \ (x, x) = x$
shows $\exists \tau. \ (\tau \models ((OclAllInstances-generic\ pre\ post\ OclAny) \rightarrow \forall X | X .oclIsTypeOf(OclAny))))$

<proof>

lemma OclAny-allInstances-at-post-oclIsTypeOfOclAny:
$\exists \tau. \ (\tau \models (OclAny .allInstances() \rightarrow \forall X | X .oclIsTypeOf(OclAny))))$

<proof>

lemma OclAny-allInstances-at-pre-oclIsTypeOfOclAny:
$\exists \tau. \ (\tau \models (OclAny .allInstances@pre() \rightarrow \forall X | X .oclIsTypeOf(OclAny))))$

<proof>

lemma OclAny-allInstances-generic-oclIsTypeOfOclAny:
assumes [simp]: $\forall x.\ \text{pre-post}(x, x) = x$
shows $\exists \tau. \ (\tau \models \neg ((OclAllInstances-generic\ pre\ post\ OclAny) \rightarrow \forall X | X .oclIsTypeOf(OclAny))))$

<proof>

lemma Person-allInstances-generic-oclIsTypeOfPerson:
$\tau \models ((OclAllInstances-generic\ pre\ post\ Person) \rightarrow \forall X | X .oclIsTypeOf(Person)))$

<proof>

lemma Person-allInstances-at-post-oclIsTypeOfPerson:
$\tau \models (Person .allInstances() \rightarrow \forall X | X .oclIsTypeOf(Person)))$

<proof>

lemma Person-allInstances-at-pre-oclIsTypeOfPerson:
$\tau \models (Person .allInstances@pre() \rightarrow \forall X | X .oclIsTypeOf(Person)))$

<proof>

OclIsKindOf

lemma OclAny-allInstances-generic-oclIsKindOfOclAny:
$\tau \models ((OclAllInstances-generic\ pre\ post\ OclAny) \rightarrow \forall X | X .oclIsKindOf(OclAny))))$

<proof>

lemma OclAny-allInstances-at-post-oclIsKindOfOclAny:
$\tau \models (OclAny .allInstances() \rightarrow \forall X | X .oclIsKindOf(OclAny))))$

<proof>

lemma OclAny-allInstances-at-pre-oclIsKindOfOclAny:
\[ \tau \models (\text{OclAny}.\text{allInstances}@\text{pre}) -> \text{forall}(X | X.\text{oclIsKindOf}(\text{OclAny})) \]

lemma Person-allInstances-generic-oclIsKindOf\text{OclAny}:
\[ \tau \models ((\text{OclAllInstances-generic pre-post Person}) -> \text{forall}(X | X.\text{oclIsKindOf}(\text{OclAny}))) \]

lemma Person-allInstances-at-post-oclIsKindOf\text{OclAny}:
\[ \tau \models (\text{Person}.\text{allInstances}) -> \text{forall}(X | X.\text{oclIsKindOf}(\text{OclAny})) \]

lemma Person-allInstances-at-pre-oclIsKindOf\text{OclAny}:
\[ \tau \models (\text{Person}.\text{allInstances}@\text{pre}) -> \text{forall}(X | X.\text{oclIsKindOf}(\text{OclAny})) \]

lemma Person-allInstances-generic-oclIsKindOf\text{Person}:
\[ \tau \models ((\text{OclAllInstances-generic pre-post Person}) -> \text{forall}(X | X.\text{oclIsKindOf}(\text{Person}))) \]

lemma Person-allInstances-at-post-oclIsKindOf\text{Person}:
\[ \tau \models (\text{Person}.\text{allInstances}) -> \text{forall}(X | X.\text{oclIsKindOf}(\text{Person})) \]

lemma Person-allInstances-at-pre-oclIsKindOf\text{Person}:
\[ \tau \models (\text{Person}.\text{allInstances}@\text{pre}) -> \text{forall}(X | X.\text{oclIsKindOf}(\text{Person})) \]

6.1.8. The Accessors (any, boss, salary)

Should be generated entirely from a class-diagram.

Definition (of the association Employee-Boss)

We start with a oid for the association; this oid can be used in presence of association classes to represent the association inside an object, pretty much similar to the Employee\_DesignModel\_UMLPart, where we stored an oid inside the class as “pointer.”

definition oid\text{PersonBOSS} :: oid where oid\text{PersonBOSS} = 10

From there on, we can already define an empty state which must contain for oid\text{PersonBOSS} the empty relation (encoded as association list, since there are associations with a Sequence-like structure).

definition eval-extract :: (A, (a::object) option option) val
\[ \Rightarrow (\text{oid} \Rightarrow (A, c::null) val) \]
\[ \Rightarrow (A, c::null) val \]

where eval-extract X f = (λ τ. case X τ of
\[ \downarrow \Rightarrow \text{invalid } τ \text{ (exception propagation *)} \]
\[ | \downarrow \downarrow | \Rightarrow \text{invalid } τ \text{ (dereferencing null pointer *)} \]
\[ \text{definition } \text{choose}_2 \cdot 1 = \text{fst} \]
\[ \text{definition } \text{choose}_2 \cdot 2 = \text{snd} \]
\[ \text{definition } \text{choose}_3 \cdot 1 = \text{fst} \]
\[ \text{definition } \text{choose}_3 \cdot 2 = \text{fst} \circ \text{snd} \]
\[ \text{definition } \text{choose}_3 \cdot 3 = \text{snd} \circ \text{snd} \]

\[ \text{definition } \text{deref-assocs}_2 :: \left( \mathbb{A} \text{ state} \times \mathbb{A} \text{ state} \Rightarrow \mathbb{A} \text{ state} \right) \Rightarrow \left( \text{oid} \times \text{oid} \Rightarrow \text{oid} \times \text{oid} \right) \Rightarrow \text{oid} \Rightarrow \left( \text{oid list} \Rightarrow \text{oid} \Rightarrow \left( \mathbb{A}, f::\text{null}\text{val} \right) \text{val} \right) \Rightarrow \text{oid} \Rightarrow \left( \mathbb{A}, f::\text{null}\text{val} \right) \text{val} \]

\[ \text{where } \text{deref-assocs}_2 \text{ pre-post to-from assoc-oid f oid } = \]
\[ (\lambda \tau. \text{case } (\text{assocs}_2 (\text{pre-post } \tau)) \text{ assoc-oid of } \ [ S ] \Rightarrow f (\text{map } (\text{choose}_2 \cdot 2 \circ \text{to-from}) \text{ filter } (\lambda p. \text{choose}_2 \cdot 1 ((\text{to-from } p) = \text{oid}) S)) \text{ oid } \tau \]
\[ \mid - \Rightarrow \text{invalid } \tau \]

The \textit{pre-post}-parameter is configured with \textit{fst} or \textit{snd}, the \textit{to-from}-parameter either with the identity \textit{id} or the following combinator \textit{switch}:

\[ \text{definition } \text{switch}_2 \cdot 1 = \text{id} \]
\[ \text{definition } \text{switch}_2 \cdot 2 = (\lambda (x,y). (y,x)) \]
\[ \text{definition } \text{switch}_3 \cdot 1 = \text{id} \]
\[ \text{definition } \text{switch}_3 \cdot 2 = (\lambda (x,y,z). (x,z,y)) \]
\[ \text{definition } \text{switch}_3 \cdot 3 = (\lambda (x,y,z). (y,x,z)) \]
\[ \text{definition } \text{switch}_3 \cdot 4 = (\lambda (x,y,z). (y,z,x)) \]
\[ \text{definition } \text{switch}_3 \cdot 5 = (\lambda (x,y,z). (z,x,y)) \]
\[ \text{definition } \text{switch}_3 \cdot 6 = (\lambda (x,y,z). (z,y,x)) \]

\[ \text{definition } \text{select-object } :: \left( (\mathbb{A}, b::\text{null}\text{val}) \right) \Rightarrow \left( (\mathbb{A}, 'b)\text{val} \Rightarrow (\mathbb{A}, 'c)\text{val} \Rightarrow (\mathbb{A}, 'b)\text{val} \right) \Rightarrow \left( (\mathbb{A}, 'b)\text{val} \Rightarrow (\mathbb{A}, 'd)\text{val} \right) \Rightarrow \left( \text{oid } \Rightarrow (\mathbb{A}, 'c::\text{null}\text{val}) \right) \Rightarrow \text{oid list} \Rightarrow \text{oid} \Rightarrow (\mathbb{A}, 'd)\text{val} \]

\[ \text{where } \text{select-object } \text{mt incl smash deref l oid } = \text{smash}(\text{foldl incl mt } (\text{map deref } l)) \]
\[(\star \text{smash returns null with mt in input (in this case, object contains null pointer) } \star)\]

The continuation \textit{f} is usually instantiated with a smashing function which is either the identity \textit{id} or, for 0..1 cardinalities of associations, the \textit{OclANY}-selector which also handles the \textit{null}-cases appropriately. A standard use-case for this combinator is for example:

\[ \text{term } (\text{select-object } \text{mtSet OclIncluding OclANY } f \ l \ oid )::(\mathbb{A}, 'a::\text{null}\text{val}) \]

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definition \(\text{deref_oid}_{\text{Person}}:: (\mathcal{A} \text{ state} \times \mathcal{A} \text{ state} \Rightarrow \mathcal{A} \text{ state}) \Rightarrow (\text{type} \text{Person} \Rightarrow (\mathcal{A} \text{, 'c::null}val) \Rightarrow \text{oid} \Rightarrow (\mathcal{A} \text{, 'c::null}val)

\text{where} \text{deref_oid}_{\text{Person}} \text{fst-snd f oid} = (\lambda x. \text{case (heap (fst-snd \ x)) oid of}
|\text{ inPerson obj } \Rightarrow f \text{ obj } \tau
|\text{-} \Rightarrow \text{invalid } \tau)

\text{definition} \text{deref_oid}_{\text{OclAny}}:: (\mathcal{A} \text{ state} \times \mathcal{A} \text{ state} \Rightarrow \mathcal{A} \text{ state}) \Rightarrow (\text{type} \text{OclAny} \Rightarrow (\mathcal{A} \text{, 'c::null}val) \Rightarrow \text{oid} \Rightarrow (\mathcal{A} \text{, 'c::null}val)

\text{where} \text{deref_oid}_{\text{OclAny}} \text{fst-snd f oid} = (\lambda x. \text{case (heap (fst-snd \ x)) oid of}
|\text{ inOclAny obj } \Rightarrow f \text{ obj } \tau
|\text{-} \Rightarrow \text{invalid } \tau)

pointer undefined in state or not referencing a type conform object representation

\text{definition} \text{select}_{\text{OclAnyANY}} f = (\lambda x. \text{case } \text{X } \text{of}
(\text{mkOclAny - ⊥}) \Rightarrow \text{null}
| (\text{mkOclAny - } \text{any}) \Rightarrow f (\lambda x. \text{-} [\text{x}]) \text{ any})

\text{definition} \text{select}_{\text{PersonBOSS}} f = \text{select-object mtSet } \text{OclIncluding } \text{OclANY} \ (f (\lambda x. \text{-} [\text{x}]))

\text{definition} \text{select}_{\text{PersonSALARY}} f = (\lambda x. \text{case } \text{X } \text{of}
(\text{mkPerson - ⊥}) \Rightarrow \text{null}
| (\text{mkPerson - } \text{salary}) \Rightarrow f (\lambda x. \text{-} [\text{x}]) \text{ salary})

\text{definition} \text{deref-assocs}_{2}\text{BOSS fst-snd f} = (\lambda \text{mkPerson oid} - \Rightarrow \text{deref-assocs}_{2} \text{fst-snd switch}_{2-1} \text{oid}_{\text{PersonBOSS}} f \text{ oid})

\text{definition} \text{in-pre-state} = \text{fst}
\text{definition} \text{in-post-state} = \text{snd}

\text{definition} \text{reconst-basetype} = (\lambda \text{ convert } \text{x. convert } \text{x})

\text{definition} \text{dot}_{\text{OclAnyANY}}:: \text{OclAny} \Rightarrow - (((1(-).any) 50)
\text{where} \text{X}. \text{any} = \text{eval-extract } \text{X}
\text{(deref-oid}_{\text{OclAny}} \text{ in-post-state}
\text{(select}_{\text{OclAnyANY}} \text{any}
\text{reconst-basetype}))

\text{definition} \text{dot}_{\text{PersonBOSS}}:: \text{Person} \Rightarrow \text{Person} ((1(-).boss) 50)
\text{where} (\text{X}). \text{boss} = \text{eval-extract } \text{X}
\text{(deref-oid}_{\text{Person}} \text{ in-post-state}

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(deref-assocs2boss in-post-state
  (select person boss
  (deref-oid person in-post-state))))

definition dot(salary) person :: person \rightarrow integer ((1 \cdot .salary) 50)
where (x).salary = eval-extract x
  (deref-oid person in-post-state
  (select person salary
  reconst-basetype))

definition dot(\text{any}@\text{at-pre}) :: ocl\text{any} \Rightarrow - ((1 \cdot .\text{any}@\text{pre}) 50)
where (x).\text{any}@\text{pre} = eval-extract x
  (deref-oid ocl\text{any} in-pre-state
  (select ocl\text{any} \text{any}
  reconst-basetype))

definition dot(boss@pre) person :: person \Rightarrow person ((1 \cdot .boss@pre) 50)
where (x).boss@pre = eval-extract x
  (deref-oid person in-pre-state
  (deref-assocs2boss in-pre-state
   (select person boss
   (deref-oid person in-pre-state))))

definition dot(salary@pre) person :: person \Rightarrow integer ((1 \cdot .salary@pre) 50)
where (x).salary@pre = eval-extract x
  (deref-oid person in-pre-state
  (select person salary
  reconst-basetype))

lemmas simp =
  dot(\text{any}@\text{at-pre}) def
  dot(boss@pre) person def
  dot(salary@pre) person def
  dot(\text{any}@\text{at-pre}) def
  dot(boss@pre) person def
  dot(salary@pre) person def

\textbf{Context Passing}

lemmas simp = eval-extract-def

lemma cp-dot(\text{any}@\text{at-pre}) : ((x).any) \tau = ((\lambda\cdot . x \tau).any) \tau \langle proof \rangle
lemma cp-dot(boss@pre) person : ((x).boss) \tau = ((\lambda\cdot . x \tau).boss) \tau \langle proof \rangle
lemma cp-dot(salary@pre) person : ((x).salary) \tau = ((\lambda\cdot . x \tau).salary) \tau \langle proof \rangle

lemma cp-dot(\text{any}@\text{at-pre}) : ((x).\text{any}@\text{pre}) \tau = ((\lambda\cdot . x \tau).\text{any}@\text{pre}) \tau \langle proof \rangle
lemma cp-dot(boss@pre) person : ((x).boss@pre) \tau = ((\lambda\cdot . x \tau).boss@pre) \tau \langle proof \rangle
lemma cp-dot(salary@pre) person : ((x).salary@pre) \tau = ((\lambda\cdot . x \tau).salary@pre) \tau \langle proof \rangle
lemmas cp-dot_OclAny\ANY-\[I \text{ [simp, intro]} = \\
cp-dot_OclAny\ANY[\text{THEN allI[THEN all]}], \\
of \lambda X \cdot X \lambda - \tau. \tau, \text{ THEN cpI}]
lemmas cp-dot_OclAny\ANY-\text{at-pre-I} \text{ [simp, intro]} = \\
cp-dot_OclAny\ANY-\text{at-pre}[\text{THEN allI[THEN all]}], \\
of \lambda X \cdot X \lambda - \tau. \tau, \text{ THEN cpI}]
lemmas cp-dot_{\text{Person}}\BOSS-\[I \text{ [simp, intro]} = \\
cp-dot_{\text{Person}}\BOSS[\text{THEN allI[THEN all]}], \\
of \lambda X \cdot X \lambda - \tau. \tau, \text{ THEN cpI}]
lemmas cp-dot_{\text{Person}}\BOSS-\text{at-pre-I} \text{ [simp, intro]} = \\
cp-dot_{\text{Person}}\BOSS-\text{at-pre}[\text{THEN allI[THEN all]}], \\
of \lambda X \cdot X \lambda - \tau. \tau, \text{ THEN cpI}]
lemmas cp-dot_{\text{Person}}\SALARY-\[I \text{ [simp, intro]} = \\
cp-dot_{\text{Person}}\SALARY[\text{THEN allI[THEN all]}], \\
of \lambda X \cdot X \lambda - \tau. \tau, \text{ THEN cpI}]
lemmas cp-dot_{\text{Person}}\SALARY-\text{at-pre-I} \text{ [simp, intro]} = \\
cp-dot_{\text{Person}}\SALARY-\text{at-pre}[\text{THEN allI[THEN all]}], \\
of \lambda X \cdot X \lambda - \tau. \tau, \text{ THEN cpI}]

\textbf{Execution with Invalid or Null as Argument}

\textbf{lemma} \textit{dot_OclAny\ANY-nullstrict} \text{ [simp]}: (null).\text{any = invalid} \\
\langle \text{proof} \rangle
\textbf{lemma} \textit{dot_OclAny\ANY-at-pre-nullstrict} \text{ [simp]}: (null).\text{any@pre = invalid} \\
\langle \text{proof} \rangle
\textbf{lemma} \textit{dot_OclAny\ANY-strict} \text{ [simp]}: (invalid).\text{any = invalid} \\
\langle \text{proof} \rangle
\textbf{lemma} \textit{dot_OclAny\ANY-at-pre-strict} \text{ [simp]}: (invalid).\text{any@pre = invalid} \\
\langle \text{proof} \rangle

\textbf{lemma} \textit{dot_{\text{Person}}\BOSS-nullstrict} \text{ [simp]}: (null).\text{boss = invalid} \\
\langle \text{proof} \rangle
\textbf{lemma} \textit{dot_{\text{Person}}\BOSS-at-pre-nullstrict} \text{ [simp]}: (null).\text{boss@pre = invalid} \\
\langle \text{proof} \rangle
\textbf{lemma} \textit{dot_{\text{Person}}\BOSS-strict} \text{ [simp]}: (invalid).\text{boss = invalid} \\
\langle \text{proof} \rangle
\textbf{lemma} \textit{dot_{\text{Person}}\BOSS-at-pre-strict} \text{ [simp]}: (invalid).\text{boss@pre = invalid} \\
\langle \text{proof} \rangle

\textbf{lemma} \textit{dot_{\text{Person}}\SALARY-nullstrict} \text{ [simp]}: (null).\text{salary = invalid} \\
\langle \text{proof} \rangle
\textbf{lemma} \textit{dot_{\text{Person}}\SALARY-at-pre-nullstrict} \text{ [simp]}: (null).\text{salary@pre = invalid} \\
\langle \text{proof} \rangle
\textbf{lemma} \textit{dot_{\text{Person}}\SALARY-strict} \text{ [simp]}: (invalid).\text{salary = invalid} \\
\langle \text{proof} \rangle
6.1.9. A Little Infrastructure on Example States

The example we are defining in this section comes from the figure 6.2.

\[ \text{definition } \text{OclInt1000} \equiv (1000) \quad \text{where } \text{OclInt1000} = (\lambda \cdot [1000]) \]
\[ \text{definition } \text{OclInt1200} \equiv (1200) \quad \text{where } \text{OclInt1200} = (\lambda \cdot [1200]) \]
\[ \text{definition } \text{OclInt1300} \equiv (1300) \quad \text{where } \text{OclInt1300} = (\lambda \cdot [1300]) \]
\[ \text{definition } \text{OclInt1800} \equiv (1800) \quad \text{where } \text{OclInt1800} = (\lambda \cdot [1800]) \]
\[ \text{definition } \text{OclInt2600} \equiv (2600) \quad \text{where } \text{OclInt2600} = (\lambda \cdot [2600]) \]
\[ \text{definition } \text{OclInt2900} \equiv (2900) \quad \text{where } \text{OclInt2900} = (\lambda \cdot [2900]) \]
\[ \text{definition } \text{OclInt3200} \equiv (3200) \quad \text{where } \text{OclInt3200} = (\lambda \cdot [3200]) \]
\[ \text{definition } \text{OclInt3500} \equiv (3500) \quad \text{where } \text{OclInt3500} = (\lambda \cdot [3500]) \]

\[ \text{definition } \text{oid0} \equiv 0 \]
\[ \text{definition } \text{oid1} \equiv 1 \]
\[ \text{definition } \text{oid2} \equiv 2 \]
\[ \text{definition } \text{oid3} \equiv 3 \]
\[ \text{definition } \text{oid4} \equiv 4 \]
\[ \text{definition } \text{oid5} \equiv 5 \]
\[ \text{definition } \text{oid6} \equiv 6 \]
\[ \text{definition } \text{oid7} \equiv 7 \]
\[ \text{definition } \text{oid8} \equiv 8 \]

\[ \text{definition } \text{person1} \equiv \text{mkPerson} \text{oid0} [1300] \]
\[ \text{definition } \text{person2} \equiv \text{mkPerson} \text{oid1} [1800] \]
\[ \text{definition } \text{person3} \equiv \text{mkPerson} \text{oid2} \text{None} \]
\[ \text{definition } \text{person4} \equiv \text{mkPerson} \text{oid3} [2900] \]
\[ \text{definition } \text{person5} \equiv \text{mkPerson} \text{oid4} [3500] \]
\[ \text{definition } \text{person6} \equiv \text{mkPerson} \text{oid5} [2500] \]
\[ \text{definition } \text{person7} \equiv \text{mkOclAny} \text{oid6} [3200] \]
\[ \text{definition } \text{person8} \equiv \text{mkOclAny} \text{oid7} \text{None} \]
\[ \text{definition } \text{person9} \equiv \text{mkPerson} \text{oid8} [0] \]

\[ \text{definition } \]

\[ \text{lemma } \text{dotPerson.SALARY}@-\text{at-strict} \quad \text{simp} : (\text{invalid}.\text{salary}@\text{pre} = \text{invalid} \langle \text{proof} \rangle \]
\[ \sigma_1 \equiv \langle \text{heap} = \text{empty}(oid0 \mapsto \text{in\_person}(mk\_\text{person} oid0 \mid 1000)) \]
\[ (oid1 \mapsto \text{in\_person}(mk\_\text{person} oid1 \mid 1200)) \]
\[ (*oid2*) \]
\[ (oid3 \mapsto \text{in\_person}(mk\_\text{person} oid3 \mid 2600)) \]
\[ (oid4 \mapsto \text{in\_person} \text{person5}) \]
\[ (oid5 \mapsto \text{in\_person}(mk\_\text{person} oid5 \mid 2300)) \]
\[ (*oid6*) \]
\[ (*oid7*) \]
\[ (oid8 \mapsto \text{in\_person} \text{person9}) \]
\[ \text{assocs}_2 = \text{empty}(oid\_\text{person}\text{BOSS} \mapsto [\langle(oid0, oid1), (oid3, oid4), (oid5, oid3)\rangle]) \]
\[ \text{assocs}_3 = \text{empty} \]

**Definition**

\[ \sigma_1' \equiv \langle \text{heap} = \text{empty}(oid0 \mapsto \text{in\_person} \text{person1}) \]
\[ (oid1 \mapsto \text{in\_person} \text{person2}) \]
\[ (oid2 \mapsto \text{in\_person} \text{person3}) \]
\[ (oid3 \mapsto \text{in\_person} \text{person4}) \]
\[ (*oid4*) \]
\[ (oid5 \mapsto \text{in\_person} \text{person6}) \]
\[ (oid6 \mapsto \text{oCl\_\text{any}} \text{person7}) \]
\[ (oid7 \mapsto \text{oCl\_\text{any}} \text{person8}) \]
\[ (oid8 \mapsto \text{in\_person} \text{person9}) \]
\[ \text{assocs}_2 = \text{empty}(oid\_\text{person}\text{BOSS} \mapsto \langle(oid0, oid1), (oid3, oid4), (oid5, oid6)\rangle) \]
\[ \text{assocs}_3 = \text{empty} \]

**Lemma** basic-\(\tau\)-wff: \(\text{WFF}(\sigma_1, \sigma_1')\)

**Proof**

**Lemma** [simp.code-unfold]: \(\text{dom}(\text{heap} \sigma_1) = \{ oid0, oid1, (*, oid2*), oid3, oid4, oid5 (*, oid6, oid7*), oid8 \} \)

**Proof**

**Lemma** [simp.code-unfold]: \(\text{dom}(\text{heap} \sigma_1') = \{ oid0, oid1, oid2, oid3, (*, oid4*), oid5, oid6, oid7, oid8 \} \)

**Proof**

**Definition** \(\text{X\_\text{person}1} :: \text{Person} \equiv \lambda - .[\text{person1}] \)

**Definition** \(\text{X\_\text{person}2} :: \text{Person} \equiv \lambda - .[\text{person2}] \)

**Definition** \(\text{X\_\text{person}3} :: \text{Person} \equiv \lambda - .[\text{person3}] \)

**Definition** \(\text{X\_\text{person}4} :: \text{Person} \equiv \lambda - .[\text{person4}] \)

**Definition** \(\text{X\_\text{person}5} :: \text{Person} \equiv \lambda - .[\text{person5}] \)

**Definition** \(\text{X\_\text{person}6} :: \text{Person} \equiv \lambda - .[\text{person6}] \)

**Definition** \(\text{X\_\text{person}7} :: \text{Ocl\_\text{any}} \equiv \lambda - .[\text{person7}] \)

**Definition** \(\text{X\_\text{person}8} :: \text{Ocl\_\text{any}} \equiv \lambda - .[\text{person8}] \)

**Definition** \(\text{X\_\text{person}9} :: \text{Person} \equiv \lambda - .[\text{person9}] \)

**Lemma** [code-unfold]: \((x:\text{Person}) \equiv y) = \text{Strict\_Ref\_Eq\_\text{object}} x y\) (proof)
lemma \texttt{[code-unfold]}: \((x::\text{OclAny}) \doteq y) = \text{StrictRefEq} @ x y \langle proof \rangle

lemmas \texttt{[simp,code-unfold]} =
\text{OclAsType}\text{OclAny-}\text{OclAny}
\text{OclAsType}\text{OclAny-Person}
\text{OclAsType}\text{Person-}\text{OclAny}
\text{OclAsType}\text{Person-Person}
\text{OclIsTypeOf}\text{OclAny-}\text{OclAny}
\text{OclIsTypeOf}\text{OclAny-Person}
\text{OclIsTypeOf}\text{Person-}\text{OclAny}
\text{OclIsTypeOf}\text{Person-Person}
\text{OclIsKindOf}\text{OclAny-}\text{OclAny}
\text{OclIsKindOf}\text{OclAny-Person}
\text{OclIsKindOf}\text{Person-}\text{OclAny}
\text{OclIsKindOf}\text{Person-Person}

value \vee \alpha \cdot (\alpha,\sigma) \vdash (\text{Person}.\text{salary} <\text{\textless} 1000)
value \vee \alpha \cdot (\alpha,\sigma) \vdash (\text{Person}.\text{salary} \geq 1300)
value \vee \alpha \cdot (\alpha,\sigma) \vdash (\text{Person}.\text{salary}@pre \geq 1000)
value \vee \alpha \cdot (\alpha,\sigma) \vdash (\text{Person}.\text{salary}@pre <\text{\textless} 1300)

\text{lemma} \quad (\sigma,\sigma) \vdash (\text{Person}.\text{oc1IsMaintained})
\langle proof \rangle

\text{lemma} \quad (\sigma,\sigma) \vdash ((\text{Person}.\text{oclAsType}(\text{OclAny}) .\text{oclAsType}(\text{Person}))
\doteq \text{X Person})
\langle proof \rangle

value \vee \alpha \cdot (\alpha,\sigma) \vdash (\text{Person}.\text{oc1IsTypeOf}(\text{OclAny}))
value \vee \alpha \cdot (\alpha,\sigma) \vdash (\text{X Person}.\text{oclIsKindOf}(\text{OclAny}))
value \vee \alpha \cdot (\alpha,\sigma) \vdash (\text{X Person}.\text{oclIsKindOf}(\text{OclAny}))
value \vee \alpha \cdot (\alpha,\sigma) \vdash (\text{X Person}.\text{oclAsType}(\text{OclAny}) .\text{oclIsTypeOf}(\text{OclAny}))

value \vee \alpha \cdot (\alpha,\sigma) \vdash (\text{Person}.\text{salary} \geq 1800)
value \vee \alpha \cdot (\alpha,\sigma) \vdash (\text{Person}.\text{salary}@pre \geq 1200)
\text{lemma} \quad (\sigma,\sigma) \vdash (\text{Person}.\text{oc1IsMaintained})
\langle proof \rangle

value \vee \alpha \cdot (\alpha,\sigma) \vdash (\text{Person}.\text{salary} \doteq \text{null})
value \vee \alpha \cdot (\alpha,\sigma) \vdash (\text{Person}.\text{salary}@pre \doteq \text{null})
\text{lemma} \quad (\sigma,\sigma) \vdash (\text{Person}.\text{oc1IsNew})
\langle proof \rangle
lemma \( (\sigma_1, \sigma_1') \models (X_{\text{Person}4}.\text{oclIsMaintained}) \)

(proof)

value \( \forall s_{\text{pre}}. \; (s_{\text{pre}}, \sigma_1') \models \neg v(X_{\text{Person}5}.\text{salary}) \)

value \( \forall s_{\text{post}}. \; (\sigma_1, s_{\text{post}}) \models (X_{\text{Person}5}.\text{salary} \& \neg \pre \models 3500) \)

lemma \( (\sigma_1, \sigma_1') \models (X_{\text{Person}5}.\text{oclIsDeleted}) \)

(proof)

lemma \( (\sigma_1, \sigma_1') \models (X_{\text{Person}6}.\text{oclIsMaintained}) \)

(proof)

value \( \forall s_{\text{pre}} s_{\text{post}}. \; (s_{\text{pre}}, s_{\text{post}}) \models v(X_{\text{Person}7}.\text{oclAsType}(\text{Person})) \)

lemma \( \forall s_{\text{pre}} s_{\text{post}}. \; (s_{\text{pre}}, s_{\text{post}}) \models ((X_{\text{Person}7}.\text{oclAsType}(\text{Person}) \& \text{oclAsType}(\text{OclAny})) \)

\( \models (X_{\text{Person}7}.\text{oclAsType}(\text{Person})) \)

(proof)

lemma \( (\sigma_1, \sigma_1') \models (X_{\text{Person}7}.\text{oclIsNew}) \)

(proof)

value \( \forall s_{\text{pre}} s_{\text{post}}. \; (s_{\text{pre}}, s_{\text{post}}) \models (X_{\text{Person}8} \iff X_{\text{Person}7}) \)

value \( \forall s_{\text{pre}} s_{\text{post}}. \; (s_{\text{pre}}, s_{\text{post}}) \models \neg v(X_{\text{Person}8}.\text{oclAsType}(\text{Person})) \)

value \( \forall s_{\text{pre}} s_{\text{post}}. \; (s_{\text{pre}}, s_{\text{post}}) \models (X_{\text{Person}8}.\text{oclIsTypeOf}(\text{OclAny})) \)

value \( \forall s_{\text{pre}} s_{\text{post}}. \; (s_{\text{pre}}, s_{\text{post}}) \models \neg v(X_{\text{Person}8}.\text{oclIsTypeOf}(\text{Person})) \)

value \( \forall s_{\text{pre}} s_{\text{post}}. \; (s_{\text{pre}}, s_{\text{post}}) \models \neg v(X_{\text{Person}8}.\text{oclIsKindOf}(\text{Person})) \)

value \( \forall s_{\text{pre}} s_{\text{post}}. \; (s_{\text{pre}}, s_{\text{post}}) \models (X_{\text{Person}8}.\text{oclIsKindOf}(\text{OclAny})) \)

lemma \( \sigma\text{-modifiedonly}: \; (\sigma_1, \sigma_1') \models (\text{Set}\{X_{\text{Person}1}.\text{oclAsType}(\text{OclAny}) \)

\( \cup X_{\text{Person}2}.\text{oclAsType}(\text{OclAny}) \)

\( \cup X_{\text{Person}3}.\text{oclAsType}(\text{OclAny}) \}

\( \cup X_{\text{Person}4}.\text{oclAsType}(\text{OclAny}) \}

\( \cup X_{\text{Person}5}.\text{oclAsType}(\text{OclAny}) \}

\( \cup X_{\text{Person}6}.\text{oclAsType}(\text{OclAny}) \}

\( \cup X_{\text{Person}7}.\text{oclAsType}(\text{OclAny}) \}

\( \cup X_{\text{Person}8}.\text{oclAsType}(\text{OclAny}) \}

\( \cup X_{\text{Person}9}.\text{oclAsType}(\text{OclAny}) \} \rightarrow \text{oclIsModifiedOnly}() \)

(proof)

lemma \( (\sigma_1, \sigma_1') \models (X_{\text{Person}9} \& \pre (\lambda x. \; \text{oclAsType}\text{Person} \boxdot X \; x)) \iff X_{\text{Person}9} \)

(proof)

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lemma \( (\sigma_1, \sigma_1') \models ((X_{\text{Person}} \circ \text{post} (\lambda x. [\text{OclAsType}_{\text{Person}}\cdot\emptyset x])) \triangleq X_{\text{Person}}) \) 
\langle proof \rangle

lemma \( (\sigma_1, \sigma_1') \models \left( ((X_{\text{Person}} \circ \text{oclAsType}(\text{OclAny})) \circ \text{pre} (\lambda x. [\text{OclAsType}_{\text{OclAny}}\cdot\emptyset x])) \triangleq ((X_{\text{Person}} \circ \text{oclAsType}(\text{OclAny})) \circ \text{post} (\lambda x. [\text{OclAsType}_{\text{OclAny}}\cdot\emptyset x])) \right) \) 
\langle proof \rangle

lemma perm-\( \sigma_1 : \sigma_1' = \emptyset \) heap = empty 
(oid8 \mapsto \text{in}_{\text{Person}} \text{person9}) 
(oid7 \mapsto \text{in}_{\text{OclAny}} \text{person8}) 
(oid6 \mapsto \text{in}_{\text{OclAny}} \text{person7}) 
(oid5 \mapsto \text{in}_{\text{Person}} \text{person6}) 
(*oid4*) 
(oid3 \mapsto \text{in}_{\text{Person}} \text{person4}) 
(oid2 \mapsto \text{in}_{\text{Person}} \text{person3}) 
(oid1 \mapsto \text{in}_{\text{Person}} \text{person2}) 
(oid0 \mapsto \text{in}_{\text{Person}} \text{person1}) 
, assocs2 = assocs2 \sigma_1' 
, assocs3 = assocs3 \sigma_1' 
\langle proof \rangle

declare const-ss [simp]

lemma \( \bigwedge \sigma_1 \). 
\( (\sigma_1, \sigma_1') \models (\text{Person} . \text{allInstances()} \subseteq \text{Set} \{ X_{\text{Person}1}, X_{\text{Person}2}, X_{\text{Person}9}, X_{\text{Person}4}(\ast, X_{\text{Person}5\ast}), X_{\text{Person}6}, X_{\text{Person}7} . \text{oclAsType}(\text{Person})(\ast, X_{\text{Person}8\ast}), X_{\text{Person}9} \}) \) 
\langle proof \rangle

lemma \( \bigwedge \sigma_1 \). 
\( (\sigma_1, \sigma_1') \models (\text{OclAny} . \text{allInstances()} \subseteq \text{Set} \{ X_{\text{Person}1} . \text{oclAsType}(\text{OclAny}), X_{\text{Person}2} . \text{oclAsType}(\text{OclAny}), X_{\text{Person}3} . \text{oclAsType}(\text{OclAny}), X_{\text{Person}4} . \text{oclAsType}(\text{OclAny})(\ast, X_{\text{Person}5\ast}), X_{\text{Person}6} . \text{oclAsType}(\text{OclAny}), X_{\text{Person}7}, X_{\text{Person}8}, X_{\text{Person}9} . \text{oclAsType}(\text{OclAny}) \) \) 
\langle proof \rangle

end

6.2. The Employee Analysis Model (OCL)

theory Employee-AnalysisModel-OCLPart
imports Employee-AnalysisModel-UMLPart
begin
6.2.1. Standard State Infrastructure

Ideally, these definitions are automatically generated from the class model.

6.2.2. Invariant

These recursive predicates can be defined conservatively by greatest fix-point constructions—automatically. See [4, 6] for details. For the purpose of this example, we state them as axioms here.

\[
\text{axiomatization inv-Person} :: \text{Person} \to \text{Boolean}
\]

where \( A : (\tau \models (\delta \text{self})) \mapsto (\tau \models \text{inv-Person(self)}) \).

\[
(\tau \models \text{inv-Person(self)}) = \bigvee (\tau \models (\text{self.boss} = \text{null})) \land (\tau \models ((\text{self.salary} \leq (\text{self.boss.salary}))) \land (\tau \models \text{inv-Person(self.boss)}))))
\]

\[
\text{axiomatization inv-Person-at-pre} :: \text{Person} \to \text{Boolean}
\]

where \( B : (\tau \models (\delta \text{self})) \mapsto (\tau \models \text{inv-Person-at-pre(self)}) \).

\[
(\tau \models \text{inv-Person-at-pre(self)}) = \bigvee (\tau \models (\text{self.boss<>null})) \land (\tau \models ((\text{self.boss.salary} \leq \text{self.salary}))) \land (\tau \models \text{inv-Person-at-pre(self.boss<>null)})))
\]

A very first attempt to characterize the axiomatization by an inductive definition - this can not be the last word since too weak (should be equality!)

\[
\text{coinductive inv} :: \text{Person} \to (\exists \text{st} \Rightarrow \text{bool}} \text{ where}
\]

\[
(\tau \models (\delta \text{self})) \mapsto ((\tau \models (\text{self.boss=} \text{null})) \land (\tau \models (\text{inv\{self.boss\}})))
\]

6.2.3. The Contract of a Recursive Query

The original specification of a recursive query:

\[
\text{context Person :: contents() : Set(Integer)}
\]

\[
\text{post: result = if self.boss = null then Set{i} else self.boss.contents()->including(i) endif}
\]

\[
\text{consts dot-contents :: Person \to Set-Integer ((I(-).contents') 50)}
\]

\[
\text{axiomatization where dot-contents-def:}
\]

\[
(\tau \models ((\text{self.contents()}) \triangleq \text{result})) = (\text{if} (\delta \text{self}) \tau = \text{true} \tau \text{ then} ((\tau \models \text{true}) \land
\]

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\[(\tau \models (\text{result} \triangleq \text{if} \ (\text{self} \ . \ \text{boss} \ \neq \ \text{null}) \ \text{then} \ \{\text{self} \ . \ \text{salary}\}) \ \text{else} \ (\text{self} \ . \ \text{boss} \ . \ \text{contents()} - \ \text{including}(\text{self} \ . \ \text{salary})) \ \text{endif}))\]

\[
\text{else } \tau \models \text{result} \triangleq \text{invalid})
\]

\text{consts} \ \text{dot-contents-AT-pre} :: \text{Person} \Rightarrow \text{Set-Integer} \ ((1 \ . \ \text{contents}@\text{pre'}(\text{'}) \ 50)

\text{axiomatization where} \ \text{dot-contents-AT-pre-def}:
\[\tau \models (\text{self} \ . \ \text{contents}@\text{pre}()) \triangleq \text{result} = \]
\quad \text{if} \ (\delta \ \text{self}) \ \tau = \text{true} \ \tau
\quad \text{then} \ \tau \models \text{true} \ \text{\&} \quad (\ast \ \text{pre} \ast)
\quad \tau \models (\text{result} \triangleq \text{if} \ (\text{self} \ . \ \text{boss}@\text{pre} \ \neq \ \text{null} \ \text{\&} \ \text{post} \ \ast)
\quad \text{then} \ \{\text{self} \ . \ \text{salary}@\text{pre}\}
\quad \text{else} \ (\text{self} \ . \ \text{boss}@\text{pre} \ . \ \text{contents}@\text{pre}()) - \ \text{including}(\text{self} \ . \ \text{salary}@\text{pre})
\quad \text{endif}
\quad \text{else} \ \tau \models \text{result} \triangleq \text{invalid})
\]

These @\text{pre} variants on methods are only available on queries, i.e., operations without side-effect.

6.2.4. The Contract of a Method

The specification in high-level OCL input syntax reads as follows:

```plaintext
context Person::insert(x: Integer)
post: contents():Set(Integer)
contents() = contents@pre()->including(x)

consts dot-insert :: Person \Rightarrow \text{Integer} \Rightarrow \text{Void} \ ((1 \ . \ \text{insert'}(\text{'}) \ 50)

axiomatization where dot-insert-def:
\[\tau \models ((\text{self} \ . \ \text{insert}(x) \ \triangleq \text{result})) = \]
\quad (if \ (\delta \ \text{self}) \ \tau = \text{true} \ \text{\&} \ (v \ x) \ \tau = \text{true} \ \tau
\quad \text{then} \ \tau \models \text{true} \ \text{\&}
\quad \tau \models ((\text{self} \ . \ \text{contents()} \ \triangleq (\text{self} \ . \ \text{contents}@\text{pre}()) - \ \text{including}(x))
\quad \text{else} \ \tau \models ((\text{self} \ . \ \text{insert}(x) \ \triangleq \text{invalid}))
```

end
7. The Employee Design Model

7.1. The Employee Design Model (UML)

theory 
   Employee-DesignModel-UMLPart
imports 
   ../OCL-main
begin

7.1.1. Introduction

For certain concepts like classes and class-types, only a generic definition for its resulting semantics can be given. Generic means, there is a function outside HOL that “compiles” a concrete, closed-world class diagram into a “theory” of this data model, consisting of a bunch of definitions for classes, accessors, method, casts, and tests for actual types, as well as proofs for the fundamental properties of these operations in this concrete data model.

Such generic function or “compiler” can be implemented in Isabelle on the ML level. This has been done, for a semantics following the open-world assumption, for UML 2.0 in [4, 7]. In this paper, we follow another approach for UML 2.4: we define the concepts of the compilation informally, and present a concrete example which is verified in Isabelle/HOL.

Outlining the Example

We are presenting here a “design-model” of the (slightly modified) example Figure 7.3, page 20 of the OCL standard [33]. To be precise, this theory contains the formalization of the data-part covered by the UML class model (see Figure 7.1):

This means that the association (attached to the association class EmployeeRanking) with the association ends boss and employees is implemented by the attribute boss and the operation employees (to be discussed in the OCL part captured by the subsequent theory).

7.1.2. Example Data-Universe and its Infrastructure

Ideally, the following is generated automatically from a UML class model.
Our data universe consists in the concrete class diagram just of node’s, and implicitly
of the class object. Each class implies the existence of a class type defined for the
corresponding object representations as follows:

\[
\text{datatype type}_{\text{Person}} = \text{mk}_{\text{Person}} \text{ oid}
\]
\[
\begin{align*}
&\quad \text{int option} \\
&\quad \text{oid option}
\end{align*}
\]

\[
\text{datatype type}_{\text{OclAny}} = \text{mk}_{\text{OclAny}} \text{ oid}
\]
\[
\begin{align*}
&\quad (\text{int option} \times \text{oid option}) \text{ option}
\end{align*}
\]

Now, we construct a concrete “universe of OclAny types” by injection into a sum type
containing the class types. This type of OclAny will be used as instance for all respective
type-variables.

\[
\text{datatype } \mathbb{A} = \text{in}_{\text{Person}} \text{ type}_{\text{Person}} | \text{in}_{\text{OclAny}} \text{ type}_{\text{OclAny}}
\]

Having fixed the object universe, we can introduce type synonyms that exactly corre-
spond to OCL types. Again, we exploit that our representation of OCL is a “shallow em-
bedding” with a one-to-one correspondance of OCL-types to types of the meta-language
HOL.

\[
\begin{align*}
\text{type-synonym } \text{Boolean} & = \mathbb{A} \text{ Boolean} \\
\text{type-synonym } \text{Integer} & = \mathbb{A} \text{ Integer} \\
\text{type-synonym } \text{Void} & = \mathbb{A} \text{ Void} \\
\text{type-synonym } \text{OclAny} & = (\mathbb{A}, \text{type}_{\text{OclAny}} \text{ option option}) \text{ val} \\
\text{type-synonym } \text{Person} & = (\mathbb{A}, \text{type}_{\text{Person}} \text{ option option}) \text{ val} \\
\text{type-synonym } \text{Set-Integer} & = (\mathbb{A}, \text{int option option}) \text{ Set} \\
\text{type-synonym } \text{Set-Person} & = (\mathbb{A}, \text{type}_{\text{Person}} \text{ option option}) \text{ Set}
\end{align*}
\]

Just a little check:

\[
\text{typ } \text{Boolean}
\]

To reuse key-elements of the library like referential equality, we have to show that the
object universe belongs to the type class “oclany,” i.e., each class type has to provide a
function \text{oid-of} yielding the object id (oid) of the object.
instantiation type Person :: object
begin
definition oid-of-typePerson-def: oid-of x = (case x of mkPerson oid - \Rightarrow oid)
instance ⟨proof⟩
end

instantiation type OclAny :: object
begin
definition oid-of-typeOclAny-def: oid-of x = (case x of mkOclAny oid - \Rightarrow oid)
instance ⟨proof⟩
end

instantiation A :: object
begin
definition oid-of-A-def: oid-of x = (case x of
    inPerson person \Rightarrow oid-of person
    \mid inOclAny oclany \Rightarrow oid-of oclany)
instance ⟨proof⟩
end

7.1.3. Instantiation of the Generic Strict Equality

We instantiate the referential equality on Person and OclAny

defs(overloaded) StrictRefEqObject*Person : (x::Person) \equiv y \equiv StrictRefEqObject*OclAny
defs(overloaded) StrictRefEqObject*OclAny : (x::OclAny) \equiv y \equiv StrictRefEqObject*Person

lemmas
cp-StrictRefEqObject[of x::Person y::Person \tau,
simplified StrictRefEqObject*Person[symmetric]]
cp-intro(9) [of P::Person \Rightarrow PersonQ::Person \Rightarrow Person,
simplified StrictRefEqObject*Person[symmetric] ]
StrictRefEqObject-def[of x::Person y::Person,
simplified StrictRefEqObject*Person[symmetric]]
StrictRefEqObject-defargs[of x::Person y::Person,
simplified StrictRefEqObject*Person[symmetric]]
StrictRefEqObject-strict1[of x::Person,
simplified StrictRefEqObject*Person[symmetric]]
StrictRefEqObject-strict2[of x::Person,
simplified StrictRefEqObject*Person[symmetric]]

For each Class C, we will have a casting operation .oclAsType(C), a test on the actual type .oclIsTypeOf(C) as well as its relaxed form .oclIsKindOf(C) (corresponding exactly to Java’s instanceof-operator.

Thus, since we have two class-types in our concrete class hierarchy, we have two operations to declare and to provide two overloading definitions for the two static types.
7.1.4. OclAsType

Definition

consts OclAsTypeOclAny :: \( \alpha \Rightarrow \text{OclAny} ((-) . \text{oclAsType} (\text{OclAny})) \)
consts OclAsTypePerson :: \( \alpha \Rightarrow \text{Person} ((-) . \text{oclAsType} (\text{Person})) \)

definition OclAsTypeOclAny-\( \mathfrak{A} \)= (\( \lambda u \). case u of inOclAny a \( \Rightarrow a \)

| inPerson (mkPerson oid a b) \( \Rightarrow mkOclAny oid [(a,b)] \))

lemma OclAsTypeOclAny-\( \mathfrak{A} \)-some: OclAsTypeOclAny-\( \mathfrak{A} \) \( x \neq \text{None} \)

⟨proof⟩
defs (overloaded) OclAsTypeOclAny-OclAny:
(\( X :: \text{OclAny} \) . oclAsType(\( \text{OclAny} \)) \( \equiv X \)

defs (overloaded) OclAsTypeOclAny-Person:
(\( X :: \text{Person} \) . oclAsType(\( \text{OclAny} \)) \( \equiv \)

(\( \lambda \tau . \text{case X } \tau \text{ of} \)
| \( \bot \Rightarrow \text{invalid } \tau \)
| \( [\bot] \Rightarrow \text{null } \tau \)
| \( [\text{mkPerson oid a b}] \Rightarrow [\text{mkOclAny oid [(a,b)] }] \))

definition OclAsTypePerson-\( \mathfrak{A} \)= (\( \lambda u . \text{case u of inPerson p } \Rightarrow [p] \)

| inOclAny (mkOclAny oid [(a,b)]) \( \Rightarrow [\text{mkPerson oid a b}] \)
| - \( \Rightarrow \text{None} \))
defs (overloaded) OclAsTypePerson-OclAny:
(\( X :: \text{OclAny} \) . oclAsType(\( \text{Person} \)) \( \equiv \)

(\( \lambda \tau . \text{case X } \tau \text{ of} \)
| \( \bot \Rightarrow \text{invalid } \tau \)
| \( [\bot] \Rightarrow \text{null } \tau \)
| \( [\text{mkOclAny oid } \bot] \Rightarrow \text{invalid } \tau \) (* down-cast exception *)
| \( [\text{mkOclAny oid (a,b)}] \Rightarrow [\text{mkPerson oid a b}] \))
defs (overloaded) OclAsTypePerson-Person:
(\( X :: \text{Person} \) . oclAsType(\( \text{Person} \)) \( \equiv X \)

lemmas [simp] =
OclAsTypeOclAny-OclAny
OclAsTypePerson-Person

Context Passing

lemma cp-OclAsTypeOclAny-Person-Person: cp P \( \Rightarrow cp (\lambda X . (P (X :: \text{Person}) :: \text{Person}) . oclAsType(\( \text{OclAny} \))) \)
⟨proof⟩
lemma cp-OclAsTypeOclAny-OclAny-OclAny: cp P \( \Rightarrow cp (\lambda X . (P (X :: \text{OclAny}) :: \text{OclAny}) . oclAsType(\( \text{OclAny} \))) \)
Lemma cp-OclAsType_Person-Person-Person
\[
\begin{align*}
\text{cp P} & 
\implies cp(\lambda X. (P (X::\text{Person})::\text{Person})) \\
.\text{oclAsType(\text{Person})}
\end{align*}
\]

Lemma cp-OclAsType_Person-OclAny-OclAny
\[
\begin{align*}
\text{cp P} & 
\implies cp(\lambda X. (P (X::\text{OclAny})::\text{OclAny})) \\
.\text{oclAsType(\text{Person})}
\end{align*}
\]

Lemma cp-OclAsType_OclAny-Person-OclAny
\[
\begin{align*}
\text{cp P} & 
\implies cp(\lambda X. (P (X::\text{Person})::\text{OclAny})) \\
.\text{oclAsType(\text{OclAny})}
\end{align*}
\]

Lemma cp-OclAsType_OclAny-Person-Person
\[
\begin{align*}
\text{cp P} & 
\implies cp(\lambda X. (P (X::\text{OclAny})::\text{Person})) \\
.\text{oclAsType(\text{OclAny})}
\end{align*}
\]

Lemma cp-OclAsType_Person-Person-OclAny
\[
\begin{align*}
\text{cp P} & 
\implies cp(\lambda X. (P (X::\text{Person})::\text{OclAny})) \\
.\text{oclAsType(\text{Person})}
\end{align*}
\]

Lemma cp-OclAsType_Person-OclAny-Person
\[
\begin{align*}
\text{cp P} & 
\implies cp(\lambda X. (P (X::\text{OclAny})::\text{Person})) \\
.\text{oclAsType(\text{Person})}
\end{align*}
\]

Lemma cp-OclAsType_OclAny-OclAny-Person
\[
\begin{align*}
\text{cp P} & 
\implies cp(\lambda X. (P (X::\text{Person})::\text{OclAny})) \\
.\text{oclAsType(\text{OclAny})}
\end{align*}
\]

Lemma cp-OclAsType_OclAny-Person-OclAny
\[
\begin{align*}
\text{cp P} & 
\implies cp(\lambda X. (P (X::\text{Person})::\text{OclAny})) \\
.\text{oclAsType(\text{OclAny})}
\end{align*}
\]

Lemma cp-OclAsType_Person-OclAny-Person
\[
\begin{align*}
\text{cp P} & 
\implies cp(\lambda X. (P (X::\text{Person})::\text{OclAny})) \\
.\text{oclAsType(\text{Person})}
\end{align*}
\]

Execution with Invalid or Null as Argument

Lemma OclAsType_OclAny-OclAny-strict : (invalid::OclAny) .oclAsType(OclAny) = invalid

Lemma OclAsType_OclAny-OclAny-nullstrict : (null::OclAny) .oclAsType(OclAny) = null

Lemma OclAsType_OclAny-Person-strict [simp] : (invalid::Person) .oclAsType(OclAny) = invalid

Lemma OclAsType_OclAny-Person-nullstrict [simp] : (null::Person) .oclAsType(OclAny) = null

Lemma OclAsType_Person-OclAny-strict [simp] : (invalid::OclAny) .oclAsType(Person) = invalid

Lemma OclAsType_Person-OclAny-nullstrict [simp] : (null::OclAny) .oclAsType(Person) = invalid
lemma OclAsType\_Person-\_OclAny-nullstrict[simp] : (null::OclAny) .oclAsType(Person) = null ⟨proof⟩

lemma OclAsType\_Person-\_Person-strict : (invalid::Person) .oclAsType(Person) = invalid ⟨proof⟩
lemma OclAsType\_Person-\_Person-nullstrict : (null::Person) .oclAsType(Person) = null ⟨proof⟩

7.1.5. OclIsTypeOf

Definition

consts OclIsTypeOf\_OclAny :: 'α ⇒ Boolean ((-).oclIsTypeOf'(OclAny'))
consts OclIsTypeOf\_Person :: 'α ⇒ Boolean ((-).oclIsTypeOf'(Person'))

defs (overloaded) OclIsTypeOf\_OclAny-\_OclAny:
(X::OclAny) .oclIsTypeOf(OclAny) ≡
(λτ. case X τ of
  ⊥ ⇒ invalid τ
  | [⊥] ⇒ true τ (* invalid ?? *)
  | [mkOclAny oid ⊥ ] ⇒ true τ
  | [mkOclAny oid [- ] ] ⇒ false τ)

defs (overloaded) OclIsTypeOf\_OclAny-\_Person:
(X::Person) .oclIsTypeOf(OclAny) ≡
(λτ. case X τ of
  ⊥ ⇒ invalid τ
  | [⊥] ⇒ true τ (* invalid ?? *)
  | [- ] ⇒ false τ)

defs (overloaded) OclIsTypeOf\_Person-\_OclAny:
(X::OclAny) .oclIsTypeOf(Person) ≡
(λτ. case X τ of
  ⊥ ⇒ invalid τ
  | [⊥] ⇒ true τ
  | [mkOclAny oid ⊥ ] ⇒ false τ
  | [mkOclAny oid [- ] ] ⇒ true τ)

defs (overloaded) OclIsTypeOf\_Person-\_Person:
(X::Person) .oclIsTypeOf(Person) ≡
(λτ. case X τ of
  ⊥ ⇒ invalid τ
  | - ⇒ true τ)

Context Passing

lemma cp-OclIsTypeOf\_OclAny-\_Person-\_Person: cp P ⇒
  cp(λX.(P(X::Person)::Person).oclIsTypeOf(OclAny))
  ⟨proof⟩
lemma \( cp \cdot \text{OclIsTypeOf}_{\text{OclAny}, \text{OclAny}, \text{OclAny}}: \) 
  \( cp(\lambda X. (P(X::\text{OclAny})::\text{OclAny}). \text{oclIsTypeOf}(\text{OclAny})) \)
  \( \langle \text{proof} \rangle \)

lemma \( cp \cdot \text{OclIsTypeOf}_{\text{Person}, \text{Person}, \text{Person}}: \) 
  \( cp(\lambda X. (P(X::\text{Person})::\text{Person}). \text{oclIsTypeOf}(\text{Person})) \)
  \( \langle \text{proof} \rangle \)

lemma \( cp \cdot \text{OclIsTypeOf}_{\text{Person}, \text{OclAny}, \text{OclAny}}: \) 
  \( cp(\lambda X. (P(X::\text{OclAny})::\text{OclAny}). \text{oclIsTypeOf}(\text{Person})) \)
  \( \langle \text{proof} \rangle \)

lemma \( cp \cdot \text{OclIsTypeOf}_{\text{Person}, \text{Person}, \text{OclAny}}: \) 
  \( cp(\lambda X. (P(X::\text{OclAny})::\text{OclAny}). \text{oclIsTypeOf}(\text{Person})) \)
  \( \langle \text{proof} \rangle \)

lemma \( cp \cdot \text{OclIsTypeOf}_{\text{Person}, \text{OclAny}, \text{OclAny}}: \) 
  \( cp(\lambda X. (P(X::\text{Person})::\text{OclAny}). \text{oclIsTypeOf}(\text{Person})) \)
  \( \langle \text{proof} \rangle \)

lemma \( cp \cdot \text{OclIsTypeOf}_{\text{Person}, \text{OclAny}, \text{OclAny}}: \) 
  \( cp(\lambda X. (P(X::\text{Person})::\text{OclAny}). \text{oclIsTypeOf}(\text{Person})) \)
  \( \langle \text{proof} \rangle \)

lemmas \([\text{simpl}] = \) 
  \( cp \cdot \text{OclIsTypeOf}_{\text{OclAny}, \text{OclAny}, \text{Person}} \) 
  \( cp \cdot \text{OclIsTypeOf}_{\text{OclAny}, \text{OclAny}, \text{OclAny}} \) 
  \( cp \cdot \text{OclIsTypeOf}_{\text{Person}, \text{Person}, \text{Person}} \) 
  \( cp \cdot \text{OclIsTypeOf}_{\text{Person}, \text{OclAny}, \text{OclAny}} \) 

 Execution with Invalid or Null as Argument

lemma \( \text{OclIsTypeOf}_{\text{OclAny}, \text{OclAny}, \text{strict1}}[\text{simpl}]: \) 
  \( (\text{invalid}::\text{OclAny}). \text{oclIsTypeOf}(\text{OclAny}) = \text{invalid} \) 
  \( \langle \text{proof} \rangle \)

lemma \( \text{OclIsTypeOf}_{\text{OclAny}, \text{OclAny}, \text{strict2}}[\text{simpl}]: \) 
  \( (\text{null}::\text{OclAny}). \text{oclIsTypeOf}(\text{OclAny}) = \text{true} \) 
  \( \langle \text{proof} \rangle \)

lemma \( \text{OclIsTypeOf}_{\text{OclAny}, \text{Person}, \text{strict1}}[\text{simpl}]: \) 
  \( (\text{invalid}::\text{Person}). \text{oclIsTypeOf}(\text{OclAny}) = \text{invalid} \) 
  \( \langle \text{proof} \rangle \)

lemma \( \text{OclIsTypeOf}_{\text{OclAny}, \text{Person}, \text{strict2}}[\text{simpl}]: \) 
  \( (\text{null}::\text{Person}). \text{oclIsTypeOf}(\text{OclAny}) = \text{true} \) 
  \( \langle \text{proof} \rangle \)
lemma \textit{OclIsTypeOf\_Person-\_\_OclAny-strict1}[simp]:
\begin{align*}
\text{(invalid::\_\_OclAny) .oclIsTypeOf(\_\_Person) = invalid}
\end{align*}
\langle \text{proof} \rangle

lemma \textit{OclIsTypeOf\_Person-\_\_OclAny-strict2}[simp]:
\begin{align*}
\text{(null::\_\_OclAny) .oclIsTypeOf(\_\_Person) = true}
\end{align*}
\langle \text{proof} \rangle

lemma \textit{OclIsTypeOf\_Person-Person-\_\_strict1}[simp]:
\begin{align*}
\text{(invalid::\_\_Person) .oclIsTypeOf(\_\_Person) = invalid}
\end{align*}
\langle \text{proof} \rangle

lemma \textit{OclIsTypeOf\_Person-Person-\_\_strict2}[simp]:
\begin{align*}
\text{(null::\_\_Person) .oclIsTypeOf(\_\_Person) = true}
\end{align*}
\langle \text{proof} \rangle

\textbf{Up Down Casting}

lemma \textit{actualType-larger-staticType}:
\begin{align*}
\text{assumes } & \text{isdef: } \tau \models (\delta \text{ } X) \\
\text{shows } & \tau \models (\text{X::Person).oclIsTypeOf(OclAny) } \triangleq \text{ false}
\end{align*}
\langle \text{proof} \rangle

lemma \textit{down-cast-type}:
\begin{align*}
\text{assumes } & \text{isOclAny: } \tau \models (X::\text{OclAny}) .oclIsTypeOf(\text{OclAny}) \\
\text{and } & \text{non-null: } \tau \models (\delta \text{ } X) \\
\text{shows } & \tau \models (X .oclAsType(\text{Person})) \triangleq \text{ invalid}
\end{align*}
\langle \text{proof} \rangle

lemma \textit{down-cast-type\_'}:
\begin{align*}
\text{assumes } & \text{isOclAny: } \tau \models (X::\text{OclAny}) .oclIsTypeOf(\text{OclAny}) \\
\text{and } & \text{non-null: } \tau \models (\delta \text{ } X) \\
\text{shows } & \tau \models \text{not (v (X .oclAsType(\text{Person})))}
\end{align*}
\langle \text{proof} \rangle

lemma \textit{up-down-cast} :
\begin{align*}
\text{assumes } & \text{isdef: } \tau \models (\delta \text{ } X) \\
\text{shows } & \tau \models ((\text{X::Person).oclAsType(OclAny) .oclAsType(Person) } \triangleq \text{ X})
\end{align*}
\langle \text{proof} \rangle

lemma \textit{up-down-cast-Person-\_\_OclAny-Person} [simp]:
\begin{align*}
\text{shows } ((\text{X::Person).oclAsType(OclAny) .oclAsType(Person) = X})
\end{align*}
\langle \text{proof} \rangle

lemma \textit{up-down-cast-Person-\_\_OclAny-Person\_'} : \textbf{assumes } \tau \models \nu \text{ } X
\begin{align*}
\text{shows } & \tau \models ((\text{X::Person).oclAsType(OclAny) .oclAsType(Person) } \triangleright \text{ X})
\end{align*}
\langle \text{proof} \rangle

lemma \textit{up-down-cast-Person-\_\_OclAny-Person\_\_} : \textbf{assumes } \tau \models \nu \text{ } (X::\text{Person})
\begin{align*}
\text{shows } & \tau \models (\text{X .oclIsTypeOf(Person) implies (X .oclAsType(OclAny) .oclAsType(Person) } \triangleright \text{ X})
\end{align*}
7.1.6. OclIsKindOf

Definition

consts OclIsKindOf : 'α ⇒ Boolean
(OclAny) ≡
(λτ. case X τ of
   ⊥ ⇒ invalid τ
   | - ⇒ true τ)

consts OclIsKindOf Person : 'α ⇒ Boolean
(Person) ≡
(λτ. case X τ of
   ⊥ ⇒ invalid τ
   | - ⇒ true τ)

defs (overloaded) OclIsKindOf OclAny-OclAny:
   (X::OclAny) .oclIsKindOf(OclAny) ≡
   (λτ. case X τ of
      ⊥ ⇒ invalid τ
      | - ⇒ true τ)

defs (overloaded) OclIsKindOf OclAny-Person:
   (X::Person) .oclIsKindOf(Person) ≡
   (λτ. case X τ of
      ⊥ ⇒ invalid τ
      | - ⇒ true τ)

defs (overloaded) OclIsKindOf Person-OclAny:
   (X::OclAny) .oclIsKindOf(Person) ≡
   (λτ. case X τ of
      ⊥ ⇒ invalid τ
      | [mkOclAny oid ⊥] ⇒ false τ
      | [mkOclAny oid [·]] ⇒ true τ)

defs (overloaded) OclIsKindOf Person-Person:
   (X::Person) .oclIsKindOf(Person) ≡
   (λτ. case X τ of
      ⊥ ⇒ invalid τ
      | - ⇒ true τ)

Context Passing

lemma cp-OclIsKindOf OclAny-Person-Person:
   cp(λX.(P(X::Person)::Person).oclIsKindOf(OclAny))
   (proof)

lemma cp-OclIsKindOf OclAny-OclAny:
   cp(λX.(P(X::OclAny)::OclAny).oclIsKindOf(OclAny))
   (proof)

lemma cp-OclIsKindOf Person-Person-Person:
   cp(λX.(P(X::Person)::Person).oclIsKindOf(Person))
   (proof)

lemma cp-OclIsKindOf Person-OclAny-OclAny:
   cp(λX.(P(X::OclAny)::OclAny).oclIsKindOf(Person))
   (proof)
lemma \[cp-\text{OclIsKindOf}\_\text{OclAny-Person-OclAny}: \quad \text{cp}(\lambda X.(P(X::\text{Person})::\text{OclAny}).\text{oclIsKindOf}(\text{OclAny})) \quad \langle \text{proof} \rangle \]

lemma \[cp-\text{OclIsKindOf}\_\text{OclAny-OclAny-Person}: \quad \text{cp}(\lambda X.(P(X::\text{OclAny})::\text{Person}).\text{oclIsKindOf}(\text{OclAny})) \quad \langle \text{proof} \rangle \]

lemma \[cp-\text{OclIsKindOf}\_\text{Person-Person-OclAny}: \quad \text{cp}(\lambda X.(P(X::\text{Person})::\text{OclAny}).\text{oclIsKindOf}(\text{Person})) \quad \langle \text{proof} \rangle \]

lemma \[cp-\text{OclIsKindOf}\_\text{Person-OclAny-Person}: \quad \text{cp}(\lambda X.(P(X::\text{OclAny})::\text{Person}).\text{oclIsKindOf}(\text{Person})) \quad \langle \text{proof} \rangle \]

lemmas \[\text{[simp]} = \]

\[ \text{cp-\text{OclIsKindOf}\_\text{OclAny-Person-Person}} \]
\[ \text{cp-\text{OclIsKindOf}\_\text{OclAny-OclAny-OclAny}} \]
\[ \text{cp-\text{OclIsKindOf}\_\text{Person-Person-Person}} \]
\[ \text{cp-\text{OclIsKindOf}\_\text{Person-OclAny-OclAny}} \]
\[ \text{cp-\text{OclIsKindOf}\_\text{OclAny-Person-Person}} \]
\[ \text{cp-\text{OclIsKindOf}\_\text{OclAny-Person-Person}} \]
\[ \text{cp-\text{OclIsKindOf}\_\text{Person-Person-OclAny}} \]
\[ \text{cp-\text{OclIsKindOf}\_\text{Person-Person-OclAny}} \]

\[\text{Execution with Invalid or Null as Argument}\]

lemma \[\text{OclIsKindOf}\_\text{OclAny-OclAny-strict1 [simp]} : \quad (\text{invalid}::\text{OclAny}).\text{oclIsKindOf}(\text{OclAny}) = \text{invalid} \quad \langle \text{proof} \rangle \]

lemma \[\text{OclIsKindOf}\_\text{OclAny-OclAny-strict2 [simp]} : \quad (\text{null}::\text{OclAny}).\text{oclIsKindOf}(\text{OclAny}) = \text{true} \quad \langle \text{proof} \rangle \]

lemma \[\text{OclIsKindOf}\_\text{OclAny-Person-strict1 [simp]} : \quad (\text{invalid}::\text{Person}).\text{oclIsKindOf}(\text{OclAny}) = \text{invalid} \quad \langle \text{proof} \rangle \]

lemma \[\text{OclIsKindOf}\_\text{OclAny-Person-strict2 [simp]} : \quad (\text{null}::\text{Person}).\text{oclIsKindOf}(\text{OclAny}) = \text{true} \quad \langle \text{proof} \rangle \]

lemma \[\text{OclIsKindOf}\_\text{Person-OclAny-strict1 [simp]} : \quad (\text{invalid}::\text{OclAny}).\text{oclIsKindOf}(\text{Person}) = \text{invalid} \quad \langle \text{proof} \rangle \]

lemma \[\text{OclIsKindOf}\_\text{Person-OclAny-strict2 [simp]} : \quad (\text{null}::\text{OclAny}).\text{oclIsKindOf}(\text{Person}) = \text{true} \quad \langle \text{proof} \rangle \]

lemma \[\text{OclIsKindOf}\_\text{Person-Person-strict1 [simp]} : \quad (\text{invalid}::\text{Person}).\text{oclIsKindOf}(\text{Person}) = \text{invalid} \quad \langle \text{proof} \rangle \]
lemma OclIsKindOf_Person-Person-strict2[simp]: (null::Person) .oclIsKindOf(Person) = true

(proof)

Up Down Casting

lemma actualKind-larger-staticKind:
assumes isdef: τ |= (δ X)
sows τ |= (X::Person) .oclIsKindOf(OclAny) ≜ true

(proof)

lemma down-cast-kind:
assumes isOclAny: ¬ τ |= (X::OclAny) .oclIsKindOf(Person)
and non-null: τ |= (δ X)
sows τ |= (X .oclAsType(Person)) ≜ invalid

(proof)

7.1.7. OclAllInstances

To denote OCL-types occuring in OCL expressions syntactically—as, for example, as “argument” of oclAllInstances()—we use the inverses of the injection functions into the object universes; we show that this is sufficient “characterization.”

definition Person ≡ OclAsType_Person-\mathcal{A}
definition OclAny ≡ OclAsType_OclAny-\mathcal{A}
lemmas [simp] = Person-def OclAny-def

lemma OclAllInstances-genericOclAny-exec: OclAllInstances-generic pre-post OclAny =
(λτ. Abs-Set-0 [[ Some ' OclAny ' ran (heap (pre-post τ)) ]])

(proof)

lemma OclAllInstances-at-postOclAny-exec: OclAny.allInstances() =
(λτ. Abs-Set-0 [[ Some ' OclAny ' ran (heap (snd τ)) ]])

(proof)

lemma OclAllInstances-at-preOclAny-exec: OclAny.allInstances@pre() =
(λτ. Abs-Set-0 [[ Some ' OclAny ' ran (heap (fst τ)) ]])

(proof)

OclIsTypeOf

lemma OclAny-allInstances-generic-oclIsTypeOfOclAnyI:
assumes [simp]: ∀x. pre-post (x, x) = x
shows ∃τ. (τ |= (OclAllInstances-generic pre-post OclAny) → forAll(X|X .oclIsTypeOf(OclAny))))

(proof)

lemma OclAny-allInstances-at-post-oclIsTypeOfOclAnyI:
\[ \exists \tau. \ (\tau \models (\text{OclAny} \ . \text{allInstances}() \rightarrow \forall X | X \ . \text{oclIsTypeOf}(\text{OclAny}))) \]

\text{lemma OclAny-allInstances-at-pre-oclIsTypeOf OclAny}^{1}: 
\[ \exists \tau. \ (\tau \models (\text{OclAny} \ . \text{allInstances}@\text{pre}() \rightarrow \forall X | X \ . \text{oclIsTypeOf}(\text{OclAny}))) \]

\text{lemma OclAny-allInstances-generic-oclIsTypeOf OclAny}^{2}: 
\text{assumes} \ [\text{simp}]: \ \forall x. \ \text{pre} \ (x, x) = x 
\text{shows} \ \exists \tau. \ (\tau \models \neg ((\text{OclAllInstances-generic pre-post OclAny}) \rightarrow \forall X | X \ . \text{oclIsTypeOf}(\text{OclAny}))) 

\text{lemma OclAny-allInstances-at-post-oclIsTypeOf OclAny}^{2}: 
\[ \exists \tau. \ (\tau \models \neg (\text{OclAny} \ . \text{allInstances}() \rightarrow \forall X | X \ . \text{oclIsTypeOf}(\text{OclAny}))) \]

\text{lemma Person-allInstances-generic-oclIsTypeOf OclAny}: 
\tau \models ((\text{OclAllInstances-generic pre-post Person}) \rightarrow \forall X | X \ . \text{oclIsTypeOf}(\text{Person}))) 

\text{lemma Person-allInstances-at-post-oclIsTypeOf Person}: 
\tau \models (\text{Person} \ . \text{allInstances}() \rightarrow \forall X | X \ . \text{oclIsTypeOf}(\text{Person}))) 

\text{lemma Person-allInstances-at-pre-oclIsTypeOf Person}: 
\tau \models (\text{Person} \ . \text{allInstances}@\text{pre}() \rightarrow \forall X | X \ . \text{oclIsTypeOf}(\text{Person}))) 

\text{OclIsKindOf}

\text{lemma OclAny-allInstances-generic-oclIsKindOf OclAny}: 
\tau \models ((\text{OclAllInstances-generic pre-post OclAny}) \rightarrow \forall X | X \ . \text{oclIsKindOf}(\text{OclAny}))) 

\text{lemma OclAny-allInstances-at-post-oclIsKindOf OclAny}: 
\tau \models (\text{OclAny} \ . \text{allInstances}() \rightarrow \forall X | X \ . \text{oclIsKindOf}(\text{OclAny}))) 

\text{lemma OclAny-allInstances-at-pre-oclIsKindOf OclAny}: 
\tau \models (\text{OclAny} \ . \text{allInstances}@\text{pre}() \rightarrow \forall X | X \ . \text{oclIsKindOf}(\text{OclAny}))) 

\text{lemma Person-allInstances-generic-oclIsKindOf OclAny}: 
\tau \models ((\text{OclAllInstances-generic pre-post Person}) \rightarrow \forall X | X \ . \text{oclIsKindOf}(\text{OclAny}))) 

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lemma Person-allInstances-at-post-oclIsKindOf_OclAny:
\[ \tau \models (\text{Person}.\text{allInstances}() - \forall X. X.\text{oclIsKindOf}(\text{OclAny})) \]

lemma Person-allInstances-at-pre-oclIsKindOf_OclAny:
\[ \tau \models (\text{Person}.\text{allInstances}@\text{pre}() - \forall X. X.\text{oclIsKindOf}(\text{OclAny})) \]

lemma Person-allInstances-generic-oclIsKindOf_Person:
\[ \tau \models ((\text{OclAllInstances-generic pre-post Person}) - \forall X. X.\text{oclIsKindOf}(\text{Person})) \]

lemma Person-allInstances-at-post-oclIsKindOf_Person:
\[ \tau \models (\text{Person}.\text{allInstances}() - \forall X. X.\text{oclIsKindOf}(\text{Person})) \]

lemma Person-allInstances-at-pre-oclIsKindOf_Person:
\[ \tau \models (\text{Person}.\text{allInstances}@\text{pre}() - \forall X. X.\text{oclIsKindOf}(\text{Person})) \]

7.1.8. The Accessors (any, boss, salary)

Should be generated entirely from a class-diagram.

Definition

definition eval-extract :: (\mathcal{F}.('a::object) \text{ option option}) \text{ val} \\

\Rightarrow (\text{oid} \Rightarrow (\mathcal{F}.,'c::null \text{ val})) \\
\Rightarrow (\mathcal{F}.,'c::null \text{ val}) \\

where eval-extract X f = (\lambda \tau. \text{ case } X \tau \text{ of} \\
\quad \bot \Rightarrow \text{ invalid } \tau \quad (* \text{ exception propagation } *) \\
\quad \mid \bot \Rightarrow \text{ invalid } \tau \quad (* \text{ dereferencing null pointer } *) \\
\quad \mid \text{ oid } \Rightarrow f \text{ (oid-of } \text{ obj } \text{ ) } \tau)


definition deref-oid_{\text{Person}} :: (\mathcal{A} \text{ state } \times \mathcal{A} \text{ state } \Rightarrow \mathcal{A} \text{ state}) \\

\Rightarrow (\text{type}_{\text{Person}} \Rightarrow (\mathcal{A},'c::null)\text{val}) \\
\Rightarrow \text{ oid} \\
\Rightarrow (\mathcal{A},'c::null)\text{val} \\

where deref-oid_{\text{Person}} \text{ fst-snd } f \text{ oid } = (\lambda \tau. \text{ case } (\text{heap (fst-snd } \tau)) \text{ oid of} \\
\quad \mid \text{ in}_{\text{Person}} \text{ obj } \Rightarrow f \text{ obj } \tau \\
\quad \mid \text{- } \Rightarrow \text{ invalid } \tau)


definition deref-oid_{\text{OclAny}} :: (\mathcal{A} \text{ state } \times \mathcal{A} \text{ state } \Rightarrow \mathcal{A} \text{ state})
⇒ (type\textsubscript{OclAny} ⇒ (A, 'c::null)val)
⇒ oid
⇒ (A, 'c::null)val

where deref-oid\textsubscript{OclAny} fst-snd f oid = (λτ. case (heap (fst-snd τ)) oid of
  [ inOclAny obj ] ⇒ f obj τ
  | - ⇒ invalid τ)

pointer undefined in state or not referencing a type conform object representation

\textbf{definition} select\textsubscript{OclAny}\textdecimalsym{ANY} f = (λ X. case X of
  (mk\textsubscript{OclAny} - ⊥) ⇒ null
  | (mk\textsubscript{OclAny} - [any]) ⇒ f (λx - [[x]]) any)

\textbf{definition} select\textsubscript{Person}\textdecimalsym{BOSS} f = (λ X. case X of
  (mk\textsubscript{Person} - - ⊥) ⇒ null (* object contains null pointer *)
  | (mk\textsubscript{Person} - - [boss]) ⇒ f (λx - [[x]]) boss)

\textbf{definition} select\textsubscript{Person}\textdecimalsym{SALARY} f = (λ X. case X of
  (mk\textsubscript{Person} - - ⊥) ⇒ null
  | (mk\textsubscript{Person} - - [salary]) ⇒ f (λx - [[x]]) salary)

\textbf{definition} in-pre-state = fst
\textbf{definition} in-post-state = snd

\textbf{definition} reconst-basetype = (λ convert x. convert x)

\textbf{definition} dot\textsubscript{OclAny}\textdecimalsym{ANY} :: OclAny ⇒ - ((1(-).any) 50)
where (X).any = eval-extract X
  (deref-oid\textsubscript{OclAny} in-post-state
   (select\textsubscript{OclAny}\textdecimalsym{ANY}
    reconst-basetype))

\textbf{definition} dot\textsubscript{Person}\textdecimalsym{BOSS} :: Person ⇒ Person ((1(-).boss) 50)
where (X).boss = eval-extract X
  (deref-oid\textsubscript{Person} in-post-state
   (select\textsubscript{Person}\textdecimalsym{BOSS}
    (deref-oid\textsubscript{Person} in-post-state)))

\textbf{definition} dot\textsubscript{Person}\textdecimalsym{SALARY} :: Person ⇒ Integer ((1(-).salary) 50)
where (X).salary = eval-extract X
  (deref-oid\textsubscript{Person} in-post-state
   (select\textsubscript{Person}\textdecimalsym{SALARY}
    reconst-basetype))

\textbf{definition} dot\textsubscript{OclAny}\textdecimalsym{ANY}-at-pre :: OclAny ⇒ - ((1(-).any@pre) 50)
where (X).any@pre = eval-extract X
  (deref-oid\textsubscript{OclAny} in-pre-state
(selectOclAny\(\text{ANY}\)
reconst-basetype))

**definition** dot\_{Person}BOSS-at-pre\::\ Person \Rightarrow \ Person \ ((1\ (-).boss@pre) 50)

where \((X).boss@pre = \text{eval-extract} X\)
(deref-oid\_{Person} in-pre-state
(select\_{Person}BOSS
(deref-oid\_{Person} in-pre-state)))

**definition** dot\_{Person}SALARY-at-pre\::\ Person \Rightarrow \ Integer \ ((1\ (-).salary@pre) 50)

where \((X).salary@pre = \text{eval-extract} X\)
(deref-oid\_{Person} in-pre-state
(select\_{Person}SALARY
reconst-basetype))

**lemmas** \[\text{simpl} =

dotOclAny\(\text{ANY}\)-def
dot\_{Person}BOSS-def
dot\_{Person}SALARY-def
dotOclAny\(\text{ANY}\)-at-pre-def
dot\_{Person}BOSS-at-pre-def
dot\_{Person}SALARY-at-pre-def\]

**Context Passing**

**lemmas** \[\text{simpl} = \text{eval-extract-def}\]

**lemma** \(\text{cp-dotOclAny\(\text{ANY}\)}\): \((X).any) \tau = ((\lambda . X \tau).any) \tau \langle \text{proof} \rangle

**lemma** \(\text{cp-dotPersonBOSS}\): \((X).boss) \tau = ((\lambda . X \tau).boss) \tau \langle \text{proof} \rangle

**lemma** \(\text{cp-dotPersonSALARY}\): \((X).salary) \tau = ((\lambda . X \tau).salary) \tau \langle \text{proof} \rangle

**lemma** \(\text{cp-dotOclAny\(\text{ANY}\)}\)-at-pre\::\ ((X).any@pre) \tau = ((\lambda . X \tau).any@pre) \tau \langle \text{proof} \rangle

**lemma** \(\text{cp-dotPersonBOSS}\)-at-pre\::\ ((X).boss@pre) \tau = ((\lambda . X \tau).boss@pre) \tau \langle \text{proof} \rangle

**lemma** \(\text{cp-dotPersonSALARY}\)-at-pre\::\ ((X).salary@pre) \tau = ((\lambda . X \tau).salary@pre) \tau \langle \text{proof} \rangle
lemmas cp-dot_PersonSALARY-one [simp, intro]=
cp-dot_PersonSALARY-I[THEN allI[THEN allI],
of λ X -. X λ - τ. τ, THEN cpII]

lemmas cp-dot_PersonSALARY-at-pre-one [simp, intro]=
cp-dot_PersonSALARY-at-pre[I[THEN allI[THEN allI],
of λ X -. X λ - τ. τ, THEN cpII]

Execution with Invalid or Null as Argument

lemma dot_OclAnyANY-nullstrict [simp]: (null).any = invalid
  ⟨proof⟩
lemma dot_OclAnyANY-at-pre-nullstrict [simp]: (null).any@pre = invalid
  ⟨proof⟩
lemma dot_OclAnyANY-strict [simp]: (invalid).any = invalid
  ⟨proof⟩
lemma dot_OclAnyANY-at-pre-strict [simp]: (invalid).any@pre = invalid
  ⟨proof⟩

lemma dot_PersonBOSS-nullstrict [simp]: (null).boss = invalid
  ⟨proof⟩
lemma dot_PersonBOSS-at-pre-nullstrict [simp]: (null).boss@pre = invalid
  ⟨proof⟩
lemma dot_PersonBOSS-strict [simp]: (invalid).boss = invalid
  ⟨proof⟩
lemma dot_PersonBOSS-at-pre-strict [simp]: (invalid).boss@pre = invalid
  ⟨proof⟩

lemma dot_PersonSALARY-nullstrict [simp]: (null).salary = invalid
  ⟨proof⟩
lemma dot_PersonSALARY-at-pre-nullstrict [simp]: (null).salary@pre = invalid
  ⟨proof⟩
lemma dot_PersonSALARY-strict [simp]: (invalid).salary = invalid
  ⟨proof⟩
lemma dot_PersonSALARY-at-pre-strict [simp]: (invalid).salary@pre = invalid
  ⟨proof⟩

7.1.9. A Little Infra-structure on Example States

The example we are defining in this section comes from the figure 7.2.

definition OclInt1000 (1000) where OclInt1000 = (λ - . [1000])
definition OclInt1200 (1200) where OclInt1200 = (λ - . [1200])
definition OclInt1300 (1300) where OclInt1300 = (λ - . [1300])
definition OclInt1800 (1800) where OclInt1800 = (λ - . [1800])
definition OclInt2600 (2600) where OclInt2600 = (λ - . [2600])
definition OclInt2900 (2900) where OclInt2900 = (λ - . [2900])
definition OclInt3200 (3200) where OclInt3200 = (λ - . [3200])
definition OclInt3500 (3500) where OclInt3500 = (λ - . [3500])
Figure 7.2.: (a) pre-state $\sigma_1$ and (b) post-state $\sigma'_1$.

definition $oid0 \equiv 0$
definition $oid1 \equiv 1$
definition $oid2 \equiv 2$
definition $oid3 \equiv 3$
definition $oid4 \equiv 4$
definition $oid5 \equiv 5$
definition $oid6 \equiv 6$
definition $oid7 \equiv 7$
definition $oid8 \equiv 8$
definition $person1 \equiv \text{mkPerson}(oid0 [1300] [oid1])$
definition $person2 \equiv \text{mkPerson}(oid1 [1800] [oid1])$
definition $person3 \equiv \text{mkPerson}(oid2 \text{None} \text{None})$
definition $person4 \equiv \text{mkPerson}(oid3 [2900] \text{None})$
definition $person5 \equiv \text{mkPerson}(oid4 [3500] \text{None})$
definition $person6 \equiv \text{mkPerson}(oid5 [2500] [oid6])$
definition $person7 \equiv \text{mkOclAny}(oid6 [[3200], [oid6]])$
definition $person8 \equiv \text{mkOclAny}(oid7 \text{None})$
definition $person9 \equiv \text{mkPerson}(oid8 [0] \text{None})$
definition $\sigma_1 \equiv (\langle heap = \text{empty}(oid0 \mapsto \text{inPerson}(\text{mkPerson}(oid0 [1000] [oid1]))) (oid1 \mapsto \text{inPerson}(\text{mkPerson}(oid1 [1200] \text{None})) (+oid2*) (oid3 \mapsto \text{inPerson}(\text{mkPerson}(oid3 [2600] [oid4]))) (oid4 \mapsto \text{inPerson}(\text{person5}) (oid5 \mapsto \text{inPerson}(\text{mkPerson}(oid5 [2300] [oid3]))) (+oid6*) (+oid7*) (oid8 \mapsto \text{inPerson}(\text{person9}),$

assocs$_2 = \text{empty},$

assocs$_3 = \text{empty} \rangle)$
definition $\sigma'_1 \equiv (\langle heap = \text{empty}(oid0 \mapsto \text{inPerson}(\text{person1})$
(oid1 \mapsto \textit{inPerson} \textit{person2})
(oid2 \mapsto \textit{inPerson} \textit{person3})
(oid3 \mapsto \textit{inPerson} \textit{person4})
(*oid4*)
(oid5 \mapsto \textit{inPerson} \textit{person6})
(oid6 \mapsto \textit{inOclAny} \textit{person7})
(oid7 \mapsto \textit{inOclAny} \textit{person8})
(oid8 \mapsto \textit{inPerson} \textit{person9}),
\text{assocs}_2 = \text{empty},
\text{assocs}_3 = \text{empty}

\text{definition} \quad \sigma_0 \equiv \left\{ \begin{array}{c}
\text{heap} = \text{empty}, \text{assocs}_2 = \text{empty}, \text{assocs}_3 = \text{empty}
\end{array} \right\}

\text{lemma basic-}\tau\text{-wff}: WFF(\sigma_1, \sigma_1')

⟨proof⟩

\text{lemma [simp, code-unfold]}: \text{dom}(\text{heap} \sigma_1) = \{ \text{oid0, oid1, (*, oid2*) oid3, oid4, oid5(*, oid6, oid7*), oid8} \}

⟨proof⟩

\text{lemma [simp, code-unfold]}: \text{dom}(\text{heap} \sigma_1') = \{ \text{oid0, oid1, oid2, oid3, (*, oid4*) oid5, oid6, oid7, oid8} \}

⟨proof⟩

\text{definition} \quad X_{\text{Person}1} :: \text{Person} \equiv \lambda - \llbracket \text{person1} \rrbracket
\text{definition} \quad X_{\text{Person}2} :: \text{Person} \equiv \lambda - \llbracket \text{person2} \rrbracket
\text{definition} \quad X_{\text{Person}3} :: \text{Person} \equiv \lambda - \llbracket \text{person3} \rrbracket
\text{definition} \quad X_{\text{Person}4} :: \text{Person} \equiv \lambda - \llbracket \text{person4} \rrbracket
\text{definition} \quad X_{\text{Person}5} :: \text{Person} \equiv \lambda - \llbracket \text{person5} \rrbracket
\text{definition} \quad X_{\text{Person}6} :: \text{Person} \equiv \lambda - \llbracket \text{person6} \rrbracket
\text{definition} \quad X_{\text{Person}7} :: \text{OclAny} \equiv \lambda - \llbracket \text{person7} \rrbracket
\text{definition} \quad X_{\text{Person}8} :: \text{OclAny} \equiv \lambda - \llbracket \text{person8} \rrbracket
\text{definition} \quad X_{\text{Person}9} :: \text{Person} \equiv \lambda - \llbracket \text{person9} \rrbracket

\text{lemma [code-unfold]}: ((x :: \text{Person}) \mapsto y) = \text{StrictRefEq}\text{Object} x y \quad \langle proof⟩
\text{lemma [code-unfold]}: ((x :: \text{OclAny}) \mapsto y) = \text{StrictRefEq}\text{Object} x y \quad \langle proof⟩

\text{lemmas [simp, code-unfold]} =
OclAsType\text{OclAny}-\text{OclAny}
OclAsType\text{OclAny}-\text{Person}
OclAsType\text{Person}-\text{OclAny}
OclAsType\text{Person}-\text{Person}
OclIsTypeOf\text{OclAny}-\text{OclAny}
OclIsTypeOf\text{OclAny}-\text{Person}
OclIsTypeOf\text{Person}-\text{OclAny}
OclIsTypeOf\text{Person}-\text{Person}
OclIsKindOf\text{OclAny}-\text{OclAny}
OclIsKindOf\text{OclAny}-\text{Person}
\( OclIsKindOf \_person \_OclAny \)
\( OclIsKindOf \_person \_Person \)

\[
\begin{align*}
\text{value} \; \land s_{pre} \cdot (s_{pre}, \sigma_1) &\models (X_{\text{Person}1} \_salary \; \llq 1000) \\
\text{value} \; \land s_{pre} \cdot (s_{pre}, \sigma_1) &\models (X_{\text{Person}1} \_salary \; \geq 1300) \\
\text{value} \; \land s_{post} \cdot (\sigma_1, s_{post}) &\models (X_{\text{Person}1} \_salary \; \geq \; pre \; \geq 1300) \\
\text{value} \; \land s_{pre} \cdot (s_{pre}, \sigma_1) &\models (X_{\text{Person}1} \_boss \; \llq X_{\text{Person}1}) \\
\text{value} \; \land s_{pre} \cdot (s_{pre}, \sigma_1) &\models (X_{\text{Person}1} \_boss \_salary \; \geq 1800) \\
\text{value} \; \land s_{pre} \cdot (s_{pre}, \sigma_1) &\models (X_{\text{Person}1} \_boss \_boss \; \llq X_{\text{Person}1}) \\
\text{value} \; \land s_{pre} \cdot (s_{pre}, \sigma_1) &\models (X_{\text{Person}1} \_boss \_boss \_salary \; \geq X_{\text{Person}2}) \\
\text{value} \; \land s_{post} \cdot (\sigma_1, s_{post}) &\models (X_{\text{Person}1} \_boss \_pre \_salary \; \geq 1800) \\
\text{value} \; \land s_{post} \cdot (\sigma_1, s_{post}) &\models (X_{\text{Person}1} \_boss \_pre \_salary \_pre \; \geq 1200) \\
\text{value} \; \land s_{post} \cdot (\sigma_1, s_{post}) &\models (X_{\text{Person}1} \_boss \_pre \_salary \_pre \; \geq 1800) \\
\text{value} \; \land s_{post} \cdot (\sigma_1, s_{post}) &\models (X_{\text{Person}1} \_boss \_pre \_salary \_pre \; \geq X_{\text{Person}2}) \\
\text{value} \; \land s_{post} \cdot (\sigma_1, s_{post}) &\models (X_{\text{Person}1} \_boss \_pre \_boss \_pre \; \geq X_{\text{Person}2}) \\
\text{value} \; \land s_{post} \cdot (\sigma_1, s_{post}) &\models (X_{\text{Person}1} \_boss \_pre \_boss \_pre \_pre \; \geq \; null) \\
\text{value} \; \land s_{post} \cdot (\sigma_1, s_{post}) &\models \; not(v\; (X_{\text{Person}1} \_boss \_pre \_boss \_pre \_pre \_pre)) \\
\text{lemma} \; (\sigma_1, \sigma_1) &\models (X_{\text{Person}1} \_oclIsMaintained()) \\
\text{proof} \\
\text{lemma} \; \land s_{pre} \cdot s_{post} \cdot (s_{pre}, s_{post}) &\models ((X_{\text{Person}1} \_oclAsType(OclAny) \_oclAsType(Person)) \; \geq X_{\text{Person}1}) \\
\text{proof} \\
\text{value} \; \land s_{pre} \cdot s_{post} \cdot (s_{pre}, s_{post}) &\models (X_{\text{Person}1} \_oclIsKindOf(Person)) \\
\text{value} \; \land s_{pre} \cdot s_{post} \cdot (s_{pre}, s_{post}) &\models not(X_{\text{Person}1} \_oclIsKindOf(OclAny)) \\
\text{value} \; \land s_{pre} \cdot s_{post} \cdot (s_{pre}, s_{post}) &\models (X_{\text{Person}1} \_oclIsKindOf(OclAny)) \\
\text{value} \; \land s_{pre} \cdot s_{post} \cdot (s_{pre}, s_{post}) &\models not(X_{\text{Person}1} \_oclAsType(OclAny) \_oclAsType(OclAny)) \\
\text{value} \; \land s_{pre} \cdot s_{post} \cdot (s_{pre}, s_{post}) &\models (X_{\text{Person}2} \_salary \; \geq 1800) \\
\text{value} \; \land s_{post} \cdot (\sigma_1, s_{post}) &\models (X_{\text{Person}2} \_salary \_pre \; \geq 1200) \\
\text{value} \; \land s_{pre} \cdot (s_{pre}, \sigma_1) &\models (X_{\text{Person}2} \_boss \; \geq X_{\text{Person}2}) \\
\text{value} \; \land s_{post} \cdot (\sigma_1, s_{post}) &\models (X_{\text{Person}2} \_boss \_salary \_pre \; \geq 1200) \\
\text{value} \; \land s_{post} \cdot (\sigma_1, s_{post}) &\models (X_{\text{Person}2} \_boss \_pre \; \geq \; null) \\
\text{value} \; \land s_{post} \cdot (\sigma_1, s_{post}) &\models (X_{\text{Person}2} \_boss \_pre \; \geq \; X_{\text{Person}2}) \\
\text{value} \; \land s_{post} \cdot (\sigma_1, s_{post}) &\models (X_{\text{Person}2} \_boss \_pre \_boss \; \geq \; null) \\
\text{value} \; \land s_{post} \cdot (\sigma_1, s_{post}) &\models (X_{\text{Person}2} \_boss \_pre \_boss \_boss \; \geq \; X_{\text{Person}2}) \\
\text{value} \; \land s_{post} \cdot (\sigma_1, s_{post}) &\models not(v\; (X_{\text{Person}2} \_boss \_pre \_boss \_boss)) \\
\text{value} \; \land s_{post} \cdot (\sigma_1, s_{post}) &\models not(v\; (X_{\text{Person}2} \_salary \_pre \_pre)) \\
\text{lemma} \; (\sigma_1, \sigma_1) &\models (X_{\text{Person}2} \_oclIsMaintained()) \\
\text{proof} \\
\text{value} \; \land s_{pre} \cdot (s_{pre}, \sigma_1) &\models (X_{\text{Person}3} \_salary \; \geq \; null) \\
\text{value} \; \land s_{post} \cdot (\sigma_1, s_{post}) &\models not(v\; (X_{\text{Person}3} \_salary \_pre)) \\
\end{align*}
\]

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\[
\begin{align*}
\text{value } & \wedge s_{\text{pre}} . (s_{\text{pre}},s_1) \models (X_{\text{Person}}^3\text{.boss} \wedge \text{null}) \\
\text{value } & \wedge s_{\text{pre}} . (s_{\text{pre}},s_1) \models \neg(v(X_{\text{Person}}^3\text{.boss }\text{salary})) \\
\text{value } & \wedge s_{\text{post}} . (s_1,s_{\text{post}}) \models \neg(v(X_{\text{Person}}^3\text{.boss@pre})) \\
\text{lemma } & \wedge s_{\text{post}} . (s_1,s_{\text{post}}) \models (X_{\text{Person}}^3\text{.oclIsNew}()) \\
\langle \text{proof} \rangle & \\
\text{value } & \wedge s_{\text{post}} . (s_1,s_{\text{post}}) \models (X_{\text{Person}}^4\text{.boss@pre} \wedge X_{\text{Person}}^5) \\
\text{value } & \wedge s_{\text{post}} . (s_1,s_{\text{post}}) \models \neg(v(X_{\text{Person}}^4\text{.boss@pre}.\text{salary})) \\
\text{value } & \wedge s_{\text{pre}} . (s_{\text{pre}},s_1) \models \neg(v(X_{\text{Person}}^5\text{.boss})) \\
\text{lemma } & \wedge s_{\text{pre}} . (s_{\text{pre}},s_1) \models (X_{\text{Person}}^5\text{.oclIsDeleted}()) \\
\langle \text{proof} \rangle & \\
\text{value } & \wedge s_{\text{pre}} . (s_{\text{pre}},s_1) \models \neg(v(X_{\text{Person}}^5\text{.salary})) \\
\text{value } & \wedge s_{\text{post}} . (s_1,s_{\text{post}}) \models (X_{\text{Person}}^5\text{.salary@pre} \wedge X_{\text{Person}}^5) \\
\text{value } & \wedge s_{\text{pre}} . (s_{\text{pre}},s_1) \models \neg(v(X_{\text{Person}}^5\text{.salary})) \\
\text{lemma } & \wedge s_{\text{pre}} . (s_{\text{pre}},s_1) \models (X_{\text{Person}}^5\text{.oclIsNew}()) \\
\langle \text{proof} \rangle & \\
\text{value } & \wedge s_{\text{pre}} . (s_{\text{pre}},s_1) \models \neg(v(X_{\text{Person}}^6\text{.boss }\text{salary@pre})) \\
\text{value } & \wedge s_{\text{post}} . (s_1,s_{\text{post}}) \models (X_{\text{Person}}^6\text{.boss@pre} \wedge X_{\text{Person}}^4) \\
\text{value } & \wedge s_{\text{post}} . (s_1,s_{\text{post}}) \models (X_{\text{Person}}^6\text{.boss@pre}.\text{salary}@pre \wedge X_{\text{Person}}^4) \\
\text{value } & \wedge s_{\text{post}} . (s_1,s_{\text{post}}) \models (X_{\text{Person}}^6\text{.boss@pre}.\text{salary}@pre \wedge X_{\text{Person}}^4) \\
\text{lemma } & \wedge s_{\text{post}} . (s_1,s_{\text{post}}) \models (X_{\text{Person}}^6\text{.oclIsNew}()) \\
\langle \text{proof} \rangle & \\
\text{value } & \wedge s_{\text{pre}} . s_{\text{post}} . (s_{\text{pre}},s_{\text{post}}) \models v(X_{\text{Person}}^7\text{.oclAsType(Person)}) \\
\text{value } & \wedge s_{\text{post}} . (s_1,s_{\text{post}}) \models \neg(v(X_{\text{Person}}^7\text{.oclAsType(Person)} . \text{boss@pre})) \\
\text{lemma } & \wedge s_{\text{pre}} s_{\text{post}} . (s_{\text{pre}},s_{\text{post}}) \models (X_{\text{Person}}^7\text{.oclAsType(Person)} . \text{oclAsType(OclAny)} . \text{oclAsType(Person)})) \\
& \models (X_{\text{Person}}^7\text{.oclAsType(Person)})) \\
\langle \text{proof} \rangle & \\
\text{lemma } & (s_1,s_1) \models (X_{\text{Person}}^7\text{.oclIsNew}()) \\
\langle \text{proof} \rangle & \\
\text{value } & \wedge s_{\text{pre}} s_{\text{post}} . (s_{\text{pre}},s_{\text{post}}) \models (X_{\text{Person}}^8 \leftrightarrow X_{\text{Person}}^7) \\
\text{value } & \wedge s_{\text{pre}} s_{\text{post}} . (s_{\text{pre}},s_{\text{post}}) \models \neg(v(X_{\text{Person}}^8\text{.oclAsType(Person)})) \\
\text{value } & \wedge s_{\text{pre}} s_{\text{post}} . (s_{\text{pre}},s_{\text{post}}) \models (X_{\text{Person}}^8\text{.oclIsTypeOf(OclAny)})) \\
\text{value } & \wedge s_{\text{pre}} s_{\text{post}} . (s_{\text{pre}},s_{\text{post}}) \models (X_{\text{Person}}^8\text{.oclIsTypeOf(Person)})) \\
\text{value } & \wedge s_{\text{pre}} s_{\text{post}} . (s_{\text{pre}},s_{\text{post}}) \models \neg(X_{\text{Person}}^8\text{.oclIsKindOf(Person)})) \\
\end{align*}
\]
\textbf{lemma} \(\sigma\text{-modifiedonly}: (\sigma_1, \sigma_1') \models (\text{Set} \{ X_{\text{Person}1}.\text{oclAsType}(\text{OclAny}) \), 
X_{\text{Person}2}.\text{oclAsType}(\text{OclAny}) 
\), X_{\text{Person}3}.\text{oclAsType}(\text{OclAny})*) 
\), X_{\text{Person}4}.\text{oclAsType}(\text{OclAny}) 
\), X_{\text{Person}5}.\text{oclAsType}(\text{OclAny})*) 
\), X_{\text{Person}6}.\text{oclAsType}(\text{OclAny}) 
\), X_{\text{Person}7}.\text{oclAsType}(\text{OclAny})*) 
\), X_{\text{Person}8}.\text{oclAsType}(\text{OclAny})*) 
\), X_{\text{Person}9}.\text{oclAsType}(\text{OclAny})*) \models \text{oclIsModifiedOnly}()) 
\langle \text{proof} \rangle 
\text{lemma} \ (\sigma_1, \sigma_1') \models ((X_{\text{Person}9}.\text{oclAsType}(\text{OclAny})) \oplus \text{pre} (\lambda x. [\text{OclAsType}_{\text{Person}}x]) \triangleq X_{\text{Person}9}) 
\langle \text{proof} \rangle 
\text{lemma} \ (\sigma_1, \sigma_1') \models ((X_{\text{Person}9}.\text{oclAsType}(\text{OclAny})) \oplus \text{post} (\lambda x. [\text{OclAsType}_{\text{Person}}x]) \triangleq X_{\text{Person}9}) 
\langle \text{proof} \rangle 
\text{lemma} \ (\sigma_1, \sigma_1') \models (((X_{\text{Person}9}.\text{oclAsType}(\text{OclAny})) \oplus \text{pre} (\lambda x. [\text{OclAsType}_{\text{OclAny}}x])) \triangleq 
((X_{\text{Person}9}.\text{oclAsType}(\text{OclAny})) \oplus \text{post} (\lambda x. [\text{OclAsType}_{\text{OclAny}}x]))) 
\langle \text{proof} \rangle 
\text{lemma} \ \text{perm}-\sigma_1': \sigma_1' = \emptyset \\text{heap} = \text{empty} 
\begin{align*} 
\text{oid8} & \mapsto \text{in}_{\text{Person}} \text{person9} \\
\text{oid7} & \mapsto \text{in}_{\text{OclAny}} \text{person8} \\
\text{oid6} & \mapsto \text{in}_{\text{OclAny}} \text{person7} \\
\text{oid5} & \mapsto \text{in}_{\text{Person}} \text{person6} \\
\text{oid}4 & \mapsto \\
\text{oid3} & \mapsto \text{in}_{\text{Person}} \text{person4} \\
\text{oid2} & \mapsto \text{in}_{\text{Person}} \text{person3} \\
\text{oid1} & \mapsto \text{in}_{\text{Person}} \text{person2} \\
\text{oid0} & \mapsto \text{in}_{\text{Person}} \text{person1} \\
\text{assoc}2 = & \text{assoc}2 \sigma_1' \\
\text{assoc}3 = & \text{assoc}3 \sigma_1' \\
\end{align*} 
\langle \text{proof} \rangle 
\textbf{declare} \ const-\text{ss} \ \text{[simp]} 
\textbf{lemma} \ \wedge \sigma_1. 
(\sigma_1, \sigma_1') \models (\text{Person} .\text{allInstances}()) \equiv \text{Set} \{ X_{\text{Person}1}, X_{\text{Person}2}, X_{\text{Person}3}, X_{\text{Person}4}(\text{\textasteriskcentered}), X_{\text{Person}5}(\text{\textasteriskcentered}), X_{\text{Person}6}, X_{\text{Person}7}.\text{oclAsType}(\text{Person})(\text{\textasteriskcentered}, X_{\text{Person}8}(\text{\textasteriskcentered}), X_{\text{Person}9} ) \} 
\langle \text{proof} \rangle 
\textbf{lemma} \ \wedge \sigma_1. 
(\sigma_1, \sigma_1') \models (\text{OclAny} .\text{allInstances}()) \equiv \text{Set} \{ X_{\text{Person}1}.\text{oclAsType}(\text{OclAny}), X_{\text{Person}2}.\text{oclAsType}(\text{OclAny}), X_{\text{Person}3}.\text{oclAsType}(\text{OclAny}), X_{\text{Person}4}.\text{oclAsType}(\text{OclAny}), X_{\text{Person}5}.\text{oclAsType}(\text{OclAny}), 
X_{\text{Person}6}.\text{oclAsType}(\text{OclAny}), 
X_{\text{Person}7}.\text{oclAsType}(\text{OclAny}), 
X_{\text{Person}8}.\text{oclAsType}(\text{OclAny}), 
X_{\text{Person}9}.\text{oclAsType}(\text{OclAny}) \}
\langle \text{proof} \rangle
7.2. The Employee Design Model (OCL)

theory Employee-DesignModel-OCLPart

imports Employee-DesignModel-UMLPart

begin

7.2.1. Standard State Infrastructure

Ideally, these definitions are automatically generated from the class model.

7.2.2. Invariant

These recursive predicates can be defined conservatively by greatest fix-point constructions—automatically. See [4, 6] for details. For the purpose of this example, we state them as axioms here.

axiomatization inv-Person :: Person ⇒ Boolean

where A : (τ |= (δ self))  →  
  (τ |= inv-Person(self)) = 
  (τ |= (self .boss = null)) ∨ 
  (τ |= (self .boss <> null) ∧ (τ |= ((self .salary) '≤ (self .boss .salary)))) ∧ 
  (τ |= (inv-Person(self .boss))))

axiomatization inv-Person-at-pre :: Person ⇒ Boolean

where B : (τ |= (δ self))  →  
  (τ |= inv-Person-at-pre(self)) = 
  (τ |= (self .boss@pre = null)) ∨ 
  (τ |= (self .boss@pre <> null) ∧ 
  (τ |= ((self .boss@pre .salary@pre '≤ self .salary@pre)) ∧ 
  (τ |= (inv-Person-at-pre(self .boss@pre)))))

A very first attempt to characterize the axiomatization by an inductive definition - this can not be the last word since too weak (should be equality!)

coinductive inv :: Person ⇒ (∃)st ⇒ bool where

(τ |= (δ self))  →  (τ |= (self .boss = null)) ∨ 
  (τ |= (self .boss <> null) ∧ (τ |= (self .boss .salary '≤ self .salary)) ∧ 

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( \text{inv}(self\.boss)) \tau \right) \Rightarrow ( \text{inv} self \tau)

7.2.3. The Contract of a Recursive Query

The original specification of a recursive query:

context Person::contents(): Set (Integer)
post: \text{result} = \begin{cases} \text{Set}\{i\} & \text{if} self\.boss = \text{null} \\ \text{self\.boss\.contents()–including}(i) & \text{else} \end{cases}

consts dot-contents :: Person \Rightarrow \text{Set-Integer} ((1(-).contents('')) 50)

axiomatization where dot-contents-def:
(\tau \models (\text{self}.\text{contents}() \triangleq \text{result})) =
(\text{if} (\delta \text{self}) \tau = \text{true} \tau \\
\text{then} (\tau \models \text{true} \land \\
(\tau \models (\text{result} \triangleq \text{if} (\text{self}.\text{boss} \cong \text{null}) \\
\text{then} \text{Set}\{\text{self}.\text{salary}\} \\
\text{else} \text{self}.\text{boss}.\text{contents()–including}(\text{self}.\text{salary}) \\
\text{endif}) \\
\text{else} \tau \models \text{result} \triangleq \text{invalid}))

consts dot-contents-AT-pre :: Person \Rightarrow \text{Set-Integer} ((1(-).\text{contents}@pre('')) 50)

axiomatization where dot-contents-AT-pre-def:
(\tau \models (\text{self}.\text{contents}@pre() \triangleq \text{result})) =
(\text{if} (\delta \text{self}) \tau = \text{true} \tau \\
\text{then} \tau \models \text{true} \land \\
\tau \models (\text{result} \triangleq \text{if} (\text{self}.\text{boss}@pre \cong \text{null} \text{ post \text{post})} \\
\text{then} \text{Set}\{\text{self}.\text{salary}@pre\} \\
\text{else} \text{self}.\text{boss}@pre.\text{contents}@pre()–including(\text{self}.\text{salary}@pre) \\
\text{endif}) \\
\text{else} \tau \models \text{result} \triangleq \text{invalid})

These \@pre variants on methods are only available on queries, i.e., operations without side-effect.

7.2.4. The Contract of a Method

The specification in high-level OCL input syntax reads as follows:

context Person::insert(x: Integer)
post: contents(): Set (Integer)
\text{contents()} = \text{contents}@pre()–including(x)

consts dot-insert :: Person \Rightarrow \text{Integer} \Rightarrow \text{Void} ((1(-).insert('')) 50)
axiomatization where dot-insert-def:

\[(\tau \models ((\text{self}).\text{insert}(x) \triangleq \text{result})) = \]
\[(\text{if } (\delta \text{ self}) \tau = \text{true } \tau \land (\nu x) \tau = \text{true } \tau \]
\[\text{then } \tau \models \text{true } \land \]
\[\quad \tau \models ((\text{self}).\text{contents}() \triangleq (\text{self}).\text{contents}@[\text{pre}]->\text{including}(x)) \]
\[\text{else } \tau \models ((\text{self}).\text{insert}(x) \triangleq \text{invalid}) \]
\end
Part IV.

Conclusion
8. Conclusion

8.1. Lessons Learned and Contributions

We provided a typed and type-safe shallow embedding of the core of UML [31, 32] and OCL [33]. Shallow embedding means that types of OCL were injectively, i.e., mapped by the embedding one-to-one to types in Isabelle/HOL [27]. We followed the usual methodology to build up the theory uniquely by conservative extensions of all operators in a denotational style and to derive logical and algebraic (execution) rules from them; thus, we can guarantee the logical consistency of the library and instances of the class model construction, i.e., closed-world object-oriented datatype theories, as long as it follows the described methodology. Moreover, all derived execution rules are by construction type-safe (which would be an issue, if we had chosen to use an object universe construction in Zermelo-Fraenkel set theory as an alternative approach to subtyping.). In more detail, our theory gives answers and concrete solutions to a number of open major issues for the UML/OCL standardization:

1. the role of the two exception elements invalid and null, the former usually assuming strict evaluation while the latter ruled by non-strict evaluation.

2. the functioning of the resulting four-valued logic, together with safe rules (for example foundation9 − foundation12 in Section 3.5.2) that allow a reduction to two-valued reasoning as required for many automated provers. The resulting logic still enjoys the rules of a strong Kleene Logic in the spirit of the Amsterdam Manifesto [19].

3. the complicated life resulting from the two necessary equalities: the standard’s “strict weak referential equality” as default (written \(- \equiv -\) throughout this document) and the strong equality (written \(- \equiv -\), which follows the logical Leibniz principle that “equals can be replaced by equals.” Which is not necessarily the case if invalid or objects of different states are involved.

4. a type-safe representation of objects and a clarification of the old idea of a one-to-one correspondence between object representations and object-id’s, which became a state invariant.

5. a simple concept of state-framing via the novel operator \(-\rightarrow\text{oclIsModifiedOnly}()\) and its consequences for strong and weak equality.

\[\text{Our two examples of Employee DesignModel (see Chapter 7) sketch how this construction can be captured by an automated process.}\]
6. a semantic view on subtyping clarifying the role of static and dynamic type (aka apparent and actual type in Java terminology), and its consequences for casts, dynamic type-tests, and static types.

7. a semantic view on path expressions, that clarify the role of invalid and null as well as the tricky issues related to de-referentiation in pre- and post state.

8. an optional extension of the OCL semantics by infinite sets that provide means to represent “the set of potential objects or values” to state properties over them (this will be an important feature if OCL is intended to become a full-blown code annotation language in the spirit of JML [25] for semi-automated code verification, and has been considered desirable in the Aachen Meeting [15]).

Moreover, we managed to make our theory in large parts executable, which allowed us to include mechanically checked value-statements that capture numerous corner-cases relevant for OCL implementors. Among many minor issues, we thus pin-pointed the behavior of null in collections as well as in casts and the desired isKindOf-semantics of allInstances().

8.2. Lessons Learned

While our paper and pencil arguments, given in [13], turned out to be essentially correct, there had also been a lesson to be learned: If the logic is not defined as a Kleene-Logic, having a structure similar to a complete partial order (CPO), reasoning becomes complicated: several important algebraic laws break down which makes reasoning in OCL inherent messy and a semantically clean compilation of OCL formulae to a two-valued presentation, that is amenable to animators like KodKod [36] or SMT-solvers like Z3 [20] completely impractical. Concretely, if the expression not(null) is defined invalid (as is the case in the present standard [33]), than standard involution does not hold, i.e., not(not(A)) = A does not hold universally. Similarly, if null and null is invalid, then not even idempotence X and X = X holds. We strongly argue in favor of a lattice-like organization, where null represents “more information” than invalid and the logical operators are monotone with respect to this semantical “information ordering.”

A similar experience with prior paper and pencil arguments was our investigation of the object-oriented data-models, in particular path-expressions [16]. The final presentation is again essentially correct, but the technical details concerning exception handling lead finally to a continuation-passing style of the (in future generated) definitions for accessors, casts and tests. Apparently, OCL semantics (as many other “real” programming and specification languages) is meanwhile too complex to be treated by informal arguments solely.

Featherweight OCL makes several minor deviations from the standard and showed how the previous constructions can be made correct and consistent, and the DNF-normalization as well as δ-closure laws (necessary for a transition into a two-valued
presentation of OCL specifications ready for interpretation in SMT solvers (see [14] for
details)) are valid in Featherweight OCL.

8.3. Conclusion and Future Work

Featherweight OCL concentrates on formalizing the semantics of a core subset of OCL in
general and in particular on formalizing the consequences of a four-valued logic (i.e., OCL
versions that support, besides the truth values true and false also the two exception
values invalid and null).

In the following, we outline the necessary steps for turning Featherweight OCL into a
fully fledged tool for OCL, e.g., similar to HOL-OCL as well as for supporting test case
generation similar to HOL-TestGen [9]. There are essentially five extensions necessary:

• extension of the library to support all OCL data types, e.g., OrderedSet(T) or
  Sequence(T). This formalization of the OCL standard library can be used for
  checking the consistency of the formal semantics (known as “Annex A”) with the
  informal and semi-formal requirements in the normative part of the OCL standard.

• development of a compiler that compiles a textual or CASE tool representation
  (e.g., using XMI or the textual syntax of the USE tool [35]) of class models. Such
  compiler could also generate the necessary casts when converting standard OCL
to Featherweight OCL as well as providing “normalizations” such as converting
  multiplicities of class attributes to into OCL class invariants.

• a setup for translating Featherweight OCL into a two-valued representation as de-
  scribed in [14]. As, in real-world scenarios, large parts of UML/OCL specifications
  are defined (e.g., from the default multiplicity 1 of an attributes x, we can directly
  infer that for all valid states x is neither invalid nor null), such a translation
  enables an efficient test case generation approach.

• a setup in Featherweight OCL of the Nitpick animator [3]. It remains to be shown
  that the standard, Kodkod [36] based animator in Isabelle can give a similar quality
  of animation as the OCLexec Tool [24]

• a code-generator setup for Featherweight OCL for Isabelle’s code generator. For
  example, the Isabelle code generator supports the generation of F#, which would
  allow to use OCL specifications for testing arbitrary .net-based applications.

The first two extensions are sufficient to provide a formal proof environment for OCL
2.5 similar to HOL-OCL while the remaining extensions are geared towards increasing
the degree of proof automation and usability as well as providing a tool-supported test
methodology for UML/OCL.

Our work shows that developing a machine-checked formal semantics of recent OCL
standards still reveals significant inconsistencies—even though this type of research is
not new. In fact, we started our work already with the 1.x series of OCL. The reasons
for this ongoing consistency problems of OCL standard are manifold. For example, the

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consequences of adding an additional exception value to OCL 2.2 are widespread across the whole language and many of them are also quite subtle. Here, a machine-checked formal semantics is of great value, as one is forced to formalize all details and subtleties. Moreover, the standardization process of the OMG, in which standards (e.g., the UML infrastructure and the OCL standard) that need to be aligned closely are developed quite independently, are prone to ad-hoc changes that attempt to align these standards. And, even worse, updating a standard document by voting on the acceptance (or rejection) of isolated text changes does not help either. Here, a tool for the editor of the standard that helps to check the consistency of the whole standard after each and every modifications can be of great value as well.
Bibliography


